

Addendum to Robust Extenders

We present a simplified version of the methods needed for a proof of [RE] §1 Theorem 1, which says that Steel's main realizability lemma for arrays with "background certification" also holds for arrays in which the extenders satisfy the weaker condition of robustness. As in [RE] we prove the theorem in ZFC, whereas Steel worked in \underline{V}_Z for an inaccessible Ω . Our new proof is technically simpler. In particular, we dispense with Steel's functions $c(i,j)$.

§1 Basic definitions and results

We recall the definition of robustness:

Def Let $N = \langle J_{\nu}^E, F \rangle$ be an active premouse. F is robust wrt. N iff whenever $U \subset \lambda = \lambda_\nu$ and $W \subset \mathcal{P}(n) \cap W$ ($n = n_\nu$) are countable, then there is $g: U \rightarrow n$ s.t.

$$(a) \langle g(\vec{\alpha}) \rangle \in X \iff \langle \vec{\alpha} \rangle \in F(X)$$

whenever $\alpha_1, \dots, \alpha_n \in U$ and $X \in W$.
 (Here $\langle \rangle$ is Gödel's tuple function on ordinals.)

(b) Let $\sigma_1, \dots, \sigma_m \subset U$, $\tau = \text{lub } U$, $\bar{\tau} = \text{lub } g''U$.
 Let φ be a Σ_1 formula. Then
 $C_{\bar{\tau}, \kappa}^E \models \varphi[g''\sigma_1, \dots, g''\sigma_m] \iff C_{\tau, \kappa}^E \models \varphi[\sigma_1, \dots, \sigma_m]$.

(Here $C_{\tau, \alpha}^E$ is the Chang hierarchy up to α over $\langle L_\tau[E], E \cap L_\tau[E] \rangle$. The precise definition is given in [RE] §1.)

* [RE] = "Robust Extenders"

(We recall the condensation principle;
Let $C \leq_{\Sigma_1} C_{\gamma, \gamma}^E$ s.t. $\delta \in C$ and $[C]^\omega \subset C$.

Then C is isomorphic to a $C_{\delta, \bar{\gamma}}^E$ for an $\bar{\gamma} \leq \gamma$.
From this it follows that if $\alpha^\omega < \kappa$ for
all $\alpha < \kappa$, then $C_{\delta, \kappa}^E \leq_{\Sigma_1} C_{\delta, \infty}^E$ for all
 $\delta < \kappa$. This holds in particular for
 $\kappa = (2^\beta)^+$.)

As in [RE] § 2 we also define:

Def Let $N = \langle J_\nu^E, F \rangle$ be as above. Let
 $\kappa = \kappa_\nu, \lambda = \lambda_\nu$ be as above. Let $\kappa \leq \delta \leq \lambda$.
 F is robust up to δ in N iff for every
pair of countable sets $U \subset \lambda, W \subset \#(a) \cap N$
there is a $g: U \rightarrow \kappa$ s.t. (a) holds as
above and:

(b') Let $c = \text{lub}(U \cap \delta), \bar{c} = \text{lub } g''(U \cap \delta)$.

Then for all $\sigma_1, \dots, \sigma_m \in U \cap \delta$ and all Σ_1
formulae φ :

$$C_{\bar{c}, \kappa}^E \models \varphi[g''\sigma_1, \dots, g''\sigma_m] \iff C_{c, \infty}^E \models \varphi[\sigma_1, \dots, \sigma_m].$$

Def A premouse N is robust iff whenever
 $N \parallel \nu = \langle J_\nu^E, F \rangle$ is active and $\delta \in (\kappa_F, \lambda_F]$
is a cardinal in N , then F is robust
up to δ in $N \parallel \nu$.

As we remarked in [RE], if $\langle N_i : i \leq \theta \rangle$
($\theta \leq \infty$) is an array in which all extenders
added are robust, then each N_i is a
robust premouse.

We again prove only a special case of the theorem, which reduces to:

Main Claim Let N be a mouse-like, \aleph_1 -robust premouse satisfying ZFC^- . Let $\mathcal{Y} = \langle \langle P_i \rangle, \langle \nu_i \rangle, \langle \pi_i \rangle, T \rangle$ be a countable normal putative^{*} iteration of a countable P_0 without truncations, where $\sigma: P_0 \prec N$.

Then one of the following holds:

(a) $lh(\mathcal{Y}) = h+1$ and there is $\sigma': P_h \prec N$ s.t. $\sigma' \pi_{0h} = \sigma$.

(b) \mathcal{Y} has a maximal branch b , which is of limit length, and there is $\sigma': P_b \prec N$ s.t. $\sigma' \pi_b = \sigma$.

Before beginning the proof we reintroduce the notions of "world" and "enlargement". For technical reasons, some definitions will differ inessentially from the earlier ones. The normal iteration \mathcal{Y} is fixed from now on. Let ZFC^* be the theory

$$ZFC^- + \bigwedge \alpha [\alpha]^\omega \in V.$$

As before we define:

Def A world is a transitive structure W s.t. $W \models ZFC^*$ and $[\tau]^\omega \cap W = [\tau]^\omega \cap V$ for $\tau = 0_M \cap W$.

^{*} "putative" means that if $lh(\mathcal{Y}) = h+1$, then P_h does not need to be well founded.

(Note Since we shall often find our worlds in generic extensions of V , the clause $[\tau]^\omega \wedge W \subset V$ is not vacuous.)

Fix $A \subset On$ s.t. whenever $\beta > \omega$, $\beta = \overline{\overline{V}}_\beta$, then $V_\beta = L_\beta[A]$ and $L_{\beta^+}[A] \models (\beta \text{ is the largest cardinal})$

Clearly $L_{\beta^+}[A]$ is a world whenever $\beta = \overline{\overline{V}}_\beta$ and $cf(\beta) > \omega$.

Def By a standard world we mean a

$$W = \langle L_{\beta^+}[A], \epsilon, A \cap \beta^+, N, P_1, \dots, P_m \rangle \text{ s.t.}$$

$P_1, \dots, P_m \in W$ and $N \in L_\beta[A]$ is

a robust mouse-like premouse,

and $cf(\beta) > \omega$.

(Hence $(C_{\tau, \infty}^e)_W = C_{\tau, \beta^+}^e \prec C_{\tau, \infty}^e$ for $e, \tau \in W$.)

Clearly $\gamma \subset H_{\omega_1} \subset W$ for any world W .

For standard worlds we shall assume, in addition, that γ is W -definable

(e.g. $\gamma = P_1$).

[Note We generally write:

$$W = \langle L_{\beta^+}[A], N, \vec{P} \rangle \text{ or } W = \langle |W|, N, \vec{P} \rangle.]$$

If $W = \langle L_{\beta^+}[A], N, \vec{P} \rangle$ is a standard world, we define for each $\mu < \infty$

the μ -th enhancement:

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$W^{(\mu)} = \langle L_{\beta^+}^{\mu}[A], N, \vec{P}, \beta \rangle$, where
 $\beta_i =$ the i -th $\beta' \geq \beta$ s.t. $\beta = \overline{\overline{\beta}}$ and
 $cf(\beta) > \omega$.

(Hence $W^{(0)} = \langle W, \beta \rangle$.)

Note that $\langle \beta_i \mid i \leq \mu \rangle$ is uniformly
 $W^{(\mu)}$ -definable.

We then define a good world to be
 an arbitrary world which has the
 salient features of a standard world:

Def $\tilde{W} = \langle L_{\alpha}[\tilde{A}], \epsilon, \tilde{A}, \tilde{N}, P_1, \dots, P_m \rangle$ is a
good world iff \tilde{W} is a world in which

the following hold:

(a) If $\beta = \overline{\overline{\beta}}$, then $V_{\beta} = L_{\beta}[\tilde{A}]$. Moreover, if
 β^+ exists, then β is the largest cardinal
 in $L_{\beta^+}[\tilde{A}]$

(b) There is a largest cardinal β . Moreover,
 $\beta = \overline{\overline{\beta}}$ and $cf(\beta) > \omega$.

