

§ 4

e - premiss

Def Let N be an active ppm.

N is special iff there is $\mu \in [\kappa_N, \lambda_N)$ which is Woodin in N .

If M is a ppm and $E_\nu^M \neq \emptyset$, we call ν special in M iff $M \parallel \nu$ is special.

Lemma 1.1 Let N be active. The following are equivalent:

(a) N is special

(b) $\bigcup_{\kappa_N}^{E^N} \neq \emptyset$ there are arbitrarily large Woodins

(c) $\bigcup_{\lambda_N}^{E^N} \neq \emptyset$ " " " " " "

The proof is left to the reader.

Def Let N be active.

$E_N \subseteq (\kappa_N, \lambda_N]$ is defined as follows:

- If N is special, set: $E_N =$ the set of $\delta \in (\kappa_N, \lambda_N)$, s.t. δ is a limit cardinal in N and $\forall \mu (\mu \leq \delta \text{ in } N \text{ for all } \mu \in (\kappa_N, \delta))$.
- If N is not special, set: $E_N =$ the set of $\delta \in (\kappa_N, \lambda_N)$ s.t. δ is a limit cardinal in N .

If M is a ppm and $E_\nu^M \neq \emptyset$, set:

$$E_\nu^M = E^M(\nu) = E_{M \parallel \nu}.$$

This gives us the class of E - premiss.

The iteration indices are, of course, defined by:

Def Let N be active.

$$e_N = \begin{cases} \text{the least } e \in E_N \text{ s.t. } e \geq \kappa_N, \\ \text{if such } e \text{ exists;} \\ \kappa_N \text{ if not} \end{cases}$$

Following §3 we then define:

Def $e_{\nu}^M = e_{\nu}^M = e_{M \parallel \nu}$ if $E_{\nu}^M \neq \emptyset$

Def Let M be an \mathbb{E} -premouse.

M is an e -premouse iff

(a) $e_{\nu} \in E_{\nu}$ in M whenever $E_{\nu}^M \neq \emptyset$

(b) Let $E_{\nu}^M \neq \emptyset$. Let β be a cardinal in $J_{\nu}^{E^M}$ s.t. $\tilde{\beta} = \beta^+ \leq \kappa_{\nu}^M$ in $J_{\nu}^{E^M}$. Then there is $\bar{\nu} \in E(\beta, \tilde{\beta})$ s.t. $E_{\bar{\nu}}^M \neq \emptyset$ and $\text{crit}(E_{\bar{\nu}}^M) = \text{crit}(E_{\nu}^M)$.

e -premouse are closely related to the "domestic premouse" studied by Neeman and Steel in [1]. We define C_N, \tilde{C}_N for e -premouse as usual and get just as in §3:

Lemma 1.2 Let N be an active e -premouse. Then $C_N = \tilde{C}_N = \{ \mathbb{Z} \in \tilde{C}_N \mid N_{\mathbb{Z}} \in N \}$.

By an e -iteration we understand an \mathbb{E} -iteration of an e -premouse.

If $\mathcal{Y} = \langle \langle N_i \rangle, \dots, T \rangle$ is such an iteration, we, as usual, set: $e_i = e_{\nu_i}^{M_i}$, $e_i^+ = e_i^+(\nu_i)^{M_i}$.

We say " i is special in \mathcal{Y} " (or " ν_i is special")

to mean: ν_i is special in M_i .

Lemma 2.1 Let $\mathcal{Y} = \langle \langle N_i \rangle, \dots, \tau \rangle$ be a normal e -iteration. If $i < j$, i is special and $\kappa_j \leq \kappa_i$, then j is special.

prf.

κ_i is a limit of Woodin in $J_{\kappa_i}^{EN_i}$. If $\kappa_j = \kappa_i$, then κ_j is a limit of Woodin in $J_{\kappa_j}^{EN_j}$. If $\kappa_j < \kappa_i$, then some $\mu \in [\kappa_j, \kappa_i)$ is Woodin in $J_{\kappa_j}^{EN_j}$, hence in $J_{\kappa_j}^{EN_i}$. QED (2.1)

Lemma 2.2 Let \mathcal{Y} be as above. If $i < j$, j is special, and $\kappa_i \leq \kappa_j < e_i$, then i is special.

proof.

By Lemma 2.1 if $\kappa_i = \kappa_j$. Otherwise there is $\mu \in [\kappa_i, e_i)$ which is Woodin in $J_{\kappa_j}^{EN_j}$, hence in $J_{\kappa_j}^{EN_i}$. QED (2.2)

Lemma 2.3 Let \mathcal{Y} be as above. Let $i < j$ s.t. i is special. Then $\kappa_j \notin (\kappa_i, e_i)$

proof. Exactly like § 3 Lemma 2.

Cor 2.4 Let \mathcal{Y} be as above. Let $i < j \leq k$ s.t. j is special. Then $\kappa_k \notin (\kappa_i, e_i)$.

proof of Cor 2.4

Suppose not, $\kappa_j \neq \kappa_i$, since otherwise $\kappa_k \in (\kappa_j, e_j)$. Hence $\kappa_j \geq e_i$, since otherwise i is special by Lemma 2.2. Hence $\kappa_i < \kappa_k < e_i \leq \kappa_j$. Hence k is special by Lemma 2.1. Hence i is special by Lemma 2.2.

QED (Cor 2.4)

Def $\Gamma(\gamma) =$ the set of $i < \text{lh}(\gamma)$ s.t. for all $j < \text{lh}(\gamma)$, $i \leq j \rightarrow i \leq_{\gamma} j$.

Def γ is e-linear iff whenever $\lambda \leq \text{lh}(\gamma)$ is a limit of special points, then $\Gamma(\gamma \upharpoonright \lambda)$ is cofinal in λ .

Lemma 2.5 Let γ be a normal e-iteration. The following are equivalent.

- (a) γ is e-linear
- (b) For all special $\gamma < \text{lh}(\gamma)$ the set $\{i \mid \Gamma(i+1) \leq \gamma < i\}$ is finite.

proof. Just like the proof of the corresponding Fact in §3.

Lemma 2.6 Let γ be a normal e-iteration of length $k+1$. Then γ is e-linear.

proof. Like §3 Lemma 3.

Def Let $\gamma = \langle \langle N_i \rangle, \dots, T \rangle$ be a putative normal e -iteration. i is prominent in γ iff one of the following holds:

- (a) i is special
- (b) $i = l+1$ where l is special
- (c) For all $j \in D_{\tau}(i)$ there is $k \in D_{\tau}(i+1)$ s.t. some $\kappa \geq e_j$ is Woodin in $J_{e_k}^{EN_k}$.

