

§3 Universality

From now on assume:

A1 Either no $\tau < \theta$ is Woodin in an inner model or else V_θ is closed under #.

A2 Let $M \in V_\theta$ be a 1-small premouse and \mathcal{I} a normal iteration of M of length θ . Then \mathcal{I} has a cofinal branch.

A3 θ is a Mahlo cardinal.

Note A1 subsumes our earlier A0. A1 is known to imply that K^c is iterable below θ .

Note A2 holds if $V_\theta^\#$ exists or if θ is not Woodin in an inner model.

Note A3 is assumed in order to insure a large supply of inaccessible $\kappa < \theta$ which can serve as critical points for the extenders in the K^c construction.

We prove:

Thm 5 κ^c is universal wrt. 1-small
premise in \mathcal{V}_θ - i.e. if $N \in \mathcal{V}_\theta$ is a
1-small premise, then any coiteration
of N , κ^c will terminate below θ .

Note A failure of well foundedness is
considered a termination. Hence the
lemma would hold even without A1,
which is used to ensure an iteration
strategy for κ^c . The interesting
case occurs when both sides have
a normal iteration strategy, thus
preventing a failure of well found-
edness.

We now prove Thm 5. Suppose not,
let $Q \in \mathcal{V}_\theta$ be a 1-small premise
and let $\langle \gamma^Q, \gamma^K \rangle$ be a coiteration
of Q, κ^c of length θ . Let:

$$\gamma^Q = \langle \langle Q_i \rangle, \langle v_i \rangle, \langle \gamma_i^Q \rangle, \langle \pi_{ii}^Q \rangle, T^Q \rangle$$

$$\gamma^K = \langle \langle K_i \rangle, \langle v_i \rangle, \langle \gamma_i^K \rangle, \langle \pi_{ii}^K \rangle, T^K \rangle.$$

Making use of this we shall prove:

Lemma 5.1 There are a stationary $S \subset \Theta$ and a commutative system $\langle \pi_{\alpha\beta} \mid \alpha \leq \beta \text{ in } S \rangle$ s.t.

(i) α is inaccessible for $\alpha \in S$

(ii) Set $\tilde{K}_\alpha = (J_{\alpha^+}^E)^{K^c}$ ($\alpha \in S$). Then

$$\pi_{\alpha\beta} : \tilde{K}_\alpha \longrightarrow \sum_\nu \tilde{K}_\beta \text{ cofinally; } \pi_{\alpha\beta} \upharpoonright \alpha = \text{id};$$

$$\pi_{\alpha\beta}(\alpha) = \beta \text{ for } \alpha \leq \beta \text{ in } S.$$

proof.

Let b be a cofinal branch in \mathcal{Y}^Q . This has a transitive limit model Q_b and maps π_i^Q to Q_b for $i \in b$. There is

$i_0 \in b$ s.t. no truncation occurs on b above i_0 . Hence π_i^Q is total

on Q_i for $i_0 \leq i \in b$. For $i \in b$ set:

$$\kappa_i = \kappa_i^Q = \text{crit}(\pi_{i,b}^Q); \quad \tau_i = \tau_i^Q = (a_i^+)^{Q_i}$$

A standard proof shows that

there is a club $C \subset b \setminus i_0$ s.t. if

$$d \in C, \text{ then } d = \kappa_d \text{ and } \pi_{\alpha\beta}(\alpha) = \beta$$

for $\alpha, \beta \in C, \alpha \leq \beta$. It follows

$$\text{that } \pi_\alpha(\alpha) = \theta; \quad \alpha = \text{crit}(\pi_\alpha)$$

for $\alpha \in C$.

Now let $\kappa \in C$ be inaccessible. Then $\kappa \in Q_\kappa \cap K_\kappa$, since otherwise K_κ would be a segment of Q_κ and the coiteration would terminate,

$$(1) \kappa \neq \pi_{i, \kappa}^{\kappa}(\bar{\kappa}) \text{ for any } \bar{\kappa} < \kappa,$$

proof. Otherwise we could repeat the proof that coiterations terminate (§4 of [NFS]) to show that the coiteration of Q, K must have terminated below κ . QED(1)

Hence:

(2) There is no truncation on the branch $\{i \mid i \leq_{T, \kappa} \kappa\}$. (Hence $\pi_{i, \kappa}^{\kappa}$ is total for $i \leq_{T, \kappa} \kappa$.)

(Otherwise $\kappa = \pi_{i, \kappa}^{\kappa}(\bar{\kappa})$ where $\bar{\kappa}_i < \kappa$.)

