

§4 An \mathcal{L} -Forcing Equivalent to \mathbb{N}^*

From now on assume CH. We apply the theory in §3, taking $\beta = \omega_2$, $N = H_{(2^{\omega_1})^+}$, and $M = L_{\omega_2}^A = \langle L_{\omega_2}[A], \in, A \rangle$, where $A \subset \omega_2$ is s.t. $L_{\omega_1}[A] = H_{\omega_1}$ and ω_1 is the largest cardinal in $L_{\omega_2}[A]$. \mathcal{L} is the infinitary language on N which, in addition to the usual predicate, constants and basic axioms, has the further axioms:

- $\dot{B} = \langle \delta_i \mid i < \underline{\omega} \rangle$ is monotone and cofinal in $\underline{\omega}_2$
- $\forall m \wedge n \geq m \quad \delta_m \in \underline{C}$, for every club $C \subset \omega_1$
- $\text{rng}(\pi_i, \underline{\omega}_1) =$ the smallest $X < \underline{M}$ s.t. $d_0 \cup \{d_j \mid j < i\} \cup \{\delta_n \mid n < \underline{\omega}\} \subset X$ for $i \leq \omega_1$.

Lemma 1 \mathcal{L} is consistent.

proof.

Let $\langle \delta_i \mid i < \omega \rangle$ be \mathbb{N}^* -generic over V .

Set $\dot{B} = \langle \delta_i \mid i < \omega \rangle$. Define M_i, π_i as follows:

Define $\langle d_i \mid i \leq \omega_1 \rangle, \langle X_i \mid i \leq \omega_1 \rangle$ by:

$X_i =$ the smallest $X < M$ s.t.

$$\{d_j \mid j < i\} \cup \{\delta_h \mid h < \omega\} \subset X$$

$$d_i =: \omega_1 \wedge X_i.$$

Let $\pi_{i, \omega_1} : M_i \xrightarrow{\sim} M \upharpoonright X_i$, where M_i is transitive.

Set: $\pi_{ij} = \pi_{j, \omega_1}^{-1} \cdot \pi_{i, \omega_1}$ for $i \leq j < \omega_1$.

Then $\langle H_{\omega_1}^V[B], \langle M_i \mid i \leq \omega_1 \rangle, \langle \pi_{ij} \mid i \leq j \leq \omega_1 \rangle, B \rangle$
models \mathcal{L} . QED (Lemma 1)

(Note There is a more elementary proof of Lemma 1 which makes no mention of \mathbb{N}^* .)

Set: $\mathbb{P} =: \mathbb{P}_{\mathcal{L}}$. If G is \mathbb{P} -generic, then
it adds $M^G = \langle M_i^G \mid i \leq \omega_1 \rangle$, $\pi^G = \langle \pi_{ij}^G \mid i \leq j \leq \omega_1 \rangle$
and $B^G = \langle \delta_i^G \mid i < \omega \rangle$ as defined in §3.

We call $B = \langle \delta_i \mid i < \omega \rangle$ a \mathbb{P} -generic sequence

iff $B = B^G$ for a \mathbb{P} -generic G . We shall
show:

Theorem $B = \langle \delta_i \mid i < \omega \rangle$ is \mathbb{P} -generic
iff it is \mathbb{N}^* -generic.

The proof stretches over many lemmata.
We shall, in fact, define a forcing \mathbb{N}^{**}
in which \mathbb{N}^* lies dense. We then prove
the theorem with \mathbb{N}^{**} in place of \mathbb{N}^* .

We first note that we may have to deal
with structures $\langle M, A \rangle$, where $A \subset M$ lies
in V but $\langle M, A \rangle$ is not amenable.

A number of familiar arguments do not work. However, we do have:

Lemma 2 Let $\langle a, \bar{a} \rangle \in F^P$, $\bar{M} = M_{|P|}^P$.

Assume (in any extension of V) that $\bar{\sigma} : \langle \bar{M}, \bar{a} \rangle \xrightarrow{\Sigma_0} \langle M, a \rangle$ cofinally.

Then $\bar{\sigma} : \langle \bar{M}, \bar{a} \rangle \prec \langle M, a \rangle$.

proof.