Def $pr(\gamma) =$ the set of prominent $i < lh(\gamma)$
 $pr^*(\gamma) = \{i \mid \forall j \in pr(\gamma) \ i \leq_{\tau} j\}$

Lemma 2.7 $0 \in pr(\gamma)$, $pr(\gamma)$ is closed in $lh(\gamma)$.

pf. (w.l.o.g. let γ be direct)

$0 \in pr(\gamma)$ is trivial. Now let $\lambda < lh(\gamma)$ be a limit pt. of $pr(\gamma)$. Let $j < \lambda$. Let $j < \gamma \in \Gamma(\gamma|\lambda)$. Let $\gamma < i \in pr(\gamma)$. Some $\kappa \geq e_j$ is Woodin in $J_{e_{\kappa}}^{EN_{\kappa}}$ since either i satisfies (c) above or some $k \in \{i-1, i\}$ is special and $\kappa \geq e_j$ is a limit of Woodins in $J_{e_k}^{EN_k}$. QED (2.7)

Lemma 2.8 If i is prominent, $j \geq i$, and $T(i+1) < i$, then j is special.
pf. Straightforward.

Cor 2.9 The following are equivalent:

- (a) γ is e -linear
- (b) Whenever $\lambda \in lh(\gamma)$ is a limit of $pr(\gamma)$, then $\Gamma(\gamma|\lambda)$ is cofinal in λ

pf. If $i < \lambda$ is prominent and no $k \in \Gamma(\gamma|\lambda)$ is special, then $i \in \Gamma(\gamma|\lambda)$ by Lemma 2.8. QED

Lemma 2.10 Let $i = \max pr(\gamma)$. Then $i \in \Gamma(\gamma)$

pf. Straightforward by Lemma 2.8.

Now let θ be a strong Mahlo cardinal (i.e. θ is strongly inaccessible and Mahlo).

We regard V_θ as our universe and construct an array $\langle N_i \mid i \leq \theta \rangle$ of e -premises. $N_0 = K^e$ will then be the "weasel" resulting from this process. We shall prove that each N_i is weakly normally iterable. (This is the condition for N_i 's existence.

We again believe it can be shown that each N_i is a weak mouse, but haven't checked the details.) We define the sequence by specifying when a new extender is added.

Def Let $N = \langle J_\alpha^E, F \rangle$ be an active ppm. $\langle M, G \rangle$ is a background certificate for N iff

(i) M is a transitive ZFC-model; $\forall \kappa \in M$, where $\kappa = \text{crit}(F)$

(ii) G is an extender at κ on M and $\text{lh}(G) > e_N$

(iii) $\exists \pi: M \xrightarrow{G} M'$, Then $V_{e_N+2} \subset M'$

(iv) $F(x) = G(x) \cap e_N$ for $x \in \text{Pl}(M) \cap N \cap M$.

Def N is certifiable iff for every $A \in \kappa$ there is a certificate $\langle M, G \rangle$ s.t. $A \in M$.

Note This is a somewhat stronger notion of certifiability than is used in [MS]. It is taken from a set of handwritten notes written earlier by Steel, which formed the basis for [MS]. Under this definition certifiability implies that κ is regular in V . We adopted the stronger notion of certifiability for two reasons:

(a) It seems likely that the stronger notion will give us the full weak iterability of each N_i , though we don't prove it here.

(b) We need the stronger notion to prove the cardinal preservation properties of K^e at the end of this section.

We then complete the definition of $\langle N_i \mid i < \theta \rangle$ by setting:

Let $N_i = \langle J_\alpha^E, E_{w\alpha} \rangle$ be defined and weakly normally e -iterable,

If $E_{w\alpha} = \emptyset$ and there is $N' = \langle J_\beta^{E'}, F \rangle$ s.t. N' is an e -premouse, $J_\alpha^{E'} = J_\alpha^E$,

$F \neq \emptyset$, and $\alpha = d_{N'}$ and N' is

certifiable, set $N_{i+1} = N'$ for

some such N' , selecting N' to be of type B if possible.

If there is no such N' , set $N_{i+1} = \langle J_{d+1}^E, \emptyset \rangle$

If N_i is not weakly normally e -iterable, then N_{i+1} is undefined.

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We shall prove that N_i is defined for $i \leq \theta$.

Def Let $\delta: \mathcal{Q} \rightarrow N_\gamma$ where \mathcal{Q} is a countable e -premouse. Let $\mathcal{Y} = \langle \langle \mathcal{Q}_i \rangle, \dots, \tau \rangle$ be a countable putative normal e -iteration of \mathcal{Q} . Let b be a branch in \mathcal{Y} .

By a realization of b we mean $\langle \delta_i : i \in b \rangle$ s.t. for all $i \in b$:

(a) $\delta_i : \mathcal{Q} \rightarrow \sum_0^{(m)} N_{\gamma_i}$ ^{in fine} \forall for all m s.t.

$$\text{lub}_{h \in b \cap i} \kappa_h \leq \omega \rho_{\mathcal{Q}_i}^m$$

(b) If i is simple above $h \in b$, then

$$\delta_i = \delta_h, \delta_i \pi_{hi} = \delta_h$$

(c) Let $i = j+1, h = \tau(i)$. Let $\gamma_i = \bar{\beta}_m(\mathcal{Q}_h, \nu_h)$.

Set: $\delta' = \gamma_m [\delta_h, \delta_h(\nu_h)]$ and

$\sigma^{(m)} = \sigma^{(m)} [\delta_h, \delta_h(\nu_h)]$. Then $\delta_i = \delta'$

and $\delta_i \pi_{hi} = \sigma^{(m)} \delta_h$.

(d) $\delta_0 = \delta, \delta_0 = \delta$.

Note It is clear that $\langle \delta_i : i \in b \rangle$ is determined by $\langle \delta_i : i \in b \rangle$.

Lemma 2.11 Let $\langle \delta_i : i \in b \rangle$ realize a branch b in \mathcal{Y} . Let $h < i$ in b . Let $h = T(j+1)$, $j+1 \leq \frac{c}{T}$. Let $u = \text{crit}(\pi_{h,i}) = u_j$. Let $\sigma^{(m)}$ be as in (c) above (hence $\delta_{j+1} \pi_{h,j+1} = \sigma^{(m)} \delta_h$).

Then $\delta_i(u) \leq \sigma^{(m)} \delta_h(u)$, $\delta_i \upharpoonright u = \sigma^{(m)} \delta_h \upharpoonright u$, and $\delta_i(X) = \delta_i(u) \cap \sigma^{(m)} \delta_h(X)$ for $X \in \mathcal{R}(u) \cap Q_i$.

proof Exactly like (1) in the proof of §3 Lemma 4.1

Def Let $\delta, Q, \mathcal{Y}, N_\gamma, b$ be as above.

b is realizable wrt $\delta: Q \rightarrow N_\gamma$ iff b has a realization.