(c) $\tilde{N} \in V_{\beta}$ is a robust, mouse-like premouse.

If \tilde{W} is a good world, we denote its
 largest cardinal by $\beta_{\tilde{W}}$.

(Note If we were doing the full proof for
 arrays, then an array would take the
 place of N .)

Def Let $\tilde{w} = \langle L_\alpha, [\tilde{A}], \tilde{N}, \tilde{P} \rangle$ be a good world. By an enhancement of \tilde{w} we mean a good world

$$w' = \langle L_\alpha, [A'], \tilde{N}, \tilde{P}, \beta \rangle$$

s.t. $\beta = \beta_{\tilde{w}}$, $\tilde{A} = A' \cap \alpha$, $\alpha = \beta + w'$ and $V_\beta^{\tilde{w}} = V_\beta^{w'}$. If w' is an enhancement of \tilde{w} , then in w' we can define a sequence $\langle \beta_i \mid i \leq \mu \rangle$ by:

β_i is the least $\beta' \geq \beta$ s.t. $\beta' > \beta_n$ for $n < i$, $\beta' = \overline{V_{\beta'}^{w'}}$, and $cf(\beta') > \omega$.

Then $\beta_\mu^{w'} = \beta_{w'}$. By $\tilde{w}^{(i)}$ ($i \leq \mu$)

we denote $\langle L_{\beta_i^+}, [A'], \tilde{N}, \tilde{P}, \beta \rangle$,

where β_i^+ denotes $(\beta_i^+)^{w'}$.

We denote β_i by $\beta_i^{w'}$ and μ by $\mu^{w'}$.

Clearly each $\tilde{w}^{(i)}$ is a good world.

We now define the notion of enlargement. We shall demand less of enlargements than we did in [RE].

Def Let $\gamma \leq \text{lh}(\mathcal{Y})$. $\mathbb{E} = \langle \langle W_i, \sigma_i \rangle \mid i < \gamma \rangle$ is an enlargement of $\mathcal{Y} \mid \gamma$ iff

(a) $W_i = \langle L_{\alpha_i}[A_i], N_i, \vec{P} \rangle$ is a good world

(b) $\sigma_i \in \mathcal{V}$ s.t. $\sigma_i \restriction P_i \prec N_i$

(Hence $\sigma_i \in W_i$ since $\sigma_i \subset V \cap W_i$ is countable.)

(c) Set: $\delta_i =$ the largest $\delta \leq \lambda_i$ which is a cardinal in P_i (hence $\delta_i \leq \lambda_i \leq \delta_j$ for $i < j < \gamma$).

Set: $c_i = \sup \sigma_i \restriction \delta_i$, $E_i = E^{N_i}$. Then

$$\sigma_i \restriction \delta_h = \sigma_h \restriction \delta_h, \quad \bigcup_{c_h} E_i = \bigcup_{c_h} E_h \quad \text{for } h < i.$$

(d) Set: $\delta_i = \sigma_i \restriction \delta_i$. Then $\delta = \langle \delta_i \mid i < \gamma \rangle \in \mathcal{V}$.

(Hence $\langle \delta_h \mid h \leq i \rangle \in W_i$.)

(e) Set $t_i = \text{th}(\langle W_i, \sigma_i, \langle t_h \mid h < i \rangle \rangle)$

(where $\text{th}(\mathcal{M})$ is the complete theory of \mathcal{M}).

Then $t = \langle t_i \mid i < \gamma \rangle \in \mathcal{V}$. (Hence

$t \in W_i$ since $t \subset H_{\omega_1}$ is countable.)

Note The same definition gives the notion "enlargement of a phalanx", since we have not used the fact that \mathcal{Y} is an iteration but merely that $P_i \restriction \lambda_h = P_h \restriction \lambda_h$ and λ_h is a cardinal in P_i for $i < h$.

Def $\langle t, \delta \rangle$ is called the trace of \mathbb{E} .

We refer to t as the first trace and

δ as the second trace.

We write $\langle t, \delta \rangle = \text{tr}(\mathbb{E})$.

Def $\tau = \langle t, \delta \rangle$ is a trace ^{for $\gamma \mid \gamma$} iff there is a n.t. coll $(w, d) \Vdash \forall E \tau = \text{tr}(E)$.

Note that the trace of an enlargement E is always in V , even if E is not.

Def Let $N = \langle \bigcup_{i < \gamma} E_i, F \rangle$ be a premouse.
 $\tau = \langle t, \delta \rangle$ is an N -conforming trace for $\gamma \mid \gamma$ iff there is a n.t. coll $(w, d) \Vdash \forall E (\tau = \text{tr}(E) \wedge$

$$\wedge \bigcup_{c_i} E_i = \bigcup_{c_i} E \text{ for } i < \gamma).$$

Def Let N be as above and let $\sigma: P \prec N$, $\tau = \langle t, \delta \rangle$ is an N, σ -conforming trace iff τ is an N -conforming trace and $\delta_i = \sigma \upharpoonright \delta_i$ for $i < \gamma$.

(Note This definition makes sense for an arbitrary $\sigma: \sup_{i < \gamma} \delta_i \rightarrow On$.)

A very basic lemma is the following:

Lemma 1 There is a Σ_1 formula φ s.t. whenever $N = \langle J_\alpha^E, F \rangle$ is a premouse, then ε is an N -conforming trace for $\mathcal{Y} \upharpoonright \gamma$ iff the following hold:

(a) $\varepsilon = \langle t, \delta \rangle$ where $\delta = \langle \delta_i \mid i < \gamma \rangle$ and

$$\delta_i : \delta_i \rightarrow \text{On}_N \text{ for } i < \gamma$$

(b) $C_{c, \infty}^E \models \varphi[\varepsilon, \mathcal{Y} \upharpoonright \gamma]$, where

$$c = \sup_{i < \gamma} c_i, \quad c_i = \sup \delta_i \text{'' } \delta_i \dots$$

proof.

(a) is obviously a necessary condition. Now let (a) hold. For appropriate φ we show:

Claim ε is an N -conforming trace \iff (b).

The condition φ says:

There exist β, d s.t.

(a) $c < \beta < d$, $\varepsilon = \langle t, \delta \rangle \in C_{c, \beta}^E$, $C_{c, d}$ is admissible

(b) T is consistent, where $T = T_{\beta, d}(\varepsilon)$ is

the theory in $\mathcal{L}_{C_{c, d}}^E$ with:

Predicate: \dot{E}

Constants: \underline{x} ($x \in C_{c, d}$), $\dot{E}, \beta, A, N, \sigma, \dot{w}, \dot{p}$

Axioms

(1) ZFC⁻, $\wedge \sigma (\sigma \in \underline{x} \iff \forall_{z \in \underline{x}} \sigma = z)$ for $x \in C_{c, d}^E$

(2) $\dot{E} = \langle \langle \dot{w}_i, \dot{\sigma}_i \rangle \mid i < \gamma \rangle$ where

the following hold:

(ii) Each of $\dot{W}, \dot{\beta}, \dot{A}, \dot{\sigma}, \dot{P}, \dot{N}$ maps $\underline{\gamma}$ into

$$C_{c, \beta}^E$$

(iii) $\dot{\beta}_i < \underline{\beta}, \dot{A}_i < \underline{\beta}_i, \dot{\sigma}_i : \underline{P}_i < \dot{N}_i,$

$$\dot{W}_i = \langle L_{\dot{\beta}_i} [A_i], E, A_i, \dot{N}_i, \dot{P}_i \rangle \quad (i < \gamma)$$

(iii') $\dot{W}_i \cap [\dot{\beta}_i]^\omega = [\underline{\beta}]^\omega \cap \mathcal{F}(\dot{\beta}_i) \quad (i < \gamma)$

(Note $[\underline{\beta}]^\omega = \{x \mid C_{c, \beta+1}^E \models (x < \underline{\beta} \wedge \bar{x} \leq \omega)\}$;
hence $\langle [\underline{\beta}]^\omega \mid \beta < \alpha \rangle$ is $\Delta_1(C_{c, \alpha}^E)$.)