We first note that for any Σ_n (Π_n)

formula φ there is a canonical Σ_n (Π_n) formula $\bar{\varphi}$ s.t.

$$(1) \quad \forall x \in u \quad \varphi(x, \vec{y}) \iff \bar{\varphi}(u, \vec{y}).$$

holds in $\langle M, a \rangle$. This follows by iterated application of:

$$(2) \quad \begin{aligned} \bigwedge x \in u \quad \forall z \quad \psi(x, z, \vec{y}) &\iff \\ &\iff \forall w \bigwedge x \in u \quad \forall z \in w \quad \psi(x, z, \vec{y}) \end{aligned}$$

and:

$$(3) \quad \forall x \in u \quad \forall z \quad \psi(x, z, \vec{y}) \iff$$

$$\iff \forall w \quad \forall x \in w \quad \forall z \in w \quad (z \in u \wedge \psi(x, z, \vec{y}))$$

(3) is trivial, as is the direction (\leftarrow) of (2)

To prove (\rightarrow), note that, by the regularity of ω_1 , if the premise holds, then

$$\bigwedge x \in u. \forall z \in L_\delta[A] \quad \psi(x, z, \vec{y})$$

for a $\delta < \omega_1$. QED (1)

But then;

(4) (1) holds uniformly in $\langle \bar{M}, \bar{a} \rangle$,

proof.

Let \mathcal{M} be a grounded model of $\mathcal{L}(p)$.

Then $\pi_{|P|, \omega_1}^{\mathcal{M}} : \langle \bar{M}, \bar{a} \rangle \prec \langle \mathcal{M}, a \rangle$. QED

We now prove by induction on φ that if φ is Σ_n , then:

(5) $\langle \mathcal{M}, a \rangle \models \varphi[\pi(\vec{x})] \iff \langle \bar{M}, \bar{a} \rangle \models \varphi[\vec{x}]$.

We proceed by induction on n . The case $n=0$ is trivial. Now let $n=m+1$, $\varphi = \forall z \psi(z, \vec{x})$, where ψ is Π_m .

It remains to show:

Claim $\langle \mathcal{M}, a \rangle \models \varphi[\pi(\vec{x})] \rightarrow \langle \bar{M}, \bar{a} \rangle \models \varphi[\vec{x}]$

proof.

Assume $\langle \mathcal{M}, a \rangle \models \varphi[\pi(\vec{x})]$. Pick $\gamma < \beta \equiv \beta_{|P|}$ big enough that:

$$\langle \mathcal{M}, a \rangle \models \forall z \in \pi(u) \psi[z, \pi(\vec{x})]$$

where $u = L_{\gamma}^{\bar{a}}$. Then by (1):

$$\langle \mathcal{M}, a \rangle \models \bar{\psi}[\pi(u), \pi(\vec{x})]$$

where $\bar{\psi}$ is Π_m . Hence:

$$\langle \bar{M}, \bar{a} \rangle \models \bar{\psi}[u, \vec{x}]$$

Hence $\langle \bar{M}, \bar{a} \rangle \models \forall z \in u \psi[z, \vec{x}]$

QED (Lemma 2)

We now define:

Def $IN^{**} =$: the set of $T \in (\omega_2)^{<\omega}$ s.t. for every $t \in T$ there is $T' \leq T$ s.t. $t \in \text{stem}(T')$ and $T' \in IN^*$.

Since IN^* is dense in IN^{**} we have:

$$BA(IN^*) \simeq BA(IN^{**}).$$

Def Let $p \in IP$.

$$T_p =: \left\{ s \in (\omega_2)^{<\omega} \mid \text{con} \left(\mathcal{L} + \Phi_p + \bigwedge_{i < |s|} s_i = \delta_i^* \right) \right\}$$

Lemma 3 $T_p \in IN^{**}$

proof

Let $s \in T_p$. For each $n \geq |s|$ define a game G_n as follows. In the i -th move:

I play C_i where $C_i \subset \omega_2$ is club and $C_i = \omega_2$ if $i < n$.

II plays $\delta_i \in C_i$ s.t. $\langle \delta_0, \dots, \delta_i \rangle \in T_p$ and $\delta_i = s(i)$ if $i < |s|$.

I then wins at i iff II has no play.

II wins iff I does not win at any i .

Claim II has a winning strategy for some game G_n .

proof

Suppose not. Then I has a winning strategy S^n for each $n \in [|\omega|, \omega)$.

Set $C = \{ \lambda < \omega_2 \mid \Lambda_m < \omega \wedge \Lambda_n \in (\lambda)^{<\omega} \wedge \lambda \in S^m(\lambda) \}$,

Then C is club in ω_2 ,

Hence there are $p' \leq p, m < \omega$ s.t.

$$p' \Vdash \Lambda_i \geq m \ \dot{\delta}_i \in \underline{C} \wedge \Lambda_i < |\lambda| \ \dot{\delta}_i = \underline{\lambda(i)}.$$

Successively choose δ_i ($i < \omega$) s.t.

$$\text{con}(\mathcal{L}(p') + \bigwedge_{n < i} \delta_n^i = \underline{\delta}_n). \text{ Then}$$

$\langle \delta_i \mid i < \omega \rangle$ is a play by II in G_m

which defeats S^m . QED (Claim)

Now let II have a winning strategy

S for G_m . Let T' be the result of applying S to all possible plays of I.