We call \mathcal{Y} realizable wrt. δ iff one of the following holds:

- $lh(\mathcal{Y}) = k+1$ and $\{i \mid i \leq_T k\}$ is realizable.
- $lh(\mathcal{Y})$ is a limit ordinal and \mathcal{Y} has a cofinal realizable branch.

The major lemma on realizability is due to Steel and is proven in [S]. Adapted to our circumstances it reads:

Lemma 3 Let $\delta: Q \rightarrow N_\gamma^{\text{cf}}$ where Q is a countable e -premouse. Let \mathcal{Y} be a countable putative normal e -iteration of Q . Assume that δ is $\Sigma_0^{(n)}$ -preserving whenever $T(i+1) = 0$, $Q_i^* = Q$, and $\kappa_i < \wp_Q^n$.

Either \mathcal{Y} is realizable or else \mathcal{Y} has a maximal realizable branch b , where b is of limit length.

(Hence in the latter case, b is cofinal in a $\lambda < lh(\mathcal{Y})$ and $b \neq \{i \mid i \leq_T \lambda\}$.)

Def Let $\delta: Q \rightarrow N_\gamma$ where Q is a countable e-promouse. Let γ be a countable putative normal e-iteration of Q . γ has the uniqueness property wrt $\delta: Q \rightarrow N_\gamma$ iff for all limit $\lambda < lh(\gamma)$, $\{i \mid i \leq_T \lambda\}$ is the unique realizable branch cofinal in λ .

Corollary 3.1 Let $Q, \delta, N_\gamma, \gamma$ be as in Lemma 3. If γ has the uniqueness property, then γ is realizable wrt. δ .

This is the foundation upon which we shall build.

Def Let γ be a putative normal e-iteration of Q . Let $\delta: Q \rightarrow \sum^* N_\gamma$ be fine.

By a support for γ wrt δ, γ we mean a sequence $\vec{\delta} = \langle \delta_i \mid i \in pr^*(\gamma) \rangle$ s.t.

- (a) $\langle \delta_i \mid i \leq_T i \rangle$ is a realization of $\gamma \upharpoonright i+1$
- (b) If i is special, $i < j$, $n_i = n_j$, and $n = T(i+1)$, then $\delta_{i+1}(n_j) < \delta_{i+1}(n_i)$.

The triple $\langle \gamma, \vec{\delta}, \gamma \rangle$ is then called a supported iteration.

Lemma 4.1 Let $\langle \gamma, \vec{\delta}, \delta \rangle$ be supported.
Then γ is ϵ -linear.

pf.

Let $z \in \text{lh}(\gamma)$ be special. Let $z \leq i < j$
s.t. $i, j \in X = \{i \mid \tau(i+1) \leq z < i\}$. Then
 $\kappa_i \leq \kappa_z$ since $\kappa_i \notin (a_z, d_z)$. Hence i
is special and, by the same argument,
 $\kappa_j \leq \kappa_i$. Hence if X is infinite there
must be a s.t. $X' = \{i \in X \mid \kappa_i = \kappa\}$ is infinite.
Let $i < j, i, j \in X'$. Then $\delta_{i+1}^{(k)} < \delta_{i+1}^{(k)}$.
Contr! QED (4.1)

Def Let $\mathcal{Y} = \langle \gamma, \vec{\delta}, \delta \rangle$ be supported.

Let b be a branch in \mathcal{Y} .

b is realizable in \mathcal{Y} iff b has a
realization $\langle \delta'_i \mid i \in b \rangle$ s.t. $\delta'_0: Q_0 \rightarrow N_\delta$
s.t. $\delta'_i = \delta_i$ whenever $i \in b \cap \text{pr}^*(\gamma)$.

\mathcal{Y} is realizable iff one of the
following hold:

- $\text{lh}(\gamma) = k+1$ and $\{i \mid i \leq_T k\}$ is realizable
- \mathcal{Y} is of limit length and some
cofinal branch b is realizable.

- Def $\mathcal{I} = \langle \gamma, \vec{\delta}, \gamma \rangle$ is a fine iteration iff
- (a) $\mathcal{I} \upharpoonright \xi$ is a realizable supported iteration for $1 \leq \xi \leq \text{lh}(\gamma)$,
 - (b) Let $\lambda < \text{lh}(\gamma)$, $\text{lim}(\lambda)$. Suppose there is a cofinal branch b in $\mathcal{I} \upharpoonright \lambda$ s.t. $b \neq \{i \mid i \leq \frac{\gamma}{\tau}\}$ and b is realizable in $\mathcal{I} \upharpoonright \lambda$. Let $k =$ the least $k > \lambda$ s.t. $k \in D$. Then $\tilde{\kappa}_\lambda = \sup_{i < \lambda} \kappa_i$ is Woodin in $J_{\kappa_i}^{E_{Q_i}}$.

(Note The formulation of (b) is simpler if \mathcal{I} is direct.)

Lemma 4.2 Let \mathcal{I} be a supported iteration of limit length λ and let $\mathcal{I} \upharpoonright \xi$ be fine for $\xi < \lambda$. Then \mathcal{I} is fine. Moreover \mathcal{I} extends to a fine iteration \mathcal{I}' of length $\lambda + 1$.

proof.

Case 1 $\gamma = \text{suppr}(\mathcal{I}) < \lambda$,

Then $\mathcal{I} \upharpoonright [\gamma, \lambda)$ has the uniqueness property w.t. $\delta_\gamma : Q_\gamma \rightarrow N_{\delta_\gamma}$. By Cor 3.2 $\mathcal{I} \upharpoonright [\gamma, \lambda)$ has a realizable branch b . Let $\langle \delta'_i \mid \gamma \leq i \in b \rangle$ be a realization w.t. $\delta_\gamma : Q_\gamma \rightarrow N_{\delta_\gamma}$. Set $\delta'_i = \delta_i$ for $i \leq \gamma$.

Then $\langle \delta'_i \mid i \in b \rangle$ is a realization of b in \mathcal{Y} .

Finally we note that if we set $Q'_\lambda = Q_b$, $T' \restriction \{\lambda\} = b$, then $\mathcal{Y}' = \langle \langle Q'_i \rangle, \dots, T' \rangle$ is an extension of \mathcal{Y} and $\mathcal{Y}' = \langle \mathcal{Y}', \vec{\delta}', \delta \rangle$ is a fine extension of \mathcal{Y} , where $\vec{\delta}' = \vec{\delta}$ if λ is not prominent in \mathcal{Y}' and otherwise has an obvious definition.