It follows easily that:

$$(3) \pi_{i, \kappa}^{\kappa}(\bar{\zeta}) < \kappa \text{ for } \bar{\zeta} < \kappa;$$

hence:

$$(4) \pi_{i, \kappa}^{\kappa}(\kappa) = \kappa,$$

Let $\tau = \kappa + \aleph^c$. Then

$$(5) \pi_{i, \kappa}^{\kappa}(\tau) = \tau,$$

Now set: $\tilde{K}_\kappa = \int_{\Sigma} E^{K^\kappa}$, $K_\kappa^* = \pi_{0\kappa}^K(\tilde{K}_\kappa)$.

It follows easily that:

(1) $K_\kappa^* = \int_{\Sigma} E^{K_\kappa} = \int_{\Sigma} E^{Q_\kappa}$.

(Note that $\Sigma = \kappa + Q_\kappa$).

Set: $\tilde{\pi}_\kappa = \pi_{0\kappa}^K \upharpoonright \tilde{K}_\kappa$. Then

(2) $\tilde{\pi}_\kappa : \tilde{K}_\kappa \xrightarrow{\Sigma_0} K_\kappa^*$ cofinally.

But by (1):

(3) $\pi_{\kappa\kappa'}^Q : K_\kappa^* \xrightarrow{\Sigma_0} K_{\kappa'}^*$ cofinally

for $\kappa, \kappa' \in C'$, $\kappa \leq \kappa'$,

where C' = the set of inaccessible in C

We wish to define a stationary

$S \subset C'$ s.t. $\pi_{\kappa\kappa'}^Q \upharpoonright \tilde{\pi}_\kappa \in \tilde{\pi}_{\kappa'}$ for

$\kappa, \kappa' \in S$, $\kappa \leq \kappa'$. Recalling that

$\tau_\kappa = ht(\tilde{K}_\kappa)$ for $\kappa \in C'$, we set

$\tau_\kappa^* = ht(K_\kappa^*)$ and note that:

(4) $cf(\tau_\kappa) = cf(\tau_\kappa^*) = cf(\tau_{\kappa'})$

for $\kappa, \kappa' \in C'$, $\kappa \leq \kappa'$.

by (2), (3).

Let $cf(\tau_\kappa) = \delta$ for all $\kappa \in C'$,
 Pick $\langle \bar{\zeta}_i^\kappa \mid i < \delta \rangle$ monotone and
 cofinal in τ_κ for $\kappa \in C'$. Assume
 w.l.o.g. that $\kappa > \delta$ for $\kappa \in C'$.
 Then there is $\delta(\kappa) < \kappa$ s.t.

$\{ \overset{\sim}{\pi}_\kappa(\bar{\zeta}_i^\kappa) \mid i < \delta \} \subset \text{rng}(\pi_{\delta(\kappa)}^Q)$
 for $\delta = \delta(\kappa)$. Hence by Fodor
 there is a stationary $S' \subset C'$ s.t.
 $\delta(\kappa) = \delta$ is constant for $\kappa \in S'$.

Set: $h(\kappa) = \langle \pi_{\delta(\kappa)}^Q \overset{\sim}{\pi}_\kappa(\bar{\zeta}_i^\kappa) \mid i < \delta \rangle$.
 Then $h(\kappa) \in \mathcal{P}_\delta^\delta$. Hence there is
 stationary $S \subset S'$ s.t. h is constant
 on S . But then:

$$(5) \pi_{\kappa\kappa'}^Q \overset{\sim}{\pi}_\kappa(\bar{\zeta}_i^\kappa) = \overset{\sim}{\pi}_{\kappa'}(\bar{\zeta}_i^{\kappa'})$$

for $i < \delta$; $\kappa, \kappa' \in S$, $\kappa \leq \kappa'$.

$$(6) \overset{\sim}{\pi}_\kappa \upharpoonright \kappa = \overset{\sim}{\pi}_{\kappa'} \upharpoonright \kappa \text{ for } \kappa, \kappa' \in S, \kappa \leq \kappa'$$

proof

$$\overset{\sim}{\pi}_\kappa \upharpoonright \kappa = \pi_{0\kappa} \upharpoonright \kappa = \pi_{\kappa\kappa'} \pi_{0\kappa} \upharpoonright \kappa = \pi_{0\kappa'} \upharpoonright \kappa = \overset{\sim}{\pi}_{\kappa'} \upharpoonright \kappa$$

Hence:

$$(7) \text{ ring}(\pi_{\kappa\kappa'}, \tilde{\pi}_{\kappa}) \subset \text{ring}(\tilde{\pi}_{\kappa'}) \text{ for } \kappa, \kappa' \in S, \kappa \leq \kappa'$$

Proof.