(iv) $\dot{W}_i \models (ZFC^* \wedge (a) \wedge (b) \wedge (c))$, where
(a), (b), (c) are as in the definition of
"good world". $(i < \gamma)$

(v) $\underline{t}_i = \text{th}(\langle \dot{W}_i, \dot{\sigma}_i, \underline{t} \upharpoonright \dot{N}_i \rangle) \quad (i < \gamma)$

(vi) $\underline{\sigma}_i \upharpoonright \underline{\delta}_h = \underline{\delta}_h \quad ; \quad \int_{\underline{c}_h}^{E \dot{N}_i} = \int_{\underline{c}_h}^{E N}$

$(h \leq i < \gamma)$.

proof of Claim:

(\rightarrow) Let $\mathbb{E} = \langle \mathbb{E}_i \mid i < \gamma \rangle$ be an N-conforming enlargement of $\mathcal{M} \upharpoonright \gamma$ with trace τ , where $\mathbb{E} = V[G]$ and G is $\text{coll}(\omega, \aleph_1)$ -generic for some \aleph_1 . Let $\mathbb{E} \in V_\beta[G]$ and let $\alpha > \beta$ s.t. $C_{c, \alpha}^E$ is admissible. Let $\alpha \leq \mu$ s.t. $V_\mu[G] \models ZFC^-$. Then $\langle V_\mu[G], \mathbb{E} \rangle$ models $T_{\beta, \alpha}$. Hence $T_{\beta, \alpha}$ is consistent.

(\leftarrow) Let \mathcal{M} model $T_{\beta, \alpha}$ (in $\mathcal{V}[G]$, where G is coll(ω, α)-generic). We can assume w.l.o.g. that \mathcal{M} is solid (i.e. $wf_{core}(\mathcal{M})$ is transitive and $E_{\mathcal{M}} \upharpoonright wf_{core}(\mathcal{M}) = E \upharpoonright wf_{core}(\mathcal{M})$).
 Then $\underline{x}^{\mathcal{M}} = x$ for $x \in C_{c, \alpha}$. It follows easily that $\mathbb{E} = \mathbb{E}^{\mathcal{M}}$ is an enlargement with traces. QED (Lemma 1)

An easy modification of the proof yields:

Cor 1.1 " π is a trace" is uniformly $\Sigma_1(C_{0, \infty})$ -definable.

The proof of Lemma 1 can be carried out in any good world. Hence:

Cor 1.2 Let $W = \langle L_\alpha[A], N, P \rangle$ be a good world. Let $\pi \in W$. Then

$W \models \pi$ is an N -conforming trace" holds iff (a) holds and

(b) $(C_{c, \infty}^E)_W = \varphi[\pi, \gamma, \gamma]$,

where $c = \sup_{i < \gamma} \delta_i$ " δ_i ", $\pi = \langle t, \delta \rangle$.

Hence:

Cor 1.3 " π is an N -conforming trace" is absolute in standard worlds W , since $(C_{c, \infty}^E)_W \prec C_{c, \infty}^E$ for $c \in W$.

As a corollary of the proof we also get:

Cor 1.4. Let $N^* = \langle J_\gamma^{E^*}, F^* \rangle$ be a premouse, let E be an N^* -conforming enlargement of $\mathcal{M}|\mathcal{N}$. Let $W = \langle |W|, N, \rho \rangle$ be a good world. Let $J_c^{E^N} = J_c^{E^*}$, where $c = \sup_{i < \gamma} c_i$, $c_i = \sup \delta_i \text{ " } \delta_i$. Suppose moreover that $\text{sn}(E) < \text{On} \cap W$. Let $\tau = \text{tr}(E)$. Then $W \models \tau$ is an N -conforming trace.

prf.

Let $\text{sn}(E) = \beta$. Let $\alpha > \beta$ be least s.t. $C_{c, \alpha}^{E^N}$ is admissible. Then $\alpha \in W$. But then the theory $T_{\alpha, \beta}(\tau)$ in the proof of Lemma 1 has a model. Hence $T_{\alpha, \beta}(\tau)$ is consistent.

QED(1.4)

We shall confine ourselves largely to traces having the following property:

Def Let E be an enlargement of $\mathcal{M}|\mathcal{N}$. E is meat (or self justifying) iff, letting $\tau = \text{tr}(E)$, we have

$W_i \models \tau \upharpoonright i$ is an N_i -conforming trace for $i < \gamma$. (Here $\tau \upharpoonright i$ denotes $\langle t \upharpoonright i, \delta \upharpoonright i \rangle$, where $\tau = \langle t, \delta \rangle$.)

(Note It follows that

$W_i \models \tau \upharpoonright i$ is an N_i, σ_i -conforming trace.)

Def A meat trace is the trace of a meat enlargement.

Note Let $r = \langle t, \delta \rangle = \text{tr}(\mathbb{E})$. Then \mathbb{E} is meat iff t satisfies a syntactical condition of the form: $x_i \in t_i$ for $i < \gamma$.

Hence:

Lemma 1.5 Let r be a trace for γ .
(a) r is meat iff it satisfies the above syntactical condition

(b) If r is meat and \mathbb{E} is any enlargement with trace r , then \mathbb{E} is meat.

We call a good world reflective if it countenances the existence of smaller imitations of itself in the following sense:

Def Let W be a good world. W is reflective iff whenever, $e, \tau \in W$ and $t = th(\langle W, e \rangle)$, then $t \in W$ and the following holds in W :

For sufficiently large δ it is forced by $coll(W, \delta)$ that there is a good world \bar{W} s.t. $\tau \in \bar{W}$ and for some $\bar{e} \in \bar{W}$: $t = th(\langle \bar{W}, \bar{e} \rangle)$ and $J_{\tau}^{\bar{e}} = J_{\tau}^e$.

The method of proof used in Lemma 1 yields:

Lemma 2 There is a Σ_1 formula ψ s.t. for all good worlds W :

W is reflective iff whenever $e, \tau \in W$ and $t = tr(\langle W, e \rangle)$, then $t \in W$ and: $(C_{\tau, \infty}^e)_W \models \psi[t]$.

proof.

ψ says that there exist β, α s.t. $\tau < \beta < \alpha$, $C_{\tau, \alpha}^e$ is admissible, and

$T^* = T_{e, \beta, \alpha}^*(t)$ is consistent, where

T^* is the following theory in the finitary language of $C_{\tau, \alpha}^e$:

Predicate : e

Constants \underline{x} ($x \in C_{\tau, \alpha}$); $\dot{W}, \dot{e}, \dot{A}, \dot{N}, \dot{p}$

Axioms

(1) ZFC^- , $\bigwedge v (v \in \underline{x} \leftrightarrow \bigvee_{z \in \underline{x}} v = \underline{z})$ for $x \in C_{\tau, \alpha}$.

