Each  $M_i$  is a premouse for  $i < \eta$ . But this condition is uniformly  $\Pi_3(M_i)$  by Lemma 3.3.12. Hence  $M_\eta$  is a premouse. If  $M_0$  is of type 1, then  $C_{M_i} = \emptyset$ for  $i < \eta$ . But this condition is uniformly  $\Pi_2(M_i)$ ; Hence  $M_\eta$  is of type 1.

Now let  $M_0$  be of type 2 and let  $\mu_0 = \max C_{M_0}$ . Then  $M_i$  is of type 2 and  $\mu_i = \max C_{M_i}$  for  $i < \eta$ , where  $\mu_i = \prod_{0i}(\mu_0)$ . Let  $e_0 = F_0|\mu_0$  where  $M_0 = \langle J_{\nu_0}^{E_0}, F_0 \rangle$ . Then  $e_i = F_i|\mu_i$  for  $i < \eta$ , since  $e = F|\mu$  is a  $\Pi_1$  condition. Thus for  $i < \rho$  each  $M_i$  satisfies the  $\Pi_2$  condition in  $e_i, \mu_i$ :

$$e_0 = F_i | \mu_i \wedge C_{F_i} \setminus \mu_i = \emptyset.$$

Hence  $M_{\eta}$  satisfies the corresponding condition. Hence  $M_{\eta}$  is of type 2 and  $\mu_{\eta} = \max(C_{\eta})$ . Clearly  $C_{M_i} = C_{F_i} \cup \{\max C_{M_i}\}$  for  $i \leq \eta$ . Hence  $\pi_{ij}(C_{M_i}) = C_{M_i}$ .

Now assume that  $M_0$  is of type 3. Then each  $M_i(i < \eta)$  satisfies the  $\Pi_3$  condition:  $\Lambda \not\in \langle \lambda \rangle M \not\in \langle \lambda \rangle ( \not\in C_M ).$ 

$$\bigwedge \xi < \lambda_i \bigvee \zeta < \lambda_i (\xi < \zeta \in C_{M_i})$$
$$\bigwedge \zeta \in C_{M_i} \bigvee e \in J_{\lambda_i}^{E_i} e = F_i | \zeta.$$

But then  $M_{\eta}$  satisfies the corresponding conditions. Hence  $M_{\eta}$  is of type 3. QED (Theorem 3.3.23)

# 3.4 Iterating premice

#### 3.4.1 Introduction

We have stated that a mouse will be an iterable premouse, but left the meaning of the term "iterable" and "iteration" vague. Iteration turns out, indeed, to be a rather complex notion. Let us begin with the simplest example. most logicians are familiar with the iteration of a structure  $\langle M, U \rangle$ , where M is, say, a transitive ZFC<sup>-</sup> model and  $U \in M$  is a normal ultrafilter on  $\mathbb{P}(U) \cap M$ . Set:  $M_0 = M, U_0 = U$ . Applying  $U_0$  to  $M_0$  gives the ultraproduct  $\langle M_1, U_1 \rangle$ and the extension  $\Pi_{0,1} : \langle M_0, U_0 \rangle \to \langle M_1, U_1 \rangle$  by  $U_0$ . We then repeat the process at  $\langle M_1, U_1 \rangle$  to get  $\langle M_2, U_2 \rangle$  etc. After  $1 + \mu$  repetitions we get an iteration of length  $\mu$ , consisting of a sequence  $\langle \langle M_i, U_i \rangle | i < \mu \rangle$  of models and a commutative sequence  $\langle \pi_{ij} | i \leq j < \mu \rangle$  of iteration maps  $\pi_{ij} : M_i \to M_j$ . These sequences are characterized by the conditions:

- $\pi_{i,i+1}: \langle M_i, U_i \rangle \to \langle M_{i+1}, U_i \rangle$  is the extension by  $U_i$ .
- The  $\pi_{ij}$  commute i.e.  $\pi_{ij}$  = id and  $\pi_{ij}\pi_{hi} = \pi_{hj}$  for  $h \le i \le j < \mu$ .