Let $x \in \tilde{K}_{\kappa}$. Then $x \in \bigcup_{\Sigma_i} E_{\Sigma_i}^{\kappa}$ for some $i \in \Sigma$ where $E = E^{\kappa}$. Let $f = f_i^{\kappa}$ be the \tilde{K}_{κ} -least $f: \kappa \xrightarrow{\text{onto}} \bigcup_{\Sigma_i} E_{\Sigma_i}^{\kappa}$. Then $x = f(v)$ for a $v \in \kappa$. We have:

$$\begin{aligned} \tilde{\pi}_{\kappa}(f) &= \text{the } K_{\kappa}^{\times} \text{-least } f: \kappa \xrightarrow{\text{onto}} \bigcup_{\tilde{\pi}_{\kappa}(\Sigma_i)} E_{\Sigma_i}^{K_{\kappa}^{\times}} \\ \pi_{\kappa\kappa'}^{\mathcal{Q}} \tilde{\pi}_{\kappa}(f) &= \text{the } K_{\kappa'}^{\times} \text{-least } f: \kappa \xrightarrow{\text{onto}} \bigcup_{\tilde{\pi}_{\kappa'}(\Sigma_i)} E_{\Sigma_i}^{K_{\kappa'}^{\times}} \\ &= \tilde{\pi}_{\kappa'}(f_i^{\kappa'}). \end{aligned}$$

$$\begin{aligned} \text{Hence: } \pi_{\kappa\kappa'}^{\mathcal{Q}} \tilde{\pi}_{\kappa}(x) &= \pi_{\kappa\kappa'}^{\mathcal{Q}} \tilde{\pi}_{\kappa}(f(v)) = \\ &= \pi_{\kappa\kappa'}^{\mathcal{Q}} (\tilde{\pi}_{\kappa}(f)(\tilde{\pi}_{\kappa}(v))) = \tilde{\pi}_{\kappa'}(f_i^{\kappa'})(\tilde{\pi}_{\kappa'}(v)) = \\ &= \tilde{\pi}_{\kappa'}(f_i^{\kappa'}(v)) \in \text{ring}(\tilde{\pi}_{\kappa'}). \end{aligned}$$

(Note $\tilde{\pi}_{\kappa}(v) = \tilde{\pi}_{\kappa'}(v)$ by (6)), QED (7)

Now set: $\pi_{\kappa\kappa'} = \tilde{\pi}_{\kappa'}^{-1} \pi_{\kappa\kappa'}^{\mathcal{Q}} \tilde{\pi}_{\kappa}$. Then $\pi_{\kappa\kappa'}: \tilde{K}_{\kappa} \xrightarrow{\Sigma} \tilde{K}_{\kappa'}$ cofinally. Moreover $\pi_{\kappa\kappa'} \upharpoonright \kappa = \text{id}$ by (6) and $\pi_{\kappa\kappa'}(\kappa) = \kappa'$

QED (Lemma 5.1)

Now let $\tilde{K}, \langle \tilde{\pi}_\alpha \mid \alpha \in S \rangle = \lim_{\alpha \leq \beta \text{ in } S} (\tilde{K}_\alpha, \tilde{\pi}_{\alpha\beta})$

Then \tilde{K} extends K^c . Clearly $cf(\alpha + K^c) = \delta'$ for all $\alpha \in S$ where w.l.o.g. $\delta' < \min S$. Hence $cf(ht(\tilde{K})) = \delta' < \theta$. Hence

$\tilde{K} \in H_{\theta^+}$. Let $\theta \in H < H_{\theta^+}$ s.t.

$\bar{H} = \theta$, H is transitive, $\tilde{K}, \langle \tilde{\pi}_\alpha \mid \alpha \in S \rangle \in H$.

Let $f: \theta \leftrightarrow H$ and set $X_\alpha = f''\alpha$

for $\alpha \leq \theta$. Set $C = \{ \alpha \mid \alpha \leq \theta \text{ s.t. } \tilde{K}, \langle \tilde{\pi}_\alpha \mid \alpha \in S \rangle \in X_\alpha < H$

and $V_\alpha = V_\theta \cap X_\alpha$. Then C is club in θ

Pick $\kappa \in C \cap S$. Let $\sigma: N \xrightarrow{\sim} X_\kappa$.