(2) $\dot{W} = \langle L_{\underline{\beta}}[\dot{A}], \dot{e}, \dot{A}, \dot{N}, \dot{p} \rangle$, where the following hold:

(i) $\dot{A} \subset \underline{\beta}$ and $\dot{W}_i \models (ZFC^* + (a) \wedge (b) \wedge (c))$, where (a), (b), (c) as in the definition of "good world"

(ii) $\dot{e} \in \dot{W} \wedge \underline{t} = th(\langle \dot{W}, \dot{e} \rangle)$

(iii) $\underline{J}_{\underline{t}}^{\dot{e}} = \underline{J}_{\underline{t}}^e$

(iv) $\dot{W} \cap [\underline{\beta}]^{\omega} = \underline{[\beta]}^{\omega}$

(where $[\underline{\beta}]^{\omega} = \{x \mid C_{\tau, \beta+1}^e \models (x \subset \beta \wedge \bar{x} \leq \omega)\}$)

is $C_{\tau, \alpha}^e$ -definable.)

If W is reflective and $e, \dot{t} \in W$, then clearly there is β s.t. letting α = the smallest $\alpha > \beta$ s.t. $C_{\tau, \alpha}^e$ is admissible, then T^* has a model in $W[G]$ for a $coll(\omega, \aleph_1)$ -generic G .

Hence T^* is consistent. Hence

$C_{\tau, \infty}^e \models \psi[\dot{t}]$ in W . Conversely, if T^* is consistent and G is $coll(\omega, \aleph_1)$ -generic over W , then T^* has

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a solid model^{or} in $W[G]$. Set $\bar{w} = \dot{w}^{or}$, $\bar{e} = \dot{e}^{or}$. Then \bar{w}, \bar{e} have the desired properties. QED (Lemma 2)

Def A good world W is enhanceable iff it has a proper enhancement (i.e. there is a good world W' of the form $W^{(1)}$).

Note If $W = \bar{W}^{(1)}$, then $W^{(i)}$ is enhanceable for all $i < \mu$.

As a corollary of the above proof:

Cor 2.1 If W is enhanceable, then it is reflective.

prf.

Let $W' = W^{(1)}$. Let $e, \tau \in W$. Let $\beta = 0 \cap W$ and let $\alpha =$ the smallest $\alpha > \beta$ s.t. $C_{\tau, \alpha}^e$ is admissible. Then $\alpha \in W'$.

$T^* = T_{e, \beta, \alpha}^*$ has the model $\langle W', W, e \rangle$,

where W interprets \dot{w} and e interprets \dot{e} (where $t = \text{th}(\langle W, e \rangle)$). Thus

$(C_{\tau, \infty}^e)_{W'} = \varphi(t)$. But

$(C_{\tau, \infty}^e)_W \prec_{\varepsilon_1} (C_{\tau, \infty}^e)_{W'}$. Hence

$(C_{\tau, \infty}^e)_W = \varphi(t)$. QED (2.1)

Lemma 2.2 Let $\mathbb{E} = \langle \langle W_i, \sigma_i \rangle \mid i < \gamma \rangle$ be a neat enlargement of $\mathcal{Y} \upharpoonright \gamma$. Let W_i be reflective. Let $\varepsilon = \text{tr}(\mathbb{E})$. Then $W_i \varepsilon \upharpoonright (i+1)$ is an N_i, σ_i -conforming trace, p.f.

Let $\varepsilon = \langle t, \delta \rangle$. We must show that, if G is $\text{coll}(w, \delta)$ -generic over W_i for a sufficient δ , then $W_i[G]$ contains an N_i, σ_i -conforming enlargement of $\mathcal{Y} \upharpoonright (i+1)$ with trace $\varepsilon \upharpoonright (i+1)$. Since ε is neat, we can assume δ large enough that there is $\bar{\mathbb{E}} \in W_i[G]$ enlarging $\mathcal{Y} \upharpoonright i$ with trace $\varepsilon \upharpoonright i$. Writing:

$$(a, b, c) = a \times \{0\} \cup b \times \{1\} \cup c \times \{2\},$$

let $e = (\sigma_i, E^{N_i}, t \upharpoonright i)$. Clearly $t_i = \text{th}(\langle W_i, \sigma_i, t \upharpoonright i \rangle)$ is recursive in $t^* = \text{th}(\langle W_i, e \rangle)$. (In fact, t_i is uniformly recursive in t^* in the sense that $t_i = F(t^*)$ and for any good world W and any $a, b, c \in W$ s.t. $b = N^W$, we have $\text{th}(\langle W, a, c \rangle) = F(\text{th}(\langle W, (a, b, c) \rangle))$.)

If we have chosen δ large enough, it follows by reflectivity that there is a good world $\bar{w} \in W_i[G]$ and a $\bar{e} \in \bar{w}$ s.t. $J_c \bar{e} = J_c e$ and $t^* = \text{th}(\langle \bar{w}, \bar{e} \rangle)$. But then, letting $\bar{N} = N\bar{w}$, $\bar{E} = E\bar{N}$, we have: $\bar{e} = (\bar{\sigma}, \bar{E}, t\uparrow i)$, where $\bar{\sigma} : P_i \prec \bar{N}$, $J_c \bar{E} = J_c E_i$, and $\bar{\sigma} \upharpoonright \delta_i = \sigma_i \upharpoonright \delta_i$. Moreover $t_i = \text{th}(\langle \bar{w}, \bar{\sigma}, t\uparrow i \rangle) = F(t^*)$. But then $\bar{E} \langle \bar{w}, \bar{\sigma} \rangle \in W[G]$ is then an enlargement of $\gamma \uparrow i+1$ with trace $\Delta \uparrow i+1$. QED (Lemma 2, 2)

Before proceeding further we adopt the following linguistic conventions:
 We often make a statement of the form:
 "There is an enlargement with property A".
 This means that there is such an enlargement in $V[G]$ if G is $\text{coll}(\omega, \delta)$ -generic over V for a sufficient δ . If we say:
 "Let \mathbb{E} be an enlargement with property A,
 Then there is an enlargement \mathbb{F} with property B", it is understood that we are working in such a $V[G]$ and are asserting that an \mathbb{F} with property B exists in $V[G][G']$, where G' is

$\text{coll}(\omega, \delta)$ generic over $V[G]$ for a sufficiently large δ . (In fact, statements of this sort will generally turn out to be equivalent to a Σ_1 statement in $C_{\bar{c}, \kappa}^e$ for appropriate e, \bar{c} .)

We recall that the ground model V is fixed throughout the discussion and was used in defining the notions "world" and "enlargement".

.

A major tool in proving iterability will be the following "superlemma". We first define:

Def An enlargement \mathbb{E} of \mathcal{U}_{i+1} is powerful iff \mathbb{E} is neat and W_i is reflective (where $\mathbb{E} = \langle \langle W_h, \sigma_h \rangle \mid h \leq i \rangle$).

Def W' is a segment of W iff W is a good world satisfying q :
 $W = \bar{W}^{(i)}$, where \bar{W} is a good world,
and $W' = \bar{W}^{(i)}$ in W for some $i \leq \mu$.

Lemma 3 (Superlemma)

Let $\mathbb{E} = \langle \langle w_i, \sigma_i \rangle \mid i \leq \gamma \rangle$ be a powerful enlargement of $\mathcal{Y} \upharpoonright \gamma + 1$ with trace $\alpha = \langle t, \delta \rangle$. Let $\bar{\mathbb{E}} = T(\gamma + 1)$ in \mathcal{Y} and let \bar{w} be a reflective segment of $w_{\bar{\mathbb{E}}}$.

(A) $\mathbb{E} \upharpoonright \bar{\mathbb{E}}$ extends to a powerful enlargement $\bar{\mathbb{E}} = \langle \langle \bar{w}_i, \bar{\sigma}_i \rangle \mid i \leq \gamma + 1 \rangle$ of $\mathcal{Y} \upharpoonright \gamma + 2$ with trace $\bar{\alpha} = \langle \bar{t}, \bar{\delta} \rangle$ s.t.