QED (Case 1)

Case 2 Case 1 fails.

Then $\Gamma(\mathcal{Y})$ is cofinal in λ by Lemma 4.1. Set $b = \{i \mid \forall j \in \Gamma(\mathcal{Y}) \ i \leq_T j\}$. b is the unique cofinal branch. For $j \in \Gamma(\mathcal{Y})$ there is $k \geq j$ s.t. $k \in \text{pr}^*(\mathcal{Y})$. Hence $i \leq_T k \Rightarrow i \in \text{pr}^*(\mathcal{Y})$. Thus $b \subset \text{pr}^*(\mathcal{Y})$. It is easily verified that $i < j \rightarrow \delta'_i \leq \delta'_j$ for $i, j \in b$. Hence there is $i_0 \in b$ s.t. $\delta'_i = \delta'_{i_0}$ for $i \in b \setminus i_0$. Thus $\delta'_i \cdot \pi_{i_0} = \delta'_{i_0}$ for $i \leq j, i, j \in b \setminus i_0$. Define $\delta' : Q_b \rightarrow N_{\delta'_{i_0}}$ by: $\delta' \cdot \pi_{i_0} = \delta'_{i_0}$ for $i_0 \leq i \in b$. Set $Q'_\lambda = Q_b$, $\delta'_\lambda = \delta'_{i_0}$. Then $\mathcal{Y}' = \langle \langle Q'_i \rangle, \langle v_i \rangle, T' \rangle$ is an extension of \mathcal{Y} , where $T' \restriction \{\lambda\} = b$. Clearly $\lambda \in \text{pr}(\mathcal{Y}')$, so we set: $\delta'_\lambda = \delta'$, $\mathcal{Y}' = \langle \mathcal{Y}', \vec{\delta}', \delta \rangle$ is then a fine extension of \mathcal{Y} .

QED (Lemma 4.2)

Def Let $\mathcal{Y} = \langle \gamma, \vec{\delta}, \delta \rangle$ be a fine iteration of length $k+1$. Let γ' be a putative normal e -iteration of length $k+2$ extending γ . γ' is acceptable to \mathcal{Y} (or ν'_k is acceptable to \mathcal{Y}) if ν'_k does not violate (b) in the def of "fine iteration" (or if $k \notin D$).

Lemma 4.3 Let $\mathcal{Y} = \langle \gamma, \vec{\delta}, \delta \rangle$ be a fine iteration of length $k+1$ and let γ' be an acceptable putative extension of length $k+2$. Then \mathcal{Y} has a fine extension $\mathcal{Y}' = \langle \gamma', \vec{\delta}', \delta \rangle$.
 proof. (w.l.o.g. let γ' be direct)

Case 1 k is not prominent in \mathcal{Y}' .

Let $\gamma = \max(\text{pr}(\mathcal{Y})) = \max(\text{pr}(\mathcal{Y}'))$. By the fineness of \mathcal{Y} and the acceptability of \mathcal{Y}' , $\mathcal{Y}' \upharpoonright [\gamma, k+2)$ is a normal e -iteration of length $(k+2) - \gamma$ which has the uniqueness property w.r.t., $\sigma_\gamma : Q_\gamma \rightarrow N_{\delta_\gamma}$. It follows straightforwardly by Lemma 3.1 that $\mathcal{Y}' = \langle \gamma', \vec{\delta}', \delta \rangle$ is fine. QED (Case 1)

Case 2 Case 1 fails.

If $k \notin \text{pr}(\gamma)$, extend $\vec{\sigma}$ by adding the maps σ'_h , where $\langle \sigma'_h \mid h \leq_\tau k \rangle$ is a realization of γ . Otherwise leave $\vec{\sigma}$ unchanged. Then:

Case 2.1 $\tau(k+1) = k$

Let $\gamma_k = \beta_m(Q_k, \nu_k)$. Set:

$$\gamma_{k+1} = \gamma_m[\gamma_k, \sigma_k(\nu_k)] ; \sigma^{(m)} = \sigma^{(m)}[\sigma_k(\nu_k), \sigma_k(\nu_k)].$$

$$\text{Set: } F = E_{\nu_k}^{Q_k}, F^* = \sigma^{(m)} \sigma_k(F).$$

Then F^* is ω -complete, since $\sigma^{(m)} \sigma_k(\tau_k)$ is a cardinal in $N_{\gamma_{k+1}}$.

Let $g: \lambda_k \rightarrow \sigma^{(m)} \sigma_k(\kappa_k)$ s.t.

$$\begin{aligned} \langle g(\vec{\alpha}) \rangle \in \sigma^{(m)} \sigma_k(X) &\iff \\ \iff \langle \sigma^{(m)} \sigma_k(\vec{\alpha}) \rangle \in F^*(\sigma^{(m)} \sigma_k(X)) & \\ \iff \langle \vec{\alpha} \rangle \in F(X). & \end{aligned}$$

Setting: $\sigma'_k(\pi_{k,k+1}(f)(\alpha)) = \sigma^{(m)} \sigma_k(f)(g(\alpha))$

for $g \in \Gamma(\kappa_k, Q_k^*)$, $\alpha < \lambda_k$, we get $\sigma_{k+1}: Q_{k+1} \rightarrow N_{\gamma_{k+1}}$ with the right preservation properties.

Moreover, $\sigma_{k+1} \pi_{k,k+1} = \sigma^{(m)} \sigma_k$.

Then:

(A) $\langle \delta_h \mid h \leq_T, k+1 \rangle$ realizes γ'

(B) Extend $\vec{\sigma}$ by adding δ_{k+1} if $k+1 \in \text{pr}(\gamma')$; otherwise not. Then $\gamma' = \langle \gamma', \vec{\sigma}, \delta \rangle$ is a supported iteration.

(C) γ' is fine.

Note (C) is immediate from (A), (B) and the acceptability of ν_k . To prove (B) we observe that the assumption of (b) in the def. of "support" cannot hold for $j=k$, since $T(k+1) = k$. QED (Case 2.1)

Case 2.2 The above cases fail.

Then k is special in γ' . We extend $\vec{\sigma}$ (if necessary) as before. Let

$h = T(k+1)$. Let $\gamma = \bar{\beta}_m(Q_h, \nu_h)$. Set:
 $\gamma'_{k+1} = \gamma \upharpoonright_m [\delta_h, \delta_h(\nu_h)]$, $\sigma^{(m)} = \sigma^{(m)} \upharpoonright [\delta_h, \delta_h(\nu_h)]$.