• If  $\lambda < \mu$  is a limit ordinal, then  $M, \langle \pi_{i\lambda} | i < \lambda \rangle$  is the direct limit of:

$$\langle M_i | i < \lambda \rangle, \langle M_{ij} | i \le j < \lambda \rangle.$$

Now suppose we are given a structure  $\langle M, S \rangle$  where  $S = \{\langle X, \kappa \rangle | X \in U_{\kappa}\}$ and for each  $\kappa \in M$ , eiter  $U_{\kappa} = \emptyset$  or else  $\kappa$  is a measurable cardinal in Mand  $U_{\kappa} \in M$  is a normal ultrafilter on  $\mathbb{P}(\kappa) \wedge M$ . An *iteration* of  $\langle M, S \rangle$  then consists of sequences  $\langle \langle M_i, S_i \rangle | i < \mu \rangle$ ,  $\langle M_{ij} | i \leq j < \mu \rangle$  and  $\langle \kappa_i | i + 1 < \mu \rangle$ .

The first condition above is then replaced by:

$$\pi_{i,i+1} : \langle M_i, S_i \rangle \to \langle M_{i+1}, S_{i+1} \rangle \text{ is the extension by the ultrafilter}$$
$$U_i = \{ X | \langle X, \kappa_i \rangle \in S_i \}$$

The other conditions remain unchanged.  $\kappa_i | i + 1 \leq \mu \rangle$  is called the sequence of *indices*.  $\kappa_i$  must always be so chosen that  $U_i$  is an ultrafilter.

Note. Since we are allowed considerable leeway in the choice of the index  $\kappa_i$ , the purist may question whether the word "iteration" is still appropriate. In fact, the mathematical meaning of this word has rapidly changed as the structures to which it is applied have grown more complex.

An iteration is called *normal* iff the indices are increasing — i.e.  $\kappa_i < \kappa_j$  for  $i < j < \mu$ .

We now attempt to apply these ideas to premice. Let M be a premouse. An iteration of length  $\mu$  will yield a sequence  $\langle M_i | i < \mu \rangle$  of premice. In passing from  $M_i$  to  $M_{i+1}$  we apply any of the extenders  $E_{\nu}^M$  such that  $M_i || \nu = \langle J_{\nu}^E, E_{\nu} \rangle$  is active.  $\nu = \nu_i$  is then the *i*-th index. (It would be ambiguous to regard  $\kappa_i = \operatorname{crit}(E_{\nu_i})$  as the index, since  $M_i$  might have many extenders with this critical point.) In a normal iteration we have that, whenever i < j, then:

$$J_{\nu_i}^{E^{M_i}} = J_{\nu_i}^{E^{M_j}}$$
 and  $\nu_i$  is a cardinal in  $M_j$ .

(In fact,  $\nu_i = \lambda_i^{+}{}^{M_j}$ , where  $\lambda_i = E_{\nu_i}(\kappa_i)$  is inaccessible in  $M_j$ .) This follows easily by induction on j. It was originally envisaged that  $E_{\nu_0}$  would be applied directly to  $M_i$  to get  $M_{i+1}$ . It turns out, however, that such iterations are unsuitable for may purposes. (In particular, they are unsuited to use in *comparison iteration*, which we shall describe below.) The problem is that  $\kappa_i = \operatorname{crit}(E_{\nu_i})$  could be much smaller than  $\lambda_i$ , where  $\lambda_i = E_{\nu_i}(\kappa_i)$  is the largest cardinal in the model  $J_{\nu_i}^{E^{M_i}}$ . In particular, we might have  $\kappa_i < \lambda_h$  for an h < i. Since  $\lambda_h$  is an inaccessible cardinal in  $M_i$ , it follows by acceptability that:

$$\mathbb{P}(\kappa) \cap M_i = \mathbb{P}(\kappa) \cap J_{\lambda_h}^{E^{M_h}} \subset M_h.$$

Hence it should be possible to apply  $E_{\nu_i}$  to  $M_h$  rather than  $M_i$ . It turns out that it is most effective to apply  $E_{\nu_i}$  to the *smalles place possible*: we apply it to  $M_{T(i+1)}$ , where

$$T(i+1) =:$$
 the least h such that either  $h = i$   
or  $h < i$  and  $\kappa_i < \lambda_h$ .