Set $F^* = \sigma \upharpoonright \neq(\kappa)$. Then F^* is an extender on N and $\sigma: N \xrightarrow{F^*} H$.

(Hence $V_\theta \subset H = \text{Ult}(N, F^*)$) (Clearly

$N \models ZFC^-$ and $V_\kappa \in N$.)

It is easily seen that:

$$(8) \quad \sigma(\langle \tilde{\pi}_{\alpha\beta} \mid \alpha \leq \beta < \kappa \text{ in } S \rangle) = \langle \tilde{\pi}_{\alpha\beta} \mid \alpha \leq \beta \text{ in } S \rangle.$$

Hence since $\tilde{K}_\kappa, \langle \tilde{\pi}_{\alpha\kappa} \mid \alpha \in S \cap \kappa \rangle =$

$=: \lim_{\alpha \leq \beta < \kappa \text{ in } S} (\tilde{K}_\alpha, \tilde{\pi}_{\alpha\beta})$, we have:

(9) $\sigma(\pi_{\alpha\kappa}) = \pi_\alpha$ ($\alpha \in \kappa \cap S$). Hence:

(10) $\sigma \upharpoonright \tilde{K}_\kappa = \pi_\kappa$, since $\sigma \pi_{\alpha\kappa}(x) = \pi_\alpha(x) = \pi_\kappa \pi_{\alpha\kappa}(x)$

Set: $F' = F * \upharpoonright \tilde{K}_\kappa = \pi_\kappa \upharpoonright \tilde{K}_\kappa$. Then

$\langle \tilde{K}, F' \rangle$ satisfies all premouse conditions except the initial segment condition.

We shorten F' so as to satisfy this condition: Let $\lambda < \theta$ be least

s.t. $\sigma(f)(\alpha) < \lambda$ whenever $f: \kappa \rightarrow \kappa$, $f \in \tilde{K}_\kappa$ and $\alpha < \lambda$. Set: $F = F' \upharpoonright \lambda$.

Let $\bar{\pi}: \tilde{K}_\kappa \rightarrow_F \bar{K}$, $\langle \bar{K}, F \rangle$ is

easily seen to be a 1-small premouse

(11) \bar{K} is an initial segment of \tilde{K} (hence of K^c).

proof.

Let $k: \bar{K} \rightarrow \tilde{K}_\kappa$ be defined by

$k(\bar{\pi}(f)(\alpha)) = \pi_\kappa(f)(\alpha)$. Then, k

is a cofinal Σ_0 preserving map

and $k \upharpoonright \lambda = \text{id}$, $k(\lambda) = k\bar{\pi}(\alpha) = \pi_\kappa(\alpha) = \theta$.

Let $\lambda < \bar{\zeta} < \text{ht}(\bar{K})$ s.t. $\text{wp}_{\bar{K}}^{\omega} \upharpoonright \bar{\zeta} = \lambda$.

(Such $\bar{\zeta}$ are cofinal in $\bar{\zeta} \text{ht}(\bar{K})$).

Set: $\bar{Q} = \bigcup_{\bar{\zeta}} E^{\bar{K}}$, $Q = \pi_\kappa(\bar{Q}) = \bigcup_{\pi_\kappa(\bar{\zeta})} E^{\tilde{K}}$;

$k' = k \upharpoonright \bar{Q}$. Then \bar{Q}, Q are round

and $k': \bar{Q} \rightarrow \sum^* Q$ where $\omega_{\bar{Q}}^{\omega} = \text{crit}(k')$.

We apply §8 Lemma 4 of [NFS].

(a) is impossible since $k' \neq \text{id}$ & Q is round

(c) is impossible since $\omega_{\bar{Q}}^{\omega} \geq \text{crit}(k')$.

Hence (b) holds - i.e. \bar{Q} is a request of Q , hence of \tilde{K} . QED (11)

We note that λ is a limit cardinal in \tilde{K} by the definition of λ . Hence:

(12) $\bigcup_{\lambda} E_{K^c} = \bigcup_{\lambda} E_{\tilde{K}}$, where λ is a limit cardinal in K^c .

(13) $\langle N, F^* \rangle$ is a certificate for $\langle \bar{K}, F \rangle$

Moreover, if $f: \bar{\kappa} \rightarrow \text{dom}(F)$, $\bar{\kappa} < \kappa$,

then $f \in N$,

proof.

$\langle N, F^* \rangle$ is trivially a certificate.