Set: $F = E_{\nu_k}^{Q_h}$; $F^* = \delta_k(F)$. Since $\tau_k \leq e_h^+$ and $J_{e_h^+}^{E^{Q_h}} = J_{e_h^+}^{E^{Q_k}}$ and e_h^+ is a cardinal in Q_k , we know that τ_k is a cardinal in Q_k . Hence F^* is ω -complete.

Claim $\delta_k(\kappa_k) \leq \sigma^{(m)} \delta_h(\kappa_k)$; $\delta_k \upharpoonright \kappa_k = \sigma^{(m)} \delta_h \upharpoonright \kappa_k$;
 $\delta_k(X) = \delta_k(\kappa_k) \cap \sigma^{(m)} \delta_h(X)$ for $X \in \mathcal{P}(\kappa_k) \cap Q_k$.

proof of Claim.

(This is exactly like the proof of the Claim in Case 2 of the proof of § Lemma 4.1.)

$\kappa_i \geq \kappa_k$ for $h \leq i < k$, since otherwise $\kappa_k \in (\kappa_i, e_i)$. Hence $T(i+1) \geq h$ for $h < i < k$.

Hence $h \leq_T k$. Let $\kappa = \text{crit}(\pi_{hk})$. Then

$\kappa = \kappa_i$ where $h = T(i+1)$, $i+1 \leq_T k$.

But $\gamma_i \leq \gamma_k$ since $\kappa_k \leq \kappa_i$. Let

$$\gamma_i = \bar{\beta}_{m+m} [Q_h, \nu_h]. \text{ Set: } \sigma^{(m+m)} = \sigma^{(m+m)} [\delta_h, \delta_h(\nu_h)].$$

Then by Lemma 2.11:

$$(*) \quad \delta_k(\kappa) < \sigma^{(m+m)} \delta_h(\kappa),$$

$$\delta_k \upharpoonright \kappa = \sigma^{(m+m)} \delta_h \upharpoonright \kappa, \text{ and}$$

$$\delta_k(X) = \delta_k(\kappa) \cap \sigma^{(m+m)} \delta_h(X) \text{ for } X \in \mathcal{P}(\kappa) \cap Q_k$$

But $\sigma^{-(m+m)} = \tilde{\sigma} \sigma^{(m)}$, where

$$\tilde{\sigma} = \sigma^{(m)} [\delta_{k+1}, \sigma^{(m)} \delta_h(\nu_h)]. \text{ Since}$$

$\bar{\tau}' = \sigma^{(m)} \delta_h(\bar{\tau}_k)$ is a cardinal in

$N_{\delta_{k+1}}$, we conclude: $\tilde{\sigma} \upharpoonright \bar{\tau}' + 1 = \text{id}$.

The conclusion follows easily.

QED (Claim)

Since F^* is ω -complete we can choose $g: \lambda_k \rightarrow \delta_k(\kappa_k)$ s.t. for all $x \in \mathbb{R}(\kappa_k) \cap \mathbb{Q}_k$

and all $\alpha_1, \dots, \alpha_n < \lambda_k$ we have:

$$\begin{aligned} \langle \vec{\alpha} \rangle \in F(x) &\leftrightarrow \langle \delta_h(\vec{\alpha}) \rangle \in F^*(\delta_k(x)) \\ &\leftrightarrow \langle g(\vec{\alpha}) \rangle \in \delta_k(x) \\ &\leftrightarrow \langle g(\vec{\alpha}) \rangle \in \sigma^{(n)} \delta_h(x). \end{aligned}$$

Setting: $\delta_{k+1}(\pi_{h,k+1}(f)(\alpha)) = \sigma^{(n)} \delta_h(f)(g(\alpha))$ for $\alpha < \lambda_k, f \in \Gamma^*(\kappa_k, \mathbb{Q}_k^*)$

we get $\delta_{k+1}: \mathbb{Q}_k^* \rightarrow N_{\delta_{k+1}}$ with the right preservation properties s.t. $\delta_{k+1} \pi_{h,k+1} = \sigma^{(n)} \delta_h$. (This is like the corresponding proof in §3.)

Thus we have:

(A) $\langle \delta_i \mid i \leq k+1 \rangle$ realizes \mathcal{Y}'

(B) Extend $\vec{\delta}$ by adding δ_{k+1} . Then $\mathcal{Y}' = \langle \mathcal{Y}', \vec{\delta}, \delta \rangle$ is a supported iteration.

(C) \mathcal{Y} is fine.

(C) follows again by (A), (B).

In proving (B) we must show that if

$h \leq j < k$ and $\kappa_j = \kappa_k$ (hence j is special), then $\pi_{h,j+1}(\kappa_j) > \pi_{h,k+1}(\kappa_k)$.

Let j be the largest such. It suffices to prove it for this j .

Claim $\delta_k(\kappa_k) = \delta_{j+1}(\kappa_k)$.

At $k = j+1$ this is immediate. Otherwise let $\kappa = \text{crit}(\pi_{j+1, k}) > \kappa_j$. By Lemma 2.11 $\delta_k \upharpoonright \kappa = \sigma^{(m)} \delta_{j+1} \upharpoonright \kappa$ where $\sigma^{(m)} = \sigma^{(m)}[\delta_{j+1}, \nu_{j+1}]$ for some m .

But τ_j a cardinal in \mathcal{Q}_{j+1} . Hence $\delta_{j+1}(\tau_j)$ is a cardinal in $N_{\delta_{j+1}}$.

Hence $\sigma^{(m)} \upharpoonright \delta_{j+1}(\tau_j) = \text{id}$. Hence

$$\delta_k(\kappa_j) = \sigma^{(m)} \delta_{j+1}(\kappa_j) = \delta_{j+1}(\kappa_j) \quad \text{QED(claim)}$$

By our construction, however,

$$\delta_{k+1}(\kappa_k) = g(\kappa_k), \text{ where}$$

$$g: \lambda_k \xrightarrow{\text{countable}} \delta_k(\kappa_k). \quad \text{QED(Lemma 4.3)}$$

Thus \forall fine iterations can always be continued, as long as the indices chosen are acceptable.

Moreover:

Lemma 4.4 Let $\mathcal{I} = \langle \mathcal{I}, \vec{s}, \mathcal{I} \rangle$ be a fine iteration of a countable e -premouse with $\text{lh}(\mathcal{I}) = \omega_1$. Then \mathcal{I} has a cofinal branch.

(Note By the regularity of ω_1 it follows that this branch is unique and that its limit model is well founded.)

proof.

Claim 1 Let g be a generic collapse of ω_1 to ω (over V). Then \mathcal{I} is realizable in $V[g]$ (wrt. $\langle N_i \rangle$ as defined in V).

proof (sketch)

If F is the top extender of N_β and $\langle M, F^* \rangle$ is a background certificate for F ,

there is a canonical F_g^* s.t.