(i) $\bar{\mathbb{E}} \upharpoonright \bar{\mathbb{E}} = \mathbb{E} \upharpoonright \bar{\mathbb{E}}$, (ii) $t = \bar{t} \upharpoonright \gamma + 1$,

(iii) $\bar{w}_{\gamma+1} = \bar{w}$, (iv) $\bar{\sigma}_{\gamma+1} \upharpoonright_{\bar{\mathbb{E}} \upharpoonright \gamma+1} = \sigma_{\bar{\mathbb{E}}}$

(B) If \bar{w} is a proper segment of $w_{\bar{\mathbb{E}}}$, then \mathbb{E} extends to a powerful enlargement $\mathbb{E}' = \langle \langle w_i, \sigma_i \rangle \mid i \leq \gamma + 1 \rangle$ of $\mathcal{Y} \upharpoonright \gamma + 2$ with trace $\alpha' = \langle t', \delta' \rangle$ s.t.

(i) $\mathbb{E} = \mathbb{E}' \upharpoonright \gamma + 1$, (ii) $O_{w_{\gamma+1}} < O_{w_{\gamma}}$

(iii) $t'_{\gamma+1} = \bar{t}_{\gamma+1}$, where $\bar{\mathbb{E}}, \bar{\alpha} = \langle \bar{t}, \bar{\delta} \rangle$ are as in (A).

Prf. of Superlemma.

We first prove (A)

Since N_γ is robust in W_γ , there is a

$$g: \lambda_\gamma \rightarrow \sigma_\gamma(\kappa_\gamma) \text{ i.t.}$$

$$(*) \langle \vec{\alpha} \rangle \in E_{\lambda_\gamma}^{\mathbb{P}_\gamma}(x) \iff \langle g(\vec{\alpha}) \rangle \in \sigma_\gamma(x)$$

for $\alpha_1, \dots, \alpha_m < \lambda_\gamma$, $x \in \mathcal{P}(\kappa_\gamma) \cap \mathbb{P}_\gamma$

$$(**) \text{ Let } \delta_1, \dots, \delta_m \subset \delta_\gamma, c_\gamma = \sup \sigma_\gamma \text{ " } \delta_\gamma, \\ c = \sup g \text{ " } \delta_\gamma. \text{ (Hence } c < \sigma_\gamma(\kappa_\gamma) \text{).}$$

Then in W_γ we have:

$$C_{c, \infty}^{E_\gamma} \models \varphi[\sigma_\gamma \text{ " } \vec{\sigma}] \iff C_{c, \delta_\gamma(\kappa_\gamma)}^{E_\gamma} \models \varphi[g \text{ " } \vec{\sigma}]$$

for all Σ_1 formulae φ .

Then:

$$(1) g \upharpoonright \kappa_\gamma = \sigma_\gamma \upharpoonright \kappa_\gamma$$

(Apply (*) with $X = \{\aleph_3\}$ for $\aleph_3 < \kappa_\gamma$)

By (*) there is $\sigma: \mathbb{P}_{\gamma+1} < N_\gamma$ defined by:

$$(2) \sigma(\pi_{\aleph_3, \gamma+1}(f)(\alpha)) = \sigma_\gamma(f)(g(\alpha)).$$

$$\text{Hence } \sigma \upharpoonright \lambda_\gamma = g, \sigma \pi_{\aleph_3, \gamma+1} = \sigma_\gamma.$$

Let $\tau = \langle t, \delta \rangle = \text{tr}(E)$. Since W_γ is reflective we know by Lemma 2.2:

(3) $W_\gamma \models \tau$ is an N_γ, σ_γ -confirming trace.

Let $f: \omega \xrightarrow{\text{onto}} \gamma+1$. Set:

$t^* = \{ \langle i, k \rangle \mid i < \omega \wedge k \in t_{f(i)} \}$, where $\langle \cdot, \cdot \rangle$ is Gödel's pairing function on ordinals. Then $t^* \subset \omega$. (We assume $t_h \subset \omega$ for $h \leq \gamma$)

We also set:

$$u = \{ \langle i, \mu \rangle \mid i < \omega \wedge \mu < \delta_{f(i)}^1 \}$$

Then $u \subset \delta_\gamma$. Moreover $\sigma_\gamma(\langle i, \mu \rangle) = \langle i, \sigma_\gamma(\mu) \rangle$ for $\mu < \delta_{f(i)}^1, i \leq \gamma$.

Since $\delta_i = \sigma_\gamma \upharpoonright \delta_i$ for $i \leq \gamma$, we see that $\pi = \langle t, \delta \rangle$ is straightforwardly coded by $t^*, \sigma_\gamma \upharpoonright u$, and e , where i any $e \subset \omega$ coding $\gamma \upharpoonright \gamma+1$ and f .

Clearly $\sigma_\gamma \upharpoonright \omega = \text{id}$. By Lemma 2 it follows that (3) can be expressed in W_γ by:

$$(4) \ C_{c_\gamma, \infty}^{E_\gamma} = \Psi [t^*, \sigma_\gamma \upharpoonright u, e],$$

where $\Psi \in \Sigma_1$. Hence by (**):

$$(5) \ C_{c, \sigma_\gamma(u_\gamma)}^{E_\gamma} = \Psi [t^*, g \upharpoonright u, e] \dots$$

But $t^*, g \upharpoonright u, e$ code $\bar{\pi} = \langle t, \bar{\delta} \rangle$,

where $\bar{\delta}_i = g \upharpoonright \delta_i = \sigma \upharpoonright \delta_i$ for $i \leq \gamma$.

But $J_c^{E_\gamma} = J_c^{E_3}$, since

$c < \sigma_\gamma(u_\gamma) = \sigma_\gamma(u_\gamma) < c_3$. Thus

(6) $C_{c, \sigma_3}^{E_3}(u_3) \models \Psi [t^*, q^* u, e]$.

But this tells us $(\text{min } \sigma_3(u_3), \sigma \in \bar{W})$:

(7) $\bar{\pi}$ is an N_3, σ -conforming trace for $\gamma+1$ in \bar{W} .

Note that $\bar{\pi} \upharpoonright \bar{z} = \pi \upharpoonright \bar{z}$, since $\sigma \upharpoonright \delta_h = \sigma_{\bar{z}} \upharpoonright \delta_h$ for $h < \bar{z}$ by (1). By (7),

if G is $\text{coll}(w, \delta)$ -generic over \bar{W} for a sufficient δ , there is

$\bar{E}'' \in \bar{W}[G]$ which is an enlargement of $\gamma \upharpoonright \gamma+1$ with trace $\bar{\pi}$.

Set $\bar{E}' = \bar{E} \upharpoonright \bar{z} \cup \bar{E}'' \upharpoonright [\bar{z}, \gamma+1]$. Then \bar{E}' is also an enlargement of $\gamma \upharpoonright \gamma+1$ with trace $\bar{\pi}$. Finally set:

$\bar{E} = \bar{E}' \widehat{\langle \bar{W}, \sigma \rangle}$. \bar{E} is easily seen to have the desired properties.

We now prove (B1).