Since $\tilde{K}_\kappa \in N$, there is $g \in N$ s.t.

$g: \kappa \leftrightarrow \tilde{K}_\kappa$. Set $\bar{f} = g^{-1} \circ f$.

Then $\bar{f} \in V_\kappa \subset N$, since κ is regular.

Hence $f = g \circ \bar{f} \in N$. QED (13)

By § 11 Lemmas 2.2, 2.4 it follows from (13) that for $\delta = \delta(\lambda) = \sup\{\zeta \mid \tilde{\mu}_\zeta < \lambda\}$ we have:

$$(14) \quad M_\delta = N_\delta = \langle \bigcup_\lambda E^{\bar{\kappa}}, \emptyset \rangle \text{ and } \tilde{\mu}_\delta = \lambda.$$

Now let $\bar{\kappa} = \bigcup_\alpha E^{\bar{\kappa}} = \bigcup_\alpha E^{\kappa^c}$. Since λ is the largest cardinal in $\bar{\kappa}$, α is cardinally absolute in κ^c . By § 11 Lemma 2.2 we conclude that for $\delta' = \delta(\alpha)$ we have:

$$(15) \quad M_{\delta'} = N_{\delta'} = \langle \bar{\kappa}, \emptyset \rangle.$$

Hence by (13) $N_{\delta'+1} = \langle \bar{\kappa}, E \rangle$. Hence $\omega p_{N_{\delta'+1}}^\omega < \lambda$. This is a contradiction,

$$\text{since } \lambda = \tilde{\mu}_\delta = \min\{\omega p_{N_\zeta}^\omega \mid \zeta \geq \delta\}.$$

Contr! QED (Lemma 5)

Corollary 6.1 Let $Q \in \mathcal{V}_\theta$ be a normally iterable \mathcal{V} -premouse in \mathcal{V}_θ . Then Q is iterable in \mathcal{V}_θ .

proof of Corollary 6.1

Cociterate $\mathcal{Q}, \mathcal{K}^c$ to $\mathcal{Q}', \mathcal{K}'$. Since \mathcal{K}^c is preordinal, the usual proofs tell us that one side of the iteration must be simple and that the simple side must be a segment of the nonsimple side, if this occurs. \mathcal{K}' cannot be a proper segment of \mathcal{Q}' ; since the \mathcal{K} -side would then have to be non-simple and \mathcal{K}' would be unsound. Thus \mathcal{Q}' is a segment of \mathcal{K}' . Hence \mathcal{Q}' is a mouse. Hence \mathcal{Q} is a mouse, since $\pi : \mathcal{Q} \xrightarrow{\Sigma^*} \mathcal{Q}'$, where π is the iteration map. QED (6.1)

Corollary 6.2 If \mathcal{Q} is a ^{1-small} MS mouse in \mathcal{V}_θ , then \mathcal{Q} is a mouse in \mathcal{V}_θ .

We now prove a refinement of Corollary 6.1.

Thm 7 Let \mathcal{Q} be a countable 1-small premouse which is countably normally iterable. Then \mathcal{Q} is iterable in V_θ .
proof. Suppose not.

Coiterate \mathcal{Q}, κ^c , getting $\langle y^{\mathcal{Q}}, y^{\kappa} \rangle$ of length λ . On the κ -side we employ a strategy definable in V_θ . Hence there must be a failure of well foundedness on the \mathcal{Q} -side. Let:

$$y^{\mathcal{Q}} = \langle \langle \mathcal{Q}_i \rangle, \langle \nu_i \rangle, \langle \gamma_i^{\mathcal{Q}} \rangle, \langle \pi_{i,i}^{\mathcal{Q}} \rangle, T^{\mathcal{Q}} \rangle$$

$$y^{\kappa} = \langle \langle \kappa_i \rangle, \langle \nu_i \rangle, \langle \gamma_i^{\kappa} \rangle, \langle \pi_{i,i}^{\kappa} \rangle, T^{\kappa} \rangle.$$

Since we have a strategy on the κ -side, we suppose y^{κ} has been continued to length $\lambda+1$ if $\lim(\lambda)$

On the \mathcal{Q} -side we employ the "economical" strategy: If $\lim(\lambda)$, choose $\{i | i \in T^{\mathcal{Q}}\}$ if possible to be "economical" - i.e. s.t. there is no $\nu \geq \sup_{i < \lambda} \nu_i$ with $E_\nu^{M_\lambda} \neq \emptyset$. This means, in particular, that if $\{i | i \in T^{\mathcal{Q}}\}$ is non economical, then it is the ^{cofinal well founded} unique branch. Now let:

(Since otherwise \mathcal{Q} coiterates out to a segment of a mouse and is therefore a mouse in κ .)