$\langle M[g], F_g^* \rangle$ is a background certificate for F in $V[g]$. Moreover,

if $A \in V^{\text{Coll}(\omega_1, \omega)}$ s.t. $\Vdash A \subset \check{\kappa}$,

$\kappa = \text{cut}(F)$, there is $A' \subset \kappa$ s.t.

$A' \in M \rightarrow A^g \in M[g]$ for $M = \text{ZFC}$.

s.t. $\check{V}_\kappa \in M$. Hence $\langle N_i \rangle$ has the relevant properties in $V[g]$ and \mathcal{I} is still a supported iteration

in $V[g]$. We need only show

Claim \mathbb{Y} is a fine iteration in $V[g]$.

proof.

Suppose not. Then for some $\lambda < \omega_1$,

$\mathbb{Y} \upharpoonright \lambda$ has a realizable cofinal branch

in $V[g]$ w.t. $b \neq b_\lambda = \{i \mid i \leq \lambda\}$,

but has no such branch in V .

We derive a contradiction. Let

$\lambda \in X < H_{\theta^+}$ be countable. Let

$\sigma: \bar{H} \xrightarrow{\sim} X$. Then $d = \text{crit}(\sigma)$,

where $\sigma(d) = \omega_1$. Let \bar{g} be

$\text{coll}(d, \omega_1)$ -generic over \bar{H} , $g \in V$.

In \bar{H} , $\sigma^{-1}(\mathbb{Y} \upharpoonright \lambda) = \sigma^{-1}(\mathbb{Y} \upharpoonright \lambda)$ has a realizable cofinal branch b w.t.

$b \neq b_\lambda = \sigma^{-1}(b_\lambda)$. But then b is

realizable in V w.t. $\mathbb{Y} \upharpoonright \lambda$,

since if $\langle \bar{\delta}_i \mid i \in b \rangle$ is a

realization in \bar{H} , $\langle \sigma \bar{\delta}_i \mid i \in b \rangle$

is a realization of $\mathbb{Y} \upharpoonright \lambda$ and

$\sigma \bar{\delta}_i = \sigma \circ \sigma^{-1}(\bar{\delta}_i) = \bar{\delta}_i$ for $i \in p^*(\mathbb{Y} \upharpoonright \lambda)$

Contr!

QED (Claim)

We now prove the lemma. Suppose not, Then $\text{pr}(\mathcal{Y})$ is not cofinal in ω_1 by e -linearity. Let $\gamma = \sup \text{pr}(\mathcal{Y})$.

Then for each limit $\lambda \in (\gamma, \omega_1)$, $b_\lambda = \{i \mid i \leq_\tau \lambda\}$ is the unique cofinal branch in $\mathcal{Y} \upharpoonright \lambda$ which is realizable wrt $\mathcal{Y} \upharpoonright \lambda$. Let $X \subset \mathbb{H}_{\theta^+}$ be countable and let $\sigma: \bar{H} \xrightarrow{\sim} X$.

Let $\bar{\mathcal{Y}} = \langle \bar{\mathcal{Y}}, \langle \bar{\delta}_i \mid i \in \text{pr}(\bar{\mathcal{Y}}) \rangle, \bar{\delta} \rangle = \sigma^{-1}(\mathcal{Y})$. Then $\sigma \bar{\delta}_i = \delta_i$ and $\text{pr}(\bar{\mathcal{Y}}) = \text{pr}(\mathcal{Y}) \subset \gamma < \alpha = \sigma^{-1}(\omega_1)$.

Let $\langle g_0, g_1 \rangle \in V$ be a generic pair of collapsing functions over \bar{H} .

Then $\bar{\mathcal{Y}}$ has a realizable cofinal branch

b in $\bar{H}[g_0]$. Let $\langle \bar{\delta}_i \mid i \in b \rangle$ realize b wrt. $\bar{\mathcal{Y}}$ in $\bar{H}[g_0]$. Then $\langle \sigma \bar{\delta}_i \mid i \in b \rangle$ realizes b wrt $\mathcal{Y} \upharpoonright \alpha$ in V . Hence

$b = b_\lambda$. By the same argument

$b_\lambda \in \bar{H}[g_1]$. But then $b_\lambda \in \bar{H}$, since

$b_\lambda \subset \bar{H}$ and $\langle g_0, g_1 \rangle$ is a generic pair.

Hence $\sigma(b_\lambda)$ is a cofinal branch

in \mathcal{Y} . QED (Lemma 4.4)

Call an e -premouse Q realizable iff there are δ, γ s.t. $\delta: Q \rightarrow N_\gamma$. We can attempt to coiterate two countable realizable premice via fine iterations. Lemma 4.4 should then give us what we need to show that the coiteration terminates at a countable stage. However, because of the acceptability restriction on the iteration indices it is not prima facie clear that such a coiteration is possible. We show that this is the case.

Def For Q an e -premouse and $e \leq \text{ht}(Q)$ set: $\nu_e \approx$ that ν s.t. $e = e_\nu^Q$

$$\tilde{E}_e = \tilde{E}_e^Q = \begin{cases} E_{\nu_e}^Q & \text{if } \nu_e \text{ exists,} \\ \emptyset & \text{if not} \end{cases}$$

Thus, if $Q \neq Q'$, there is a least $e \leq \min(\text{ht}(Q), \text{ht}(Q'))$ s.t. $\tilde{E}_e^Q \neq \tilde{E}_e^{Q'}$.

Def If $\langle \mathcal{Y}^0, \mathcal{Y}^1 \rangle$ is a pair of fine iterations of limit length λ and b^h is a realizable cofinal branch in \mathcal{Y}^h w.t. \mathcal{Y}^h ($h=0,1$), then $e(b^0, b^1) \approx$ the least $e \leq \text{ht}(Q_b^h)$ ($h=0,1$) s.t. $\tilde{E}_e^{Q_{b^0}} \neq \tilde{E}_e^{Q_{b^1}}$.

Def Let Q^0, Q^1 be realizable. By a fine coiteration of Q^0, Q^1 with coiteration indices $\langle e_i \rangle$, we mean a pair of fine iterations $\mathcal{Y} = \langle \mathcal{Y}^0, \mathcal{Y}^1 \rangle$ of common length $\theta = \text{lh}(\mathcal{Y})$ s.t.

(a) \mathcal{Y}^h is a fine iteration of Q^h .