Then $T' \in \mathbb{N}^{\omega}$, $\text{stem}(T') = m = |\text{stem}(T)|$

and $T' \leq T_p$. QED (Lemma 3)

Lemma 4 Let $\langle \bar{\delta}_0, \dots, \bar{\delta}_i \rangle = b_p \uparrow i+1$, Let $\langle \delta_0, \dots, \delta_i \rangle \in T_p$.

There is a unique $\pi^i: J_{\bar{\delta}_i}^{AP} \hookrightarrow J_{\delta_i}^A$ defined by:

Let f be M -least s.t. $f: \omega_1 \xrightarrow{\text{onto}} \delta_i$

Let \bar{f} be $M_{|P|}^P$ -least s.t. $\bar{f}: \alpha_i^P \xrightarrow{\text{onto}} \bar{\delta}_i$

Then $\pi^i(\bar{f}(\bar{\zeta})) = f(\zeta)$ for $\zeta < \alpha_i$.

Moreover:

• $\pi^h = \pi^i \uparrow J_{\bar{\delta}_h}^{AP}$ for $h < i$ ($M_{|P|}^P = J_{\beta_p}^{AP}$)

• If $\langle a, \bar{a} \rangle \in F^P$, then

$\pi^i: \langle J_{\bar{\delta}_i}^{AP}, \bar{a} \cap J_{\bar{\delta}_i}^{AP} \rangle \hookrightarrow \langle J_{\delta_i}^A, a \cap J_{\delta_i}^A \rangle$.

proof.

It is clear that π^i , if it exists, is uniquely defined by: $\pi^i \circ \bar{f} = f$. To show existence.

let \mathcal{M} be a grounded model of $\mathcal{L} + \mathcal{C}_p$ s.t.

$\pi_{i, \omega_1}^{\mathcal{M}}(\bar{\delta}_h) = \delta_h$ for $h \leq i$. Then $\pi_{i, \omega}^{\mathcal{M}}(\bar{f}) = f$.

Hence $\pi^i = \pi_{i, \omega_1}^{\mathcal{M}} \uparrow J_{\bar{\delta}_i}^{AP}$ has all of the

above properties. QED (Lemma 4)

Lemma 5. Let $T' \leq T_p$ in \mathbb{N}^{**} . There is $q \in P$ in IP s.t. $q \Vdash \langle \delta_i: i < \omega \rangle$ is a branch in T' .

proof.

Let G be \mathbb{N}^{**} -generic, $T' \in G$. Let

$\langle \delta_i: i < \omega \rangle$ be the generic sequence given by G . Since $\langle \delta_i: i < \omega \rangle$ is a

branch in $T' \subset T_p$, there is for each $i < \omega$ the unique $\pi^i: J_{\bar{y}_i}^{A^P} \prec J_{y_i}^A$ defined as in Lemma 3. Then,

(1) $\pi^i \subset \pi^j$ for $i \leq j < \omega$,

Proof.

Let f_i be M -least s.t. $f: \omega_1 \xrightarrow{\text{onto}} y_i$. Let \bar{f}_i be M -least s.t. $f: \omega_1^P \xrightarrow{\text{onto}} \bar{y}_i$. Clearly,

if $\bar{f}_i \in L_{\bar{y}_i}^{A^P}$, then $f_i \in L_{y_i}^A$, since by the definition of T_p there is (via sufficient generic collapse) a grounded model M

of $L(p) + \bigwedge_{m \leq i} \pi_{\omega_1}^i(\bar{y}_m) = y_m$. But then,

letting $\pi = \pi_{\omega_1}^M$, we have

$\pi(\bar{f}_i) = f_i$. Hence $\pi \upharpoonright L_{\bar{y}_i}^{A^P}$ and

$\pi \upharpoonright L_{y_i}^A \supset \pi^i$. QED (1)

Set $\pi = \bigcup_i \pi^i$. Then $\pi: \langle \bar{M}, \bar{a} \rangle \xrightarrow{\Sigma_0} \langle M, a \rangle$

whenever $\bar{M} = M_{|p|}^P$ and $\langle a, \bar{a} \rangle \in F^P$,

Hence $\pi: \langle M, \bar{a} \rangle \prec \langle M, a \rangle$ by Lemma 2. QED (1)

Let $d = d_p$. Clearly $\text{rng}(\pi)$ is the smallest $X \prec M$ s.t. $d \cup \{y_i : i < \omega\} \subset X$.