$\langle y^{\bar{Q}}, y^{\bar{K}} \rangle \in X \prec V_{\theta}$, $\bar{X} = \omega$. Let

$\sigma: \bar{V} \xrightarrow{\sim} X$, where \bar{V} is transitive.

Let $\sigma(\langle y^{\bar{Q}}, y^{\bar{K}} \rangle) = \langle y^{\bar{Q}}, y^{\bar{K}} \rangle$. Let:

$y^{\bar{Q}} = \langle \langle \bar{Q}_i \rangle, \langle \bar{v}_i \rangle, \dots, T^{\bar{Q}} \rangle$.

$y^{\bar{K}} = \langle \langle \bar{K}_i \rangle, \langle \bar{v}_i \rangle, \dots, T^{\bar{K}} \rangle$, let $\sigma(\bar{\lambda}) = \lambda$.

$\langle y^{\bar{Q}}, y^{\bar{K}} \rangle$ is then a countable coiteration.

If $\delta < \bar{\lambda}$ is a limit ordinal, then

we know that $\{i \mid i \in T^{\bar{Q}} \delta\}$ is the ^{cofinal well founded} unique \forall branch through $T^{\bar{Q}} \delta$ in

the sense of \bar{V} , but not necessarily in V . Set: $\tilde{\lambda} =$ the least

$\tilde{\lambda} \leq \bar{\lambda}$ s.t. for all $\delta < \tilde{\lambda}$, $\{i \mid i \in T^{\bar{Q}} \delta\}$

is the unique cofinal well founded branch in $T^{\bar{Q}} \delta$. Then

(1) $\text{Lim}(\tilde{\lambda})$

pf. Suppose not. Let $\tilde{\lambda} = h+1$. Then

$\tilde{\lambda} = \bar{\lambda}$ and \bar{V} thinks that the coiteration cannot be continued. Since $\bar{\lambda}$ is

countable and we have followed the unique iteration strategy for Q ,

and Q is countably normally iterable,

then it can be continued. Hence

\bar{V} thinks so by an easy absoluteness

Note that $\bar{Q}_0 = Q$

argument. QED (1)

$T^{\bar{Q}}/\bar{\lambda}$ must have a well founded cofinal branch b , since Q is countably normally iterable. Choose b s.t. the limit model Q_b has minimal height.

Case 1 $ht(Q_b) \in \bar{V}$.

Case 1.1 b is the unique ^{well fnd} cofinal branch. Then $\bar{\lambda} = \bar{\lambda}$ and $\gamma^{\bar{Q}}$ has no ^{founded} cofinal branch. Hence \bar{V} thinks that $\gamma^{\bar{Q}}$ has no such branch. We derive a contradiction. Let $\bar{\lambda} < \delta \in \bar{V}$ s.t.

$L_{\delta}[\gamma^{\bar{Q}}]$ is admissible and $ht(Q_b) < \delta$.

Let $\bar{\xi} = ht(Q_b)$. Let $\mathcal{L} = L_{\delta, \bar{\xi}} =$

$= L_{\delta, \bar{\xi}, \gamma^{\bar{Q}}}$ be the following infinitary language on $L_{\delta}[\gamma^{\bar{Q}}]$: Predicates: $\in, =$; Constants: \underline{x} ($x \in L_{\delta}[\gamma^{\bar{Q}}]$), b° .

Axioms: ZF^- ; $\forall \sigma (\sigma \in \underline{x} \leftrightarrow \forall_{z \in \sigma} z \in \underline{x})$

($x \in L_{\delta}[\gamma^{\bar{Q}}]$); b° is a cofinal branch in $\gamma^{\bar{Q}}$; b° has a transitive limit model Q_b° s.t. $ht(Q_b^{\circ}) = \bar{\xi}$.

Clearly \mathcal{L} is consistent. Every model \mathcal{M} of \mathcal{L} is isomorphic to one which is good in the sense that its well founded core is a transitive \in -structure. Hence we may work only with good models.

We then have $\underline{x}^{\mathcal{M}} = x$ for $x \in L_{\delta} [y^{\bar{Q}}]$.

But then $b^{\circ \mathcal{M}}$ is really a well founded cofinal branch in $y^{\bar{Q}}$; hence $b^{\circ \mathcal{M}} = b$.