This should give us

$$\pi_{h,i+1}: M_h \to M_{i+1}.$$

Here, however, we must deal with a second problem, which can arise even when T(i+1) = i. We know that  $E_{\nu_i}$  is an extender at  $\kappa_i$  on  $J_{\nu_i}^E$ . Then  $\mathbb{P}(\kappa_i) \cap J_{\nu_i}^{E^{M_i}} = \mathbb{P}(\kappa_i) \cap J_{\tau_i}^{E^{M_i}} = \mathbb{P}(\kappa_i) \cap J_{\tau_i}^{E^{M_h}}$ , where  $\tau_i = \kappa_i^{+J_{\nu_i}^E}$ . But  $M_h$ might contain subsets of  $\kappa_i$  which do not lie in  $J_{\tau_i}^E$  (hence  $\tau_i$  is not a cardinal in  $M_h$ , by acceptability).  $E_{\nu_i}$  is then only a partial function on  $M_h$  and cannot be applied to  $M_h$ . The resolution of this difficulty is to apply  $E_{\nu_i}$  to the *largest possible segment* of  $M_h$ . We set:

$$M_i^* =: M_h || \eta_h$$
, where  $\eta_i \leq \text{On}_{M_h}$  is maximal such that  $\tau_h$  is a cardinal in  $M_h || \eta$ .

By acceptability,  $\mathbb{P}(\kappa_i) \cap M_i^* = \mathbb{P}(\kappa_i) \cap J_{\tau_i}^E$  and  $\rho_{M_i^*}^{\omega} \leq \kappa_i$  if  $\eta_i < \operatorname{On}_{M_h}$ .

We then say that  $M_h$  drops (or truncates) to  $M_i^*$ , if  $M_h \neq M_i^*$ . i + 1 is then called a drop point (or truncation point).  $\pi_{h,i+1} : M_i^* \to M_{i+1}$  is then a partial map of  $M_h$  to  $M_{i+1}$ 

This means that iteration is no longer a linear process. Previously  $\pi_{ij}$  was defined whenever  $i \leq j < \mu$ ,  $\mu$  being the length of the iteration. Now it is defined only when *i* is less than or equal to *j* in a tree *T* on  $\mu$ . (We write  $i \leq_T j$  for  $i = j \lor iT_j$ .) 0 is the unique minimal point of *T*. T(i+1) is the unique *T*-predecessor of i + 1. The  $\pi_{ij}$  are partial maps and we again have:

$$\pi_{ij} \cdot \pi_{hi} = \pi_{hj}$$
 for  $h \leq_T i \leq_T j$ .

We will always have:  $iT_j \to i < j$ , but the converse may not hold. If  $\mu = \omega$ , these conditions completely define  $T \subset \omega^2$ . But how do we then extend the iteration to an iteration of length  $\omega + 1$ ? Previously we simply took a transitivized direct limit of  $\langle M_i | i < \omega \rangle$ ,  $\langle \pi_{ij} | i \leq j < \omega \rangle$ . Now we must first find a *branch* b in T which is cofinal in  $\omega$  (i.e.  $\sup b = \omega$ ). We also require that b have at most finitely may drop points. Pick any  $i \in b$  such that  $b \setminus i$  has no drop point. Then  $\pi_{hj} : M_h \to M_j$  is a total map on  $M_h$  for  $i \leq_T h \leq \varepsilon b$ .

Form the direct limit:

$$M_b, \langle \pi_{h_i} | i \le h \in b \rangle$$

of:

$$\langle M_h | i \leq h \in b \rangle, \ \langle \pi_{hj} | i \leq_T h \leq j \in b \rangle.$$

If  $M_b$  is well founded, we call b a *well founded branch* and take  $M_b$  are being transitive. We can then continue the iteration by setting:

$$M_{\omega} =: M_b; hT_{\omega} \leftrightarrow: h \in b \text{ for } h < \omega.$$

 $\pi_{j\omega}$  is then defined for  $i \leq_T j <_T \omega$ . If hTi, we set  $\pi_{h\omega} =: \pi_{j\omega} \cdot \pi_{hi}$ .