Let $t' = \bar{E}_{\gamma+1} = \text{th}(\langle \bar{W}, \sigma, t \rangle)$. For

$c < \beta < d$ with $C_{c,d}^{E_3}$ admissible

consider the theory $T_{\beta d}(\bar{\pi})$ in the language of $C_{c,d}^{E_3}$ with:

Predicate G

Constants \underline{x} ($x \in C_{c,d}^{E_3}$), $\bar{W}, \sigma, \bar{A}, \bar{N}, \bar{P}$

Axioms

(A) ZFC^- , $\wedge x (x \in \underline{x} \leftrightarrow \bigvee_{z \in x} z = \underline{z})$ for $x \in C_{c,d}$

(B) $\dot{W} = \langle L_{\underline{\beta}}[A], \in, A, N, p \rangle$ with:

(i) $A \subset \underline{\beta}$, $W_i \models ZFC^* + (a) \wedge (b) \wedge (c)$, where (a), (b), (c) as in the def. of "good world"

(ii) $\underline{a} \in \dot{W} \wedge \underline{t}' = Th(\langle \dot{W}, \sigma, \underline{t} \rangle) \wedge \sigma \in \dot{W}_1$
 $\wedge \bigwedge_{i \leq \gamma} \underline{\sigma}_i = \sigma \upharpoonright \underline{\delta}_i \wedge \sigma \upharpoonright P_{\gamma+1} \prec N$

(iii) $\dot{W} \cap [\underline{\beta}]^\omega = [\underline{B}]^\omega$

(Recall that $\{[\underline{\beta}]^\omega\}$ is $C_{c,d}^E$ -definable in β)

(iv) \dot{W} is reflective.

Then $T = T_{\beta, d}(\bar{a})$ is consistent, since

$\langle W_{\gamma+1}, \bar{W}, \sigma, A^{\bar{W}}, N^{\bar{W}}, P^{\bar{W}} \rangle$ is a model.

The statement: "There exist β, d s.t. $c < \beta < d$, $C_{c,d}^E$ is admissible, and $T_{\beta, d}(\bar{a})$ is consistent" has the form:

(7) $C_{c, \infty}^E \models \chi[\bar{a}, t']$

Hence it has the form:

(8) $C_{c, \infty}^E \models \chi[t^*, y^* u, e, t']$, where

χ is Σ_1 .

Since (7) holds in W_3 and

$(C_{c, \infty}^E)_{\bar{W}} \prec_{\Sigma_1} (C_{c, \infty}^E)_{W_3}$, we

conclude that (7) holds in \bar{W} . But it follows just as in [RE] by the mouse-likeness and robustness of N_γ in \bar{W} that:

$$(9) C_{c, \kappa_\gamma}^E \prec_{\Sigma_1} (C_{c, \infty}^E \upharpoonright \bar{W})$$

Hence:

$$(10) C_{c, \kappa_\gamma}^E \models \chi[t^*, q \text{ "} u, e, t' \text{"}]$$

By (**1) we conclude:

$$(11) (C_{c, \infty}^{E\gamma})_{W_\gamma} \models \chi[t^*, \sigma_\gamma \text{ "} u, e, t' \text{"}] \text{ or}$$

in other words:

$$(12) (C_{c, \infty}^{E\gamma})_{W_\gamma} \models \chi'[s, t']$$

(where $s = \text{tr}(E)$).

But this means that there are $\beta, \alpha \in W_\gamma$ s.t. $c_\gamma < \beta < \alpha$, and $C_{c_\gamma, \alpha}^{E\gamma}$ is admissible, and $T_{\beta, \alpha}(s)$ consistent. Now let G be $\text{coll}(w, \alpha)$ -generic over W_γ . Then $T_{\beta, \alpha}(s)$ has

a solid model \mathcal{M} in $W_\gamma[G]$.

Set: $W' = W \upharpoonright \mathcal{M}$, $\sigma' = \sigma \upharpoonright \mathcal{M}$. Then:

(13) W' is a reflective good world;

(14) $s \in W'$, $\sigma' \in W'$ and

$\sigma' : P_{\gamma+1} < N' = N^{W'}$ with $\delta_i = \sigma' \upharpoonright \delta_i$ ($i \leq \gamma$).

Moreover

(15) $t' = \text{th}(\langle W', \sigma', \lambda \rangle)$, hence:

(16) $W' \vDash \lambda$ is an N', σ' -conforming trace.

If we set $\mathbb{E}' = \mathbb{E} \sim \langle W', \sigma' \rangle$, then \mathbb{E}' has the desired properties.

QED (Superlemma)

§ 2 Iterability

We now give the proof of the Main Claim in § 1. N is a mouse-like, robust premouse satisfying ZFC^- . We can place it inside a standard world $W = \langle L_\alpha[A], N, p \rangle$, where \mathcal{Y} is W -definable. Hence there is a W -definable bijection $m^*: lh(\mathcal{Y}) \leftrightarrow \omega$. Following Steel we define:

Def $m(i) = \min \{ m^*(j) \mid i \leq_T j \text{ in } \mathcal{Y} \}$ ($i < lh(\mathcal{Y})$)

Def i survives at j ($i \text{ surv } j$) iff $i \leq j \wedge m(i) = m(j) \wedge m(l) \geq m(i)$ for all $l \in (i, j)$.

Steel establishes:

Lemma 1

(a) $i \text{ surv } j \rightarrow i \leq_T i$

(b) $(h \text{ surv } i \text{ surv } j) \rightarrow h \text{ surv } j$

(c) $(h \text{ surv } j \wedge h \leq_T i \leq_T j) \rightarrow h \text{ surv } i \text{ surv } j$

(d) Let b be a branch of limit length in \mathcal{Y} . b is maximal in \mathcal{Y} iff

$$\bigwedge i \in b \bigvee j \in b (i < j \wedge \neg i \text{ surv } j)$$

(Hence if $b = \{ h \mid h \leq_T \lambda \}$, $\lambda \text{ limit}$, then

$$\bigvee i \in b \bigwedge j \in b (i < j \rightarrow i \text{ surv } j))$$

We now define:

Def $i <_* j \iff (i \leq_T j \wedge \neg i \text{ surv } j)$

Def i dies at $j \iff (j < lh(\mathcal{Y}) \wedge$

$i < j \wedge$ whenever $h \geq j$, $T(h) = i$, then

$$T(h) <_* h.$$

Finally:

Def i is a break point at $\delta \leq lh(\gamma)$ iff
 iff $i < \delta$ and whenever $i < h < \delta$ and
 $T(h) \leq i$, then $T(h) <_* h$. (Another words:
 An $\gamma \upharpoonright \delta$ every $l \leq i$ dies at $i+1$.)

We now turn to the proof of the main claim. We are given $\sigma : P \prec N$, where P is a putative iteration of P , and wish to show that one of the following holds:

(a) $lh(\gamma) = h+1$ and there is $\sigma' : P_h \prec N$
 s.t. $\sigma \pi_{\sigma, h} = \sigma'$.

(b) There is a maximal branch b of limit length in γ and a $\sigma' : P_b \prec N$ s.t. $\sigma \pi_b = \sigma'$.

We shall assume (b) to be false and prove (a). We begin by reformulating

(b):

Def $\Gamma = \Gamma_N$ is defined by:

$$\Gamma = \Gamma_N = \{ \sigma \mid \forall i \sigma : P_i \prec N \}$$

$$\Gamma = \{ \langle \sigma', \sigma \rangle \mid \forall i, j (i <_* j \wedge \sigma : P_i \prec N \wedge \sigma' : P_j \prec N \wedge \sigma \pi_{\sigma, i} = \sigma') \}$$

If $\sigma \in \Gamma$ we also set:

$$\Gamma^\sigma = \Gamma_N^\sigma = \{ \langle \sigma', \sigma'' \rangle \mid \sigma'' \in \Gamma \wedge (\sigma'' \in \Gamma \vee \sigma' = \sigma) \}$$

Then (b) is equivalent to:

(b') R^σ is well founded

for our given $\sigma: P_0 \prec N$. Thus we are assuming R^σ to be well founded.

