# SoPlex 

## and

Zimpl

## Thorsten Koch, ZIB

## $\min c^{\top} x$

## subject to $A x \geqslant b$

with $\quad c \in \mathbb{R}^{n}$
$x \in \mathbb{R}^{n}$
$A \in \mathbb{R}^{k \times n}$
$b \in \mathbb{R}^{k}$
and $k \geqslant n$ and $A$ has full rank

## $\max c^{\top} x$

## subject to $A x \leqslant b$


with $c \in \mathbb{R}^{n}, x \in \mathbb{R}^{n}, A \in \mathbb{R}^{k \times n}, b \in \mathbb{R}^{k}$.

## Linear program example

(4)


$$
\min -x-y
$$

subject to

$$
\begin{align*}
x_{1}-2 x_{2} & \geqslant-2  \tag{1}\\
-6 x_{1}-2 x_{2} & \geqslant-9  \tag{2}\\
-x_{1}+2 x_{2} & \geqslant-1 \\
x_{1} & \geqslant 0  \tag{4}\\
x_{2} & \geqslant 0
\end{align*}
$$

## The Sequential object-oriented simplex class library

- Implementation of the revised simplex algorithm
- Primal and dual solving routines
- Row and column based basis representation
- Can solve instances with a million non-zero elements
- C++ class library and standalone program
- Very portable: compiles with C++ compilers from

GNU, Compaq, Intel, SUN, HP, SGI, IBM, and even M\$

- Licensed to several commercial companies
- Free for non commercial academic use


## SoPlex is a linear program solver.

## It can not solve integer programs,

 i.e. $x \in \mathbb{Z}^{n}$.For this we have SCIP, please wait for the next talk.

- Currently one of the top free LP-Solvers
- Initially developed in 1996 by Roland Wunderling, as part of his PhD thesis Paralleler und Objektorientierter Simplex-Algorithmus
- SoPlex is under continous development
- While slower than, for example, CPLEX, adaquate for many tasks

Available at
http://www.zib.de/Optimization/Software/Soplex

## Zuse Institute Mathematical Programming Language

- Algebraic modeling language, like AMPL, GAMS, etc.
- Transforms mathematical descriptions of linear mixed integer models into solver input
- Easy to use
- Has been used in several industry projects and lectures
- Can generate models with more than 10 million variables
- Released under GNU GPL

A Zimpl model consists of

- Sets
- Parameters
- Variables
- Objective
- Constraints

Let $G=(V, E)$ be a complete graph, with $V$ being the set of cities and $E$ being the set of links between the cities. Introducing binary variables $x_{i j}$ for each $(i, j) \in E$ indicating if edge $(i, j)$ is part of the tour, the TSP can be written as:

$$
\begin{array}{cl}
\min \sum_{(i, j) \in E} d_{i j} x_{i j} & \text { subject to } \\
\sum_{(i, j) \in \delta_{v}} x_{i j}=2 & \text { for all } v \in V \\
\sum_{(i, j) \in E(U)} x_{i j} \leqslant|U|-1 & \text { for all } U \subseteq V, \emptyset \neq U \neq V \\
x_{i j} \in\{0,1\} & \text { for all }(i, j) \in E
\end{array}
$$

## Zimpl example: data

The data is read in from a file that lists for each city the name and the $x$ and $y$ coordinates. Distances between cities are assumed Euclidean.

| \# City | X | Y |
| :--- | ---: | ---: |
| Berlin | 5251 | 1340 |
| Frankfurt | 5011 | 864 |
| Leipzig | 5133 | 1237 |
| Heidelberg | 4941 | 867 |
| Karlsruhe | 4901 | 840 |
| Hamburg | 5356 | 998 |
| Bayreuth | 4993 | 1159 |
| Trier | 4974 | 668 |
| Hannover | 5237 | 972 |


| Stuttgart | 4874 | 909 |
| :--- | ---: | ---: |
| Passau | 4856 | 1344 |
| Augsburg | 4833 | 1089 |
| Koblenz | 5033 | 759 |
| Dortmund | 5148 | 741 |
| Bochum | 5145 | 728 |
| Duisburg | 5142 | 679 |
| Wuppertal | 5124 | 715 |
| Essen | 5145 | 701 |
| Jena | 5093 | 1158 |



The resulting linear program has
171 variables,
239,925 constraints, and
22,387,149 non-zero entries
in the constraint matrix, giving an MPS-file size of 936 mb .
An optimal tour for the data on the previous slide is Berlin, Hamburg, Hannover, Dortmund, Bochum, Wuppertal, Essen, Duisburg, Trier, Koblenz, Frankfurt, Heidelberg, Karlsruhe, Stuttgart, Augsburg, Passau, Bayreuth, Jena, Leipzig, Berlin.

- Modeling languages make it much easier to rapidly experiment with models and ideas
- The reproducibility of the results is increased
- Higher solver independency
- Solver dependent transformation of special functions are possible

For $a, b, c \in \mathbb{Z}, a \in[-15,15], b \in[-10,20], c \in[-20,10]$, maximize $5 a+3 b+c$ subject to:

$$
\text { if }(a \neq b \text { and }(|a-b|=3 c \text { or } c-a \geqslant 0))
$$

then $a+b+c \geqslant 7$
else $a+b \leqslant 1$
can be formulated as an IP. But it is rather difficult...

For $a, b, c \in \mathbb{Z}, a \in[-15,15], b \in[-10,20], c \in[-20,10]$, maximize $5 a+3 b+c$ subject to:

$$
\text { if }(a \neq b \text { and }(|a-b|=3 c \text { or } c-a \geqslant 0))
$$

then $a+b+c \geqslant 7$
else $a+b \leqslant 1$
var a integer >= -15 <= 15;
var b integer >= -10 <= 20;
var c integer >= -20 <= 10;
maximize obj: 5 * a + 3 * b + c;
subto ci: vif (a!=b and (vabs(a-b) == $3 * c$ or $c-a \quad>=0)$ )
then $a+b+c>=7$
else $a+b<=1$ end;

## Extended functions: result



## www.zib.de/koch/zimpl

## THANK YOU!

QUESTIONS?