(We set: $M_i^P = L_{d_i^P}^{A^P}$, $d_p = d_{|p|}^P$, $A^P = A_{|p|}^P$)

Define $\langle d_i \mid i \leq \omega_1 \rangle, \langle X_i \mid i \leq \omega_1 \rangle$ by:

$X_i =$ the smallest $X \prec M$ s.t.

$$d \cup \{d_j \mid j < i\} \cup \{\delta_m \mid m < \omega\} \subset X$$

$$d_i =: \omega_1 \cap X_i.$$

Clearly $X_i = \text{rng } \pi \circ \pi_{i, |\rho|}^P$ for $i \leq |\rho|$.

Define $\langle M_i \mid i \leq \omega_1 \rangle, \langle \pi_{i, j} \mid i \leq j \leq \omega_1 \rangle$

by: $\pi_{i, \omega_1} : M_i \xrightarrow{\cong} \langle X_i, \in, A \cap X_i \rangle$ when M_i is transitive

$$\pi_{i, j} = \pi_{j, \omega_1} \circ \pi_{i, \omega_1} \text{ for } i \leq i' \leq \omega_1,$$

Then $\dots, \pi = \pi_{|\rho|, \omega_1}$ and:

$$M^P = \langle M_i \mid i \leq |\rho| \rangle, \pi^P = \langle \pi_{i, j} \mid i \leq j \leq |\rho| \rangle.$$

For $a \in R^P$ set: $a_i = (\pi_{i, \omega_1}^{-1})'' a$ ($i \leq \omega_1$).

Then $\exists \lambda \geq |\rho|$ s.t. $\langle \bar{a}, a \rangle \in R^P$.

Clearly $\pi' \subset M$. For $i \leq \omega_1$ set:

$T'_i = (\pi_{i, \omega_1})^{-1}'' T'$. There is a least

$\lambda \geq |\rho|$ s.t. $\pi_{\lambda, \omega_1} : \langle M_\lambda, a_\lambda \rangle \prec \langle M, a \rangle$

for all $a \in R^P$ and $\pi_{\lambda, \omega_1} : \langle M_\lambda, T'_\lambda \rangle \prec \langle M, T' \rangle$.

Set $g_0 = \langle M \cap (\lambda+1), \pi \cap (\lambda+1)^2, \langle \delta_m^\lambda \mid m < \omega \rangle \rangle$

$$g_1 = \{ \langle a, a_\lambda \rangle \mid a \in R^P \} \cup \{ \langle T', T'_\lambda \rangle \},$$

Then $q \in \mathbb{P}$, since:

$\langle H_\theta, \langle M_i \mid i \leq \omega_1 \rangle, \langle \pi_i \mid i \leq i \leq \omega_1 \rangle, \langle \delta_m \mid m < \omega \rangle \rangle$
is a model of $\mathcal{L}(q)$ for sufficiently large regular θ in $V[G]$. Clearly, then, $q \leq p$ in \mathbb{P} .

Claim Let $G' \ni q$ be \mathbb{P} -generic. Set $\delta'_i = \delta_i^{G'}$ ($i \leq \omega$). Then $\langle \delta'_i \mid i < \omega \rangle$ is a branch through T' .

proof.

Since $\pi_{\lambda, \omega_1} : \langle M_\lambda, T'_\lambda \rangle \prec \langle M, T' \rangle$

and $\pi_{\lambda, \omega_1}(\delta_i^{\lambda}) = \delta_i$ for $i < \omega$, where $\langle \delta_i \mid i < \omega \rangle$ is a branch through T' ,

we know that $\langle \delta_i^{\lambda} \mid i < \omega \rangle$ is a branch through T'_λ - i.e. $\delta^{\lambda} \upharpoonright m \in T'_\lambda$ for $m < \omega$.

Since $\pi_{\lambda, \omega_1}^G : \langle M_\lambda, T'_\lambda \rangle \prec \langle M, T' \rangle$,

it follows that:
 $\delta' \upharpoonright m = \pi_{\lambda, \omega_1}^G(\delta^{\lambda} \upharpoonright m) \in T'$

for $m < \omega$. QED (Lemma 5)

Note that if q, T' are as in Lemma 5, then $T_q \subseteq T$. (Otherwise there is $s \in T_q$ s.t. $s \notin T'$,

But then there is $q' \leq q$ s.t.