In the following we shall always deal with good worlds $w' = \langle W', N', p' \rangle$ in which γ, m^* are w' -definable by the same definition as in our standard world w . Hence there is a relation $R_{w'} = R_N$ definable in w' as R was defined in w . With this convention we define:

Def An enlargement $\mathbb{E} = \langle \langle W_h, \sigma_h \rangle \mid h \leq i \rangle$ of $\gamma \mid i+1$ is proud iff

(a) \mathbb{E} is neat

(b) $R_h = R_{W_h}^{\sigma_h}$ is well founded for $h \leq i$

(c) W_h has the form $\overline{W_h}(\mu_h)$, where

$\mu_h \geq$ the rank of σ_h in R_h ($h \leq i$)

(d) W_i is reflective

(e) If $h < i$ does not die at $i+1$, then W_h is reflective.

Def \mathbb{E} is semi-proud iff (a)-(c) hold.

Def s is a (semi) proud trace iff

s is the trace of a (semi) proud enlargement.

Note Semiproudness is equivalent to a syntactic condition of the form:

$x \in t_h$ for $h \leq i$. Hence any enlargement of a semiproud trace is semiproud.

We prove:

Main Lemma Let $j+1 \leq lh(\gamma)$. Let $i < j$ and let \mathbb{E} be a proud enlargement of $\gamma|_{i+1}$.

(a) If i is a breakpoint at $j+1$, then \mathbb{E} extends to a proud enlargement \mathbb{F} of $\gamma|_{j+1}$ s.t. $On_{W_l} < On_{W_i}$ for $i < l \leq j$.

(b) If i survives at j , then $\mathbb{E}|_i$ extends to a proud enlargement \mathbb{F} of $\gamma|_{j+1}$ s.t. $On_{W_l^{\mathbb{F}}} < On_{W_i^{\mathbb{E}}}$

for $i \leq l < j$ and $W_j^{\mathbb{F}} = W_i^{\mathbb{E}}$, $\sigma_j^{\mathbb{F}} \pi_i = \sigma_i^{\mathbb{E}}$.

Before proving this, we show that it implies the main claim. Note that if $\mu =$ the rank of σ in \mathbb{R} ,

then $\langle W^{(\mu)}, \sigma \rangle$ is a proud enlargement of $\gamma|_1$. Suppose, first, that $lh(\gamma) = j+1$. W.l.o.g. we may suppose: $m^*(j) = 0$.

Then o survives at j . Hence $\mathbb{E} \cap O = \emptyset$ extends to a proad enlargement \mathbb{F} of \mathbb{Y} with $W_j^{\mathbb{F}} = W$ and $\sigma_j^{\mathbb{F}} \pi_{o_j} = \sigma$.

Now suppose that \mathbb{Y} has limit length. We derive a contradiction. We suppose w.l.o.g. that $m^*(o) = 0$. Hence o is a breakpoint at $lh(\mathbb{Y})$.

Define ξ_i ($i < \omega$) by: $\xi_0 = o$;

$$\xi_{i+1} = \text{that } \xi \text{ s.t. } m^*(\xi) = \min \{ m^*(h) \mid h > \xi_i \}.$$

An easy induction shows that each ξ_i is a break point at $lh(\mathbb{Y})$, hence at ξ_{i+1} . By successive applications of (a) we get enlargements \mathbb{E}_i of $\mathbb{Y} \upharpoonright \xi_{i+1}$ s.t. $W_{\xi_{i+1}}^{\mathbb{E}_i} \in W_{\xi_i}^{\mathbb{E}_i}$ for $i < \omega$. Contr!

We now prove the main lemma by induction on j . Let it hold below j . We first prove (a), assuming (b) to hold:

Case 1. There is an $h < j$ which survives at j . Then $h > i$, since i is a break point at $j+1$. By the induction hypothesis \mathbb{E} extends to a proud enlargement \mathbb{E}' of $\mathcal{J}|_{h+1}$ with $\text{On}_{W_l} < \text{On}_{W_i}$ for $i < l \leq h$. We then apply (b) to \mathbb{E}' , getting a proud \mathbb{F} extending \mathbb{E}' s.t. $\text{On}_{W_l}^{\mathbb{F}} < \text{On}_{W_h}^{\mathbb{E}'}$ $< \text{On}_{W_i}^{\mathbb{E}}$ for $h \leq l < j$ and $W_j^{\mathbb{F}} = W_h^{\mathbb{E}'}$ (hence $\text{On}_{W_j}^{\mathbb{F}} < \text{On}_{W_i}^{\mathbb{E}}$). Then \mathbb{F} extends \mathbb{E} and has the right properties.

Case 2 Case 1 fails. Then j is a successor ordinal $h+1$ and $T(j) <^* j$. By the ind. hypothesis \mathbb{E} extends to a proud \mathbb{E}' enlarging $\mathcal{J}|_j$ s.t. $\text{On}_{W_l} < \text{On}_{W_i}$ for $i < l < j$. We then apply the Superlemma. By Superlemma (A) there is $\bar{\sigma} \in W_{\bar{3}}$. ($\bar{3} = T(j)$) s.t. $\bar{\sigma}: P_{j+1} < N_{\bar{3}}$ and

$\bar{\sigma} \pi_{3, i+1} = \sigma_3$. Hence $\bar{\sigma} \in R_3 \sigma_3$ and hence $\bar{\mu} < \mu_3$, where $\bar{\mu}$ = the rank of $\bar{\sigma}$

in R_3 . We then apply Superlemma (B) to $\bar{W}_3^{(\bar{\mu})}$, getting an enlargement E'' of $\gamma|_{i+1}$ extending E' s.t.

$On_{W_j} < On_{W_h}$ and W_j is reflective.

Moreover E' is neat, $R_{j,i}$ is well founded

and $W_j = \bar{W}_j^{(\mu_j)}$ for μ_j = the rank of σ_j in R_j , since

$$\langle \bar{W}_3^{(\bar{\mu})}, \bar{\sigma} \rangle \equiv \langle W_j, \sigma_j \rangle.$$

Hence E'' is proud. $\square \text{EP}((b) \rightarrow (a))$

We now prove (b). We first note:

Fact If i survives at j , then every $h \in [i, j)$ dies at $j+1$.

prf. Suppose not.

Let $\gamma \geq j+1$, $\exists = T(\gamma)$ s.t. \exists survives at γ and $\exists \in [i, j)$. Then $m(\gamma) = m(\exists) \geq m(i) = m(j)$

since $\exists \in [i, j)$. But $m(\gamma) \neq m(j)$, since $\gamma \not\leq_T j$ and $j \not\leq_T \gamma$. Hence $m(\exists) > m(i)$.

Hence \exists does not survive at γ ,

since $m(\exists) > m(j)$, $j \in (\exists, \gamma)$. Cont!

Case 1 $\text{Lim}(j)$

Let $\langle i_m \mid m < \omega \rangle$ be monotone s.t., $i_0 = i$, $\sup_m i_m = j$. (Hence $i_m \leq i_{m+1}$ for $m < m+1$.) We first apply the induction hypothesis to get successive enlargements E_m ($m < \omega$) s.t., $E_0 = E$ and E_{m+1} extends $E_m \upharpoonright i_m$ with E_{m+1} in prod; $W_{m+1}^{E_{m+1}} = W_m^{E_m} = W_i$,

$$\sigma_{m+1} \pi_{i_m, i_{m+1}} = \sigma_m \quad \text{for } m < \omega.$$

Let $\tilde{\alpha}_m = \langle t_m, \sigma_m \rangle = \text{tr}(E_m)$. Then $\tilde{\alpha}_m$ is a semi prod trace in $W_i = W_m^{E_m}$, since E_m is semi prod (hence neat) and W_i is reflective.