(b) e_i that $e \leq \text{ht}(Q_i^h)$ ($h=0,1$) s.t.

$$J_e^{Q_i^0} = J_e^{Q_i^1} \text{ but } \tilde{E}^{Q_i^0} \neq \tilde{E}^{Q_i^1} \text{ for } i \leq e$$

(c) $v_i^h = v_{e_i}^{Q_i^h}$ if — exists

otherwise $i \notin D_\theta^h$

(d) If $\lim |\lambda|, \lambda < \theta$, then

$$b^h = \{i \mid i \leq_{T^h} \lambda\} \quad (h=0,1) \text{ are chosen}$$

s.t. $e(b^0, b^1)$ is minimal,

Def A fine coiteration is terminal if it cannot be extended to a coiteration of greater length.

By the previous lemmas we have:

Lemma 4.5 Let Q^0, Q^1 be realizable countable e -premise. Let $\mathcal{Y} = \langle \mathcal{Y}^0, \mathcal{Y}^1 \rangle$ be a terminal coiteration which is countable. Then $\text{lh}(\mathcal{Y}) = k+1$ and e_k is either undefined or $v_{e_k}^{Q^h}$ is not acceptable to \mathcal{Y}^h for an $h=0$ or 1 .

We now show that the second alternative cannot occur.

Lemma 4.6 Let $\mathcal{Y} = \langle \mathcal{Y}^0, \mathcal{Y}^1 \rangle$, Q^0, Q^1 be as above. Then e_k does not exist, proof. Suppose not.

Assume not. Then $k = \lambda$ is a limit ordinal and $\tilde{E}_{e_\lambda}^{Q_\lambda^h} \neq \emptyset$ but $\bar{\kappa} = \sup_{i < \lambda} \kappa_i^h$ is not Woodin in $J_{e_\lambda}^{E^{Q_\lambda^h}}$, although there is a $b \neq b_\lambda^h = \{i \mid i <_{T^h} \lambda\}$ which is cofinal and realizable wrt \mathcal{Y}^h .

For such a b we know that $\bar{\kappa}$ is Woodin in $(J_{\bar{\kappa}^+}^E)^{Q_\lambda^h}$ wrt. all $A \subset \bar{\kappa}$

wt. $A \in Q_\lambda^h \cap Q_b$. Hence there is a

b wt. $J_{e_\lambda}^{E^{Q_\lambda^h}} \neq J_{e_\lambda}^{E^{Q_b^h}}$. Hence there

is $\bar{e} < e_\lambda$ wt. $J_{\bar{e}}^{E^{Q_\lambda^h}} = J_{\bar{e}}^{E^{Q_b^h}}$ but

$\tilde{E}_{\bar{e}}^{Q_\lambda^h} \neq E_{\bar{e}}^{Q_b^h}$. Hence $J_{\bar{e}}^{E^{Q_b^h}} = J_{\bar{e}}^{E^{Q_\lambda^{h-1}}}$

and $\tilde{E}_{\bar{e}}^{Q_b^h} \neq \tilde{E}_{\bar{e}}^{Q_\lambda^{h-1}}$, violating the

minimal choice of b_λ^0, b_λ^1 . Contr!

QED (4.6)

A coiteration of countable Q^0, Q^1 can be extended either to a countable

terminal coiteration or to a coiteration of length ω_1 . The usual proofs then show:

Lemma 4.7 Let Q^0, Q^1 be countable realizable e-premices. There is no coiteration of length ω_1 .

Thus coiterations must terminate at a countable stage $k+1$, which means that e_k is not defined; hence Q_k^0 is a segment of Q_k^1 or conversely

The appropriate version for "double rooted iterations" will also follow.

This should be enough to prove the roundness and the condensation properties for N_i . That, in turn, enables us to define $M_i = \text{core}(N_i)$ and continue the array construction.

Unfortunately, however, we extravagantly made weak mousehood the official criterion for proceeding to M_i and N_{i+1} in the definition of "array".

Building upon what we have done here, it is, indeed, possible to prove that each N_i is a weak mouse. The proof requires a bit of extra work, however, and will, therefore, be relegated to an appendix.

For the moment we adopt a "quick fix" by amending the definition of "array" to provide that M_i, N_{i+1} are defined whenever:

(a) N_i is solid

(b) If $Q = \text{core}_p(N_i)$ and $p = p^m$, then

$$\left(\bigcup_{p^+} E\right)^Q = \left(\bigcup_{p^+} E\right)^{N_i}$$

(c) If $\bar{\alpha}$ is a cardinal in N_i and $\gamma < \text{ht}(N_i)$, $\omega_{p^+}^\omega = \alpha$, $\sigma : \bar{N} \xrightarrow{\Sigma^* N_i} N_i \parallel \gamma$, $\sigma \upharpoonright \bar{\alpha} = \text{id}$, where $\sigma(\bar{\alpha}) = \alpha$, then $\bar{N} = N_i \parallel \bar{\gamma}$ for some $\bar{\gamma}$.

(b), (c) are the only condensation properties we have made use of.)

Under this definition we have already shown that $K^e = N_\theta$ exists and can proceed to the study of its large cardinal properties.

Exactly as before (§ 3 Lemma 6.1) we get:

Lemma 5.1 Let $M_i = \text{core}(N_i) = \langle J_\alpha^E, \emptyset \rangle$,

There is at most one candidate

$$N = \langle J_\beta^{E'}, F \rangle, F \neq \emptyset \text{ for } N_{i+1}.$$

Cor 5.1.1 The def. of $\langle N_i \mid i < \delta \rangle$ is uniform over every V_η s.t. $\overline{V_\delta} = \delta$ and δ is a limit of inaccessible.

Lemma 5.2 Let $i \leq \infty$ and let δ be a cardinal in N_i . There is at most one \bar{N} s.t. $\bar{N} = \langle J_\beta^{E'}, F \rangle$, F is certified, $J_\alpha^{E^{N_i}} = J_\alpha^{E'}$ and \bar{N} is an e -premouse with $\alpha = d_{\bar{N}}^+$ where $\delta < \alpha < \delta^+$ in N_i .

An e -premouse is an \mathbb{E} -premouse satisfying the conditions (a), (b). As before (§3 Lemma 6.3) we get:

Lemma 5.3 Let $M_\alpha = \langle \bigcup_\alpha E, \emptyset \rangle$. Let $N = \langle \bigcup_\beta E', F \rangle$ be certifiable w.r.t. $\bigcup_\alpha E = \bigcup_\beta E'$, $\alpha = e_N^+$, and N is an \mathbb{E} -premouse satisfying (a). Then N is an e -premouse.
proof.

The proof is as before with one change:

In Case 1 it is not enough to show that \bar{F} is ω -complete; we must show that it is certified. Clearly $\bar{F} \upharpoonright \delta = F \upharpoonright \delta$. It follows easily that if $\langle Q, F^* \rangle$ is a certificate for N , with $A \in Q$, $A \subset \kappa$, then $\langle Q, F^* \rangle$ is also a certificate for \bar{N} .