The same procedure is applied at all limit points  $\lambda$ . We then have:

- $\lambda$  is a limit point of T
- $T''\{\lambda\}$  is cofinal in  $\lambda$
- $T''\{\lambda\}$  contains at most finitely many truncation points.

By now we have almost given a virtual definition of what is meant by a "normal iteration of a premouse". The only point left vague is what we mean by "applying" the extender  $E_{\nu_i}$  to  $M_i^*$ . We shall, in fact, take the  $\Sigma_0^{(n)}$ -ultrapower:

$$\pi: M_i^* \to_{E_{u_i}}^{(n)} M_{i+1},$$

where  $n \leq \omega$  is maximal such that  $\kappa_i < \rho_{M_i^*}^n$ .

## 3.4.2 Normal iteration

We are now ready to write out the formal definition of "normal iteration". We shall employ the following notational devices:

**Definition 3.4.1.** Let T be a tree. We set:

- $i <_T j \leftrightarrow : \circ T_j$
- $i \leq_T j \leftrightarrow : i = j \lor iT_j$
- $[i, j]_T := \{h | i \leq_T h \leq_T j\}$  (similarly for  $[i, j]_T, [i, j]_T, [i, j]_T)$
- T(i) =: The immediate *T*-predecessor of *i* (if it exists).

We can now define:

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**Definition 3.4.2.** Let M be a premouse. By a *normal iteration* of M of *length*  $\mu$  we mean:

$$\langle \langle M_i | i < \mu \rangle, \langle \nu_i | i + 1 < \mu \rangle, \langle \pi_{ij} | i \leq_T j \rangle, T \rangle$$

where.

- (a) T is a tree on  $\mu$  such that  $iT_j \rightarrow j < j$
- (b)  $M_i$  is a premouse for  $i < \mu$
- (c)  $\nu_i < \nu_j$  if i < j. Moreover  $M_i || \nu_i = \langle J_{\nu_i}^E, E_{\nu_i} \rangle$  with  $E_{\nu_i} \neq \emptyset$ . (We set:  $\kappa_i =: \operatorname{crit}(E_{\nu_i}), \tau_i =: \kappa_i^+ J_{\nu_i}^E, \ \lambda_i =: E_{\nu_i}(\kappa_i) = \text{the largest cardinal in} J_{\nu_i}^E$ .)
- (d) Let h be least such that h = i or h < i and  $\kappa_i < \lambda_h$ . Then h = T(i+1)and  $J_{\tau_i+1}^{E^{M_h}} = J_{\tau_i+1}^{E^{M_i}}$ .
- (e)  $\pi_{ij}$  is a partial map of  $M_i$  to  $M_j$ . Moreover  $\pi_{ij} \circ \pi_{hi} = \pi_{hj}$  for  $h \leq_T i \leq_T j$ .
- (f) Let h = T(i+1). Set:  $M_i^* = M_h || \eta_i$ , where  $\eta_i$  is maximal such that  $\tau_i$ is a cardinal in  $M_h || \eta_i$ . Then  $\pi_{h,i+1} : M_i^* \to_{E_{\nu_i}}^{(n)} M_{i+1}$ , where  $n \leq \omega$  is maximal such that  $\kappa_i < \rho_{M_i^*}^n$ . (We call i+1 a drop point or truncation point iff  $M_i^* \neq M_h$ )
- (g) If  $k \leq_j$  and  $(i, j]_T$  has no drop point, then  $\pi_{ij} : M_i \to M_j$  is a total function on  $M_i$ .
- (h) Let  $\lambda$  be a limit ordinal. Then  $T''\{\lambda\}$  is club in  $\lambda$  and contains at most finitely many drop points. Moreover, if  $iT\lambda$  and  $(i, \lambda)_T$  is free of drops, then:

$$M_{\lambda}, \langle \pi_{j\lambda} | i \leq_T j <_T \lambda \rangle$$

is the transitivized direct limit of:

$$\langle M_j | i \leq_T j <_T \lambda \rangle, \langle \pi_{hj} | i \leq_T h \leq_T j <_T \lambda \rangle.$$

This completes the definition.

