

Extending SCIP for solving MIQCPs

Timo Berthold Ambros M. Gleixner Stefan Heinz Thorsten Koch Stefan Vigerske

Mixed-Integer Quadratically Constrained Programming (MIQCP)

We consider optimization problems of the form

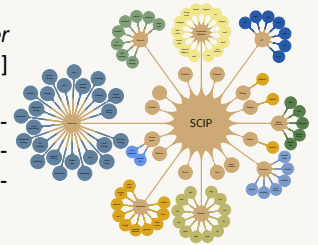
$$\begin{aligned} & \text{minimize} && b_0^T x + c_0 \\ & \text{subject to} && x^T A_j x + b_j^T x + c_j \leq 0 \quad j = 1, \dots, m \\ & && x_i \in \mathbb{Z} \quad \forall i \in I, \end{aligned}$$

where $A_j \in \mathbb{Q}^{n \times n}$, $b_j \in \mathbb{Q}^n$, $c_j \in \mathbb{Q}$, $j = 0, \dots, m$, and $I \subseteq \{1, \dots, n\}$.

In general, A_j does not need to be positive semidefinite.

Algorithm: LP-based branch-and-cut

We extend the *Constraint Integer Programming* framework SCIP [1] by MINLP-specific plugins. For nonlinear constraints, we generate a suitable linear outer approximation and apply domain propagation and primal heuristics.



Currently implemented [5]

- ▷ reformulation for products with binary variables
- ▷ recognition of convex quadratic functions
- ▷ separation for convex and nonconvex quad. constraints
- ▷ domain propagation for quadratic constraints
- ▷ handling of second-order-cone (SOC) constraints
- ▷ local search heuristic: fix integers, solve *sub-QCP* locally
- ▷ interfaces to GAMS, MPS, and ZIMPL

Nonlinear RENS Heuristic

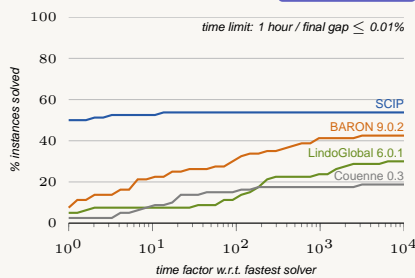
Relaxation Enforced Neighborhood Search heuristic [2]: fix integer variables that take integral value in optimal solution of LP relaxation and solve remaining *sub-MIQCP*

Undercover Heuristic [3]

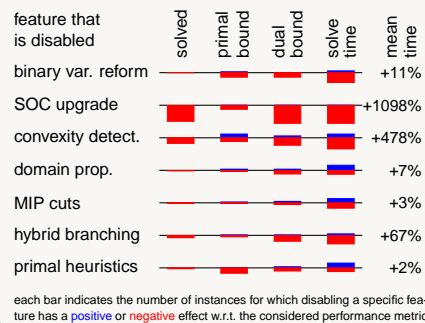
construct (and solve) *sub-MIP* by fixing as few nonlinear variables as possible to their value in an optimal solution of the LP or NLP relaxation

Computational Results

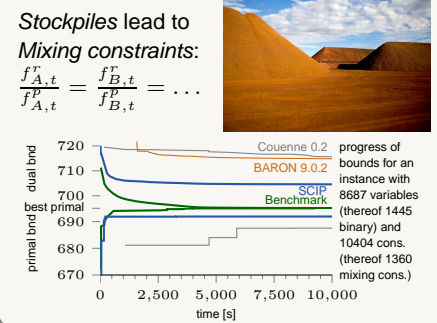
Benchmark



Impact of Single Components

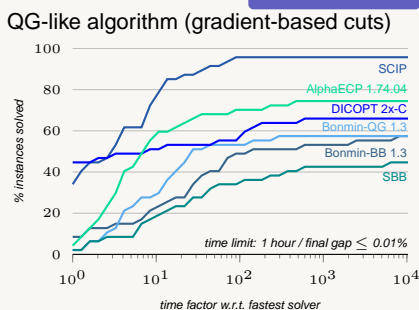


Application: Mine Prod. Scheduling [6]



Beyond MIQCP

Convex MINLP



Pseudo-Boolean Optimization [4]

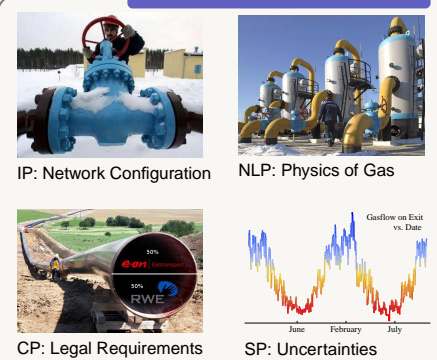
Pseudo-Boolean Constraints:

$$\sum_{j=1}^{t_i} a_{ij} \cdot \prod_{k \in I_{ij}} x_k \cdot \prod_{k \in \bar{I}_{ij}} (1 - x_k) \geq b_i$$

$$x_k \in \{0, 1\}, k \in I_{ij} \cup \bar{I}_{ij}$$

- ▷ replace multiplications by new variables and *AND-constraints*
- ▷ specialised separation and propagation for *AND-constraints* (small initial relax., separate strong ineq., always propagate)
- ▷ Winner in 3 out of 4 categories at Pseudo-Boolean Evaluation 2009

Application: Gas Transport



References

- [1] T. Achterberg. SCIP: Solving Constraint Integer Programs, *Mathematical Programming Computation* 1 (2009), pp. 1–41
- [2] T. Berthold. RENS – Relaxation Enforced Neighborhood Search, ZIB-Report 07-28
- [3] T. Berthold, A. Gleixner. Undercover – a primal heuristic for MINLP based on sub-MIPs generated by set covering, ZIB-Report 09-40
- [4] T. Berthold, S. Heinz, M.E. Pfetsch. Nonlinear pseudo-Boolean optimization: relaxation or propagation?, *Proc. of SAT 2009*, pp. 441–446
- [5] T. Berthold, S. Heinz, S. Vigerske. Extending a CIP framework to solve MIQCPs, ZIB-Report 09-23
- [6] A. Bley, A. Gleixner, T. Koch, S. Vigerske. Comparing MIQCP solvers to a specialised algorithm for mine production scheduling, ZIB-Report 09-32

Cooperations