By the completeness theorem for countable admissible it follows that:

$v \in b \iff \mathcal{L} \models v \in b^{\circ}$. Hence $b \in \bar{V}$.

Contr! QED (Case 1.1)

Case 1.2 $y^{\bar{Q}} | \bar{\lambda}$ has another well founded cofinal branch.

By 56 and the minimality of $ht(Q_b)$, Q_b

has the form J_{α}^{EN} , where $N = J_{\delta}^E$, $\delta = \sup_{i < \bar{\lambda}} \nu_i^{Q_i}$ and $N = \bigcup_{i < \bar{\lambda}} J_{\lambda_i}^{E Q_i}$. More-

over δ is Woodin in Q_b if $\alpha > \delta$.

Claim 1 $\bar{K}_{\bar{\lambda}}$ is not a proper segment of Q_b

Suppose not. Then there is no truncation on the main branch to $\bar{K}_{\bar{\lambda}}$, since

otherwise \bar{K}_λ would be unbound. Hence $ht(\bar{K}_\lambda) = \text{On} \cap \bar{V} > ht(Q_b)$. Contr!

Claim 2 Q_b is a segment of \bar{K}_λ .

Suppose not. There is ν s.t. $\delta < \nu \leq \alpha$ and $E_{\nu, \bar{K}_\lambda} \neq \emptyset$, where δ is sound in $\bigcup_a E_{\delta, \bar{K}_\lambda}$.

It follows easily that \bar{K}_λ is not 1-small. Contr! QED (Claim 2).

But then Q_b is a mouse^{in \bar{V}_α} . Q_b is a simple iterate of Q , since otherwise Q_b could not be a proper segment of \bar{K}_λ . Hence \bar{K}_λ is not a simple iterate of \bar{K} .

Contradiction! Hence $\bar{\pi}_{\alpha, b}^Q : Q \rightarrow \sum_{\ast} Q_b$ and Q is a mouse in \bar{V} . Contr!

QED (Case 1, 2)

Case 2 Case 1 fails. (Hence $\tilde{\lambda} = \bar{\lambda}$)

We repeat an argument of Woodin.

Let Γ = the set of $\delta \in \bar{V}$ which are admissible in $\mathcal{Y}^{\bar{Q}}$. For $\delta \in \Gamma$:

let $L_\delta = L_{\delta, \mathcal{Y}^{\bar{Q}}}$ be the following theory in the infinitary language of $L_\delta[\mathcal{Y}^{\bar{Q}}]$:

Predicates $\in, =$. Constants \underline{x} ($x \in L_{\gamma}[\bar{y}^{\bar{Q}}]$ and b°). Axioms ZF^{-} , $\bigwedge \alpha (\forall x \in \underline{x} \leftrightarrow \forall z \in \underline{x} \exists z')$ ($x \in L_{\gamma}[\bar{y}^{\bar{Q}}]$), b° is a cofinal well founded branch through $\bar{y}^{\bar{Q}}$, $\bar{\alpha} \in Q_{b^{\circ}}^{\circ}$ ($\bar{\alpha} < \bar{\alpha}'$).

Clearly L_{γ} is consistent. By Velle's Lemma L_{γ} has a good model \mathcal{M} whose well founded core has rank exactly $\bar{\alpha}$. (As above, good means that the well founded core is a transitive \in -model.) Let $b_{\gamma} = b^{\circ \mathcal{M}}$. Then b_{γ} is a cofinal branch through $\bar{y}^{\bar{Q}}$ and $\bar{Q}_{\gamma} = (Q_{b^{\circ}}^{\circ})^{\mathcal{M}}$ is a good limit model whose well founded core has rank exactly $\bar{\alpha}$. But then $b_{\gamma} \neq b_{\gamma'}$ for $\bar{\alpha}, \bar{\alpha}' \in \Gamma$, $\bar{\alpha} < \bar{\alpha}'$. Repeating some arguments from §6 we see that if $\delta = \sup_{i < \bar{\lambda}} \bar{v}_i$ and $N = \bigcup_{\gamma} J_{\gamma}^E = \bigcup_{i < \bar{\lambda}} J_{\bar{v}_i}^{E \bar{Q}_i}$, then δ is Woodin in J_{γ}^E . But this holds for arbitrarily large $\bar{\alpha} \in \Gamma$. Hence δ is Woodin in $L_{\bar{\theta}}^E$, where $\bar{\theta} = \text{On} \cap \bar{v}$.