In Case 2 we have $\bar{N} = \langle J_{\bar{B}}^{\bar{E}}, \bar{F} \rangle$,
 $\bar{\gamma} = e_{\bar{N}}$, where $\bar{F} \upharpoonright \bar{\gamma} = \tilde{F} \upharpoonright \bar{\gamma}$ and \tilde{F} is
 defined by: $\tilde{F} = \pi_1 \pi_0 \upharpoonright \#(u)$,
 $\pi_0 : J_C^E \xrightarrow{F} J_B^{E'}$, $\pi_1 : J_B^E \xrightarrow{E, \nu} J_{\tilde{B}}^{\tilde{E}}$.
 Let $A \in u$. We must find a
 certificate $\langle Q, F^* \rangle$ for \bar{N} with
 $A \in Q$. Let $\langle Q, F^* \rangle$ be a certificate
 for N with $A \in Q$. Let $\sigma = \sigma^* [i, \nu]$
 $\gamma = \gamma^* [i, \nu]$. Then $\sigma : N_i \upharpoonright \gamma \xrightarrow{\Sigma^*} N_\gamma$.
 Let F_1 be the top extender of N_γ .
 since $\nu \in N_i$ and $\gamma < \nu$ is a cardinal
 in N_i , we have: $\sigma \upharpoonright \gamma = id$. Let
 $\langle Q_1, F_1^* \rangle$ be a certificate for N_γ
 w.t. $B \in Q_1$, where B codes the set
 $\langle F(x) \cap \kappa_\nu \mid x \in \#(u) \cap N \rangle$. (Since $\kappa_\nu = \text{crit}(F_1)$
 $= \text{crit}(E_\nu)$ is inaccessible, we may
 assume $\bar{Q} < \kappa_\nu$.) It follows
 easily that $\langle Q, F_1^* \upharpoonright \bar{\gamma} \rangle$ is a
 certificate for \bar{N} . QED (5.3)

Virtually as before we get:

Lemma 6.1 Let $\kappa^e = \bigcup_{\theta}^E$. Let κ be Σ_2 -strong.
Then κ is E -strong.

proof like §3 Lemma 7.1 using Cor 5.2.

Lemma 6.2 Let κ be E -strong. Then
 $o(\kappa) = \infty$ in κ^e . (Hence κ is E -strong
in κ^e .)

proof. Exactly like §3 Lemma 7.2 using
the method of Lemma 5.3 to provide
the required background certificates.

Similarly:

Lemma 6.3 Let κ be Σ_2 -strong at level
 $\alpha < \theta$. Then κ is E -strong at level α , where
 $\kappa^e = \bigcup_{\theta}^E$.

Lemma 6.4 Let κ be E -strong at level
 α where $\kappa^e = \bigcup_{\theta}^E$. Then κ has the same
property in κ^e .

In addition to these results we get:

Lemma 7 Let $\alpha < \kappa$ be a limit of Woodin
cardinals. Then α is a limit of Woodin
cardinals in κ^e .

proof of Lemma 7.

Suppose not. Set:

$$\tilde{\mu} = \text{lub} \{ \mu < \alpha \mid \mu \text{ is Woodin in } \mathcal{K}^e \}$$

Let $\tilde{\mu} < \mu < \alpha$ s.t. μ is Woodin. Since μ is not Woodin in \mathcal{K}^e , there is $A \in \mathcal{P}(\mu) \cap \mathcal{K}^e$ s.t. no $\kappa < \mu$ is A -strong in J_μ^E , ($J_\mu^E = \mathcal{K}^e \mid \mu = \bigcup_{\kappa} \mathcal{K}^e$, since μ is a limit of inaccessibles.)

We assume w.l.o.g. that A also codes $E \cap J_\mu^E$ in some natural way. Let $\kappa < \mu$ be A -strong (hence E -strong) in V_μ , where $\tilde{\mu} < \kappa < \mu$. Let

$$\delta = \text{lub} \{ \delta < \mu \mid \text{There is } F \in \mathcal{K}^e \text{ s.t.}$$

$\text{lh}(F) < \mu$ and F is an A -strong extender up to δ in $\mathcal{K}^e \}$.

Then $\delta < \mu$. Let $\delta < \beta < \mu$ where β is inaccessible (in V). Let $F^* \in V_\mu$ be an extender on κ s.t. if $\pi: V \rightarrow_{F^*} W$, then $V_\beta \subset W$

and $\pi(A) \cap \beta = A \cap \beta$. (Hence $\pi(E) \cap V_\beta = E \cap V_\beta$.)

We now imitate the proof of §3 Lemma 7.2 (or Lemma 6.2 of this section). We operate as in Case 1 of that proof. (The question of whether there is $\gamma \in (\kappa, \beta)$ s.t. $o(\gamma) > \beta$ in κ^e is uninteresting, since κ is not a limit of Woodins in κ^e .) Set:

$$F = (F^* \upharpoonright \beta) \upharpoonright N_\beta. \quad (\text{Note, as before, that } N_\beta = \langle J_\beta^E, \emptyset \rangle.) \quad \text{Let:}$$

$$\pi': J_\tau^E \xrightarrow{F} J_\nu^E \quad (\tau = \kappa + \kappa^e).$$

There is $\sigma: J_\nu^E \xrightarrow{\Sigma_0} \pi(N_\beta)$ defined

$$\text{by } \sigma(\pi'(A \upharpoonright \alpha)) = \pi(A \upharpoonright \alpha) \quad (\alpha < \beta). \quad \text{Then}$$

$$\sigma \upharpoonright \beta = \text{id}. \quad \text{Set: } F' = \pi' \upharpoonright \neq(\kappa). \quad \text{Set:}$$

$$Q = \langle J_\nu^E, F' \rangle. \quad \text{We then essentially}$$

repeat the proof in §3 Lemma 7.2, Case 1, to show that $Q = \kappa^e \parallel \nu$ for

$\alpha \nu > \beta$. (The fact that κ is not

a limit of Woodins in Q enables

us to omit some of the steps.)

Thus $F = F' \upharpoonright \beta \in \kappa^e$, where F is an

ω -complete extender on κ^e . Let

$$\pi'': \kappa^e \xrightarrow{F} W'. \quad (\text{Hence } \pi' = \pi'' \upharpoonright J_\tau^E.)$$

But since $\sigma \upharpoonright \beta = \text{id}$, we have:

$$\pi''(A) \upharpoonright \beta = \pi'(A) \upharpoonright \beta = \sigma \pi'(A) \upharpoonright \beta = \pi(A) \upharpoonright \beta = A \upharpoonright \beta.$$

Contr! QED (Lemma 7)