Hence $\langle \langle \tilde{\alpha}_m, \sigma_m \rangle \mid m < \omega \rangle$ forms a descending chain in the following relation S , which is defined in W_i :

Def $D =$ the set of $\langle \alpha, \sigma \rangle$ s.t., $\sigma : P_{i_m} \prec N_V$, α is an N_i , σ - conforming, semi prod trace for $\gamma \upharpoonright i_{m+1}$, and $\alpha \upharpoonright i = \tilde{\alpha}_0 \upharpoonright i$. ($\tilde{\alpha}_0 = \text{tr}(E)$)

$S =$ the set of $\langle \langle \alpha', \sigma' \rangle, \langle \alpha, \sigma \rangle \rangle \in D^2$ s.t. for some $m < \omega$:
 $\sigma : P_{i_m} \prec W_i$, $\sigma' : P_{i_{m+1}} \prec N_i$, $\sigma' \pi_{i_m, i_{m+1}} = \sigma$,
 $\alpha' \upharpoonright i_m = \alpha \upharpoonright i_m$.

Thus S is ill founded. Let $\langle \langle s_m, \sigma_m \rangle \mid m < \omega \rangle \in W_i$ be a chain through S . Define

$\sigma : P_i < N_i$ by $\sigma \upharpoonright_{i_m} = \sigma_m$. Set:

$$s = \bigcup_{m < \omega} s_m \upharpoonright_{i_m}$$

Claim s is an N_i, σ conforming trace in W_i proof.

s is obviously σ -conforming, since $\sigma \upharpoonright_{i_h} = \sigma_m \upharpoonright_{i_h} = s_{m,h}$ for $h < i_m$. Now let μ_m be least s.t.

$\text{coll}(\omega, \mu_m) \Vdash$ (There is an N_i -conforming enlargement IF of $\mathcal{Y} \upharpoonright_{i_m}$ s.t. $s_m \upharpoonright_{i_m} = \text{tr}(\text{IF})$)

Then $\langle \mu_m \mid m < \omega \rangle$ is W_i -definable from $\langle s_m \mid m < \omega \rangle$. Set:

$\mu = \text{lub}_{m < \omega} \mu_m$. Let G be $\text{coll}(\omega, \mu)$ -generic over W_i . Then in $W_i[G]$ we

find $\langle \text{IF}_m \mid m < \omega \rangle$ s.t. IF_m is an N_i -conforming enlargement of $\mathcal{Y} \upharpoonright_{i_m}$.

Define IF by: $\text{IF} \upharpoonright_{i_0} = \text{IF}_0$,

$\text{IF} \upharpoonright_{[i_m, i_{m+1})} = \text{IF}_{m+1} \upharpoonright_{[i_m, i_{m+1})}$. Then

IF is an N_i -conforming enlargement of $\mathcal{Y} \upharpoonright_j$. QED (Claim)

Now let E' be an N_i -conforming-enlargement of $\mathcal{Y} \upharpoonright_j$. Define

E'' by $E'' \upharpoonright i = E \upharpoonright i$, $E'' \upharpoonright [i, j) = E' \upharpoonright [i, j)$.
 Then E'' is an N_i -conforming enlargement
 of $\mathcal{Y}|_j$. Note that σ is a neat en-
 -largement, since each σ_n is neat.
 Set $E^* = E'' \sim \langle W_i, \sigma \rangle$. Then E^* is
 a neat enlargement of $\mathcal{Y}|_{i+1}$ by
 the above Claim. Moreover $E^* \upharpoonright i = E \upharpoonright i$,
 and $\sigma \upharpoonright \pi_{i,j} = \sigma_i^E$. We note that E^* is
 semi prond since each σ_n is semi
 prond (hence $W_n \models R_n^{\sigma_n}$ is well founded
 and $W_n = \overline{W}_n^{\mu_n}$ for a $\mu_n \geq \text{rank of } \sigma_n$
 in R_n), and $W_i \models (R_i^{\sigma_i}$ is well
 founded and $W_i = \overline{W}_i^{\mu_i}$ for a
 $\mu_i \geq \text{rank of } \sigma_i$ in R_i). Thus
 $\mu_i \geq \text{rank of } \sigma$ in R_i , since
 whenever $\sigma \upharpoonright R_i \sigma'$ (i.e. $\sigma' \upharpoonright P_l \leq N_i$
 for an $l_* > j$ and $\sigma' \upharpoonright \pi_{j,l} = \sigma$),
 then $\sigma_i \upharpoonright R_i \sigma'$, since $\sigma' \upharpoonright \pi_{i,l} = \sigma_i$.
 Thus E^* is semi prond. But $W_i^{E^*} =$
 W_i is reflective and $W_h^{E^*} = W_h$ is
 reflective for all $h < i$ which
 does not die at j (since it then
 does not die at $i+1$). By
 the above Fact it then follows

That \mathbb{E}^* is proud, QED (Case 1)

Case 2 $j = h+1$.

Let $\xi = T(j)$. Then ξ is a break point at j , since if $k \in (\xi, j)$ and $T(k) \leq \xi$, then $T(k) \prec_* k$. (Otherwise $m(T(k)) = m(k) \geq m(\xi)$, since $k \in (\xi, j)$. But $m(k) \neq m(j) = m(\xi)$, since $k \not\prec_T j$ and $j \not\prec_T k$. Hence $m(k) > m(\xi)$, where $\xi \in [T(k), k)$. Hence $\neg T(k) \succ_* k$.)

Clearly $\xi \geq i$ and i survives at ξ . By the induction hypothesis $\mathbb{E} \upharpoonright i$ extends to a proud \mathbb{E}' enlarging $\gamma \upharpoonright \xi+1$ s.t.

$On_{W_l} \mathbb{E}' < On_{W_l} \mathbb{E}$ for $i \leq l < \xi$, $W_\xi^{\mathbb{E}'} = W_i^{\mathbb{E}}$, and $\sigma_\xi^{\mathbb{E}'} \pi_{i,\xi} = \sigma_i^{\mathbb{E}}$. Hence we can

assume w.l.o.g. that $\xi = i$. By the induction hypothesis we can apply (a) to \mathbb{E} , getting a proud \mathbb{E}' extending \mathbb{E} and enlarging $\gamma \upharpoonright j$ s.t. $On_{W_l} \mathbb{E}' < On_{W_l} \mathbb{E}$ in \mathbb{E}' for $\xi < l < j$. Applying

Superlemma (A) to \mathbb{E}' then gives the desired result: We obtain $\bar{\mathbb{E}}$ extending $\mathbb{E} \upharpoonright i$ s.t., letting

$\bar{t} = \langle \bar{t}, \bar{\sigma} \rangle = tr(\bar{\mathbb{E}})$, $t = \langle t, \sigma \rangle = tr(\mathbb{E})$, we have $t = \bar{t} \upharpoonright i$, $\bar{\sigma} \pi_{i,j} = \sigma_i$ ($\bar{\sigma}_h = \sigma_h^{\bar{\mathbb{E}}}$), and $\bar{W}_i = W_i$ ($\bar{W}_h = W_h^{\bar{\mathbb{E}}}$).

Then \bar{E} is semi prond, since s is a semi prond trace and (as in Case 1) $W_i = (R_i, \bar{\sigma}_i)$ is well founded and $\forall = \bar{W}^{(u)}$ for a $\mu \geq \text{rank of } \bar{\sigma}_i \text{ in } R_i$.

But W_i is reflective and $\bar{W}_h = W_h$ is reflective for all $h < i$ which do not die at $i+1$ (hence for $h < i$ which do not die at $j+1$). By the above Fact it follows that \bar{E} is prond, QED (Case 2).

This proves the Theorem.