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Let $\sigma(E) = E'$, $\sigma(\delta) = \delta'$. Then δ' is Woodin in $L_{\theta}^{E'}$, hence in $L^{E'}$, since E' is inaccessible. Hence by A1, V_{θ} is closed under $\#$, hence \bar{V} is closed under $\#$.

Now let $\text{ht}(Q_b) = \bar{\zeta}$ and let $\delta' > \bar{\zeta}$ be admissible in $\gamma^{\bar{Q}}$. Let $\mathcal{L} = \mathcal{L}_{\delta', \bar{\zeta}, \gamma^{\bar{Q}}}$ be as in Case 1.1.

Then \mathcal{L} is consistent. If τ is the first indiscernible for $L[\gamma^{\bar{Q}}]$ given by $(\gamma^{\bar{Q}})^{\#}$, then $L_{\tau}[\gamma^{\bar{Q}}] \prec L[\gamma^{\bar{Q}}]$.

Hence there are $\bar{\delta}, \bar{\zeta} < \tau$ s.t. $\bar{\mathcal{L}} = \mathcal{L}_{\bar{\delta}, \bar{\zeta}, \gamma^{\bar{Q}}}$ is consistent.

Let \mathcal{M} be a good model of $\bar{\mathcal{L}}$ and set $\bar{b} = b^{\mathcal{M}}$. Then \bar{b} is a cofinal well founded branch through $\gamma^{\bar{Q}}$ and $\text{ht}(Q_{\bar{b}}) = \bar{\zeta} < \bar{\theta}$.
 Contradiction!

QED (Theorem 7)

We recall that M was called a weak mouse iff whenever $\sigma: Q \rightarrow \sum^* M$ and Q is countable, then Q is countably iterable. Call M a very weak mouse iff every such Q is countably normally iterable.

Then:

Corollary 7.1 Every ^{1-small} very weak mouse is a weak mouse in V_θ .

Hence:

Corollary 7.2 Every ^{1-small} weak MS-mouse is a weak mouse.

We believe that these facts should be provable without $A1-A3$, but don't know how.

The condition A3 was imposed only to ensure enough inaccessible cardinals to verify the background conditions for the construction of K^c . As mentioned in §2 of this addendum, we can get by with a weaker background condition if we are willing to work with MS-mice. Thus it would seem that we can prove Thm 5 for the K^c of MS-mice without use of A3. However, we also made use of A3 in proving Lemma 5.1. It would be natural to prove Lemma 5.1 replacing (i) by (i') $\alpha = \overline{\overline{V}}_\alpha$ and $cf(\alpha) > \omega$. The rest of the proof would then be essentially as before. However I don't see how to prove even this version of Lemma 5.1 without adopting some new assumption in place of A3. The following will do

A4 A2 holds for $M \subset V_\theta$ s.t. M is a 1-small premouse.

(In particular, this will hold for $M = K^c$)

The details are left to the reader.
We then get:

Thm 8 Assume $A1, A2, A4$. Let K^c be the K^c of MS-mice. Then K^c is universal wrt. 1-small premice in V_θ .
We would expect this to yield a version of Corollary 6.1 for MS-mice. We must, however, be careful in formulating. If $Q \in V_\theta$ is a normally iterable 1-small premouse in V_θ , then Thm 8 shows that Q normally iterates up to a normal iterate Q' of a sound MS-mouse \tilde{Q} . Let $\pi: Q \rightarrow Q'$ be the iteration map. Q' itself may not be MS-iterable and hence the existence of π does not entitle us to assert that Q is MS-iterable. It can be shown, however, that $\text{core}(Q')$ is an MS-mouse. Suppose now that Q' is a solid core mouse. Then $\pi(p_Q) = p_{Q'}$ and π induces a map

$\pi : Q \xrightarrow{\Sigma^*} \text{core}(Q')$. Hence Q is MS-iterable. Thus we get;

Cor 8.1 Assume A_1, A_2, A_4 . Let $Q \in \mathcal{V}_\theta$ be a normally iterable 1-small solid core mouse in \mathcal{V}_θ . Then Q is an MS-mouse in \mathcal{V}_θ .

Similarly, modifying the proof of Thm 7

Thm 9 Assume A_1, A_2, A_4 . Let Q be a countable 1-small solid core mouse which is countably normally iterable. Then Q is an MS-mouse in \mathcal{V}_θ .

Hence;

Thm 10 Assume A_1, A_2, A_4 . Let Q be a very weak ^{1-small} solid core mouse. Then Q is a weak MS-mouse.