$q \Vdash \langle \dot{x}_i, i \leq n \rangle = s$, where $n = |s|$. Let $G \ni q'$ be \mathbb{P} -generic, Then $q \in G$ and $\langle \dot{x}_i, i < \omega \rangle$ is not a branch in T' . Contradiction!)

Corollary 6 Let Δ be strongly dense in $\mathbb{N}^{*\omega}$,

Then $\{p \mid T_p \in \Delta\}$ is dense in \mathbb{P} .

proof.

Let $T' \in \Delta$, $T' \in \Delta$, Pick $q \leq p$ s.t. $T_q \subseteq T'$. Then $T_q \in \Delta$. QED (Cor 6)

But then;

Corollary 7 Let $B = \langle \dot{x}_i, i < \omega \rangle$ be \mathbb{P} -generic,

Then B is $\mathbb{N}^{*\omega}$ -generic.

proof.

Let $B = B^G$ where G is \mathbb{P} -generic. Then

$G' = \{T \in \mathbb{N}^{*\omega} \mid \forall p \in G T_p \subseteq T\}$ is

$\mathbb{N}^{*\omega}$ -generic. But B is a branch in

every $T \in G'$. Hence B is the $\mathbb{N}^{*\omega}$ -generic

sequence given by G' . QED (Cor 7)

It remains only to prove the converse:

By a virtual repetition of the proof of Lemma 5 we have:

Lemma 8 Let $T \in \mathbb{N}^{**}$, There is $p \in \mathbb{P}$ s.t.,
 $T_p \leq T$.

proof:

Let $G \ni T$ be \mathbb{N}^{**} -generic. Let

$\langle \delta_i \mid i < \omega \rangle$ be the generic sequence given by G . Define d_i, X_i, M_i, π_{ij} ($i \leq j \leq \omega_1$) exactly as before but without reference to a previously chosen $p \in \mathbb{P}$. (i.e. we set:

$X_i =$ the smallest $X \leq M$ s.t.

$$\{\delta_j \mid j < i\} \cup \{\delta_n \mid n < \omega\} \subset X.$$

As before there is $\lambda < \omega_1$ s.t.

$$\pi_{\lambda, \omega_1} : \langle M_\lambda, T_\lambda \rangle \prec \langle M, T \rangle$$

where $T_\lambda = \pi_{\lambda, \omega_1}^{-1} \restriction T$. Define

$$P = \langle \langle M_i \mid i \leq \lambda \rangle, \langle \pi_{ij} \mid i \leq j \leq \lambda \rangle, B_\lambda \rangle$$

$$\text{where } B_\lambda = \pi_{\lambda, \omega_1} \restriction \{\delta_n \mid n < \omega\}$$

$$P_1 = \{\langle T, T_\lambda \rangle\}.$$

The rest of the proof is as before.

QED (Lemma 8)

Lemma 9 Let $B = \langle \delta_i : i < \omega \rangle$ be \mathbb{N}^{**} -generic.
Then $\langle \delta_i : i < \omega \rangle$ is \mathbb{P} -generic.
proof,

Suppose not. Then there is $T \in \mathbb{N}^{**}$ s.t.
 $T \Vdash \dot{B}$ is not \mathbb{P} -generic. Let $p \in \mathbb{P}$ s.t.
 $T_p \leq T$. Let $G \ni p$ be \mathbb{P} -generic. Let
 $B = B^G$. Then B is \mathbb{N}^{**} -generic and,
in fact, $B = \dot{B}^{G'}$ where $G' = \{T \mid \forall p \in G \ T_p \leq T\}$
is \mathbb{N}^{**} -generic. But $T \in G'$ and B is
 \mathbb{P} -generic. Contradiction! QED (Lemma 9)

This proves the theorem.

Note Carrying this further, we could
show that $BA(\mathbb{N}^*) = BA(\mathbb{P})$.

Note If \mathcal{L}' is like \mathcal{L} except that the
axiom:

$$\forall m \wedge n \geq m \ \delta'_m \in \bar{C} \quad \text{for all club } C \subset \omega_2$$

is replaced by:

$$\forall m \wedge n \geq m \ \delta'_{n+1} > F(\delta'_m) \quad \text{for all}$$

$$F: \omega_2 \rightarrow \omega_2,$$

then a virtual repetition of our proof
shows:

Theorem $B' = \langle \delta'_i : i < \omega \rangle$ is $\mathbb{P}_{\mathcal{L}'}$ -generic
iff it is \mathbb{N}' -generic.