**Lemma 3.4.1.** Let  $I = \langle \langle M_i \rangle, \langle \nu_i \rangle, \langle \pi_{i,j} \rangle, T \rangle$  be a normal iteration. Then

- (a)  $J_{\nu_i}^{E^{M_i}} = J_{\nu_i}^{E^{M_{i+1}}}$
- (b) In  $M_{i+1}$ ,  $\lambda_i$  is inaccessible and  $\nu_i = \lambda_i^+$ .

**Proof:**  $\tau_i$  is a cardinal in  $M_i^*$ . Since  $\kappa_i$  is inaccessible in  $J_{\tau_i}^{E^{M_i}}$  and is the largest cardinal in  $J_{\tau_i}^{E^{M_i}}$ , it follows by acceptability that:

 $\tau_i = \kappa_i^+$  and  $\kappa_i$  is inaccessible in  $M_i^*$ 

 $F = E^{M_i}{}_{\nu_i}$  is a full extender of length  $\lambda_i$  with base  $H = |J^{E^{M_i}}_{\tau_i}|$  and extension  $\langle \pi, H' \rangle$ , where  $H' = |J^{E^{M_i}}_{\nu_i}|$ . By acceptability we have:

$$\mathbb{P}(\kappa_i) \cap M_i^* = \mathbb{P}(\kappa_i) \cap J_{\tau_i}^{E^M}$$

Hence F is an extender on  $M_i^*$  (and the condition (f) makes sense). But then  $\langle M_{i+1}, \pi_{i,i+1} \rangle$  is the  $\Sigma_i^{(n)}$ -liftup of  $\langle M_i^*, \pi \rangle$ , where n is maximal such that  $\kappa_i < \rho_{M_i^*}^n$ . Hence:

$$\pi_{i,i+1}(\tau_i) = \sup \pi^{"} \tau_i = \nu_i \text{ and } \pi_{i,i+1}(\kappa_i) = \lambda_i$$

Hence (b) holds, since the corresponding statement is function of  $\kappa_i, \tau_i$  in  $M_i^*$ .

To see that (a) holds, note that each element of H' has the form  $\pi(f)(\alpha)$ , where  $\alpha < \lambda_0$  and  $f \in H$  is a function on  $\kappa$ . But then:

$$\pi(f)(\alpha) \in E^{M_i} \longleftrightarrow \pi(f)(\alpha) \in E^{M_{i+1}} \longleftrightarrow \alpha \in \pi(X)$$

where  $X = \{\xi < \kappa_i : f(\xi) \in E^{M_i}\}$ . Hence

$$E^{M_i} \cap H' = E^{M_{i+1}} \cap H^i$$
 and  $J^{E^{M_i}}_{\nu_i} = J^{E^{M_{i+1}}}_{\nu_i}$ 

QED(Lemma 3.4.1)

Using these facts we prove:

**Lemma 3.4.2.** Let  $I = \langle \langle M_i \rangle, \langle \nu_i \rangle, \langle \pi_{ij} \rangle, T \rangle$  be a normal iteration. Let h < i. Then

- (a)  $J_{\nu_h}^{E^{M_h}} = J_{\nu_h}^{E^{M_i}}$
- (b)  $\lambda_h$  is inaccessible in  $M_i$  and  $\nu_h = \lambda_h^+$  in  $M_i$
- (c) Let  $h < j <_T i$ . Then  $\lambda_h \leq \operatorname{crit}(\pi_{j,i}) < \lambda_i$ .
- (d) Let  $h <_T i$ .  $\pi_{h,i}$  is a total function on  $M_h$  iff  $[H,i]_T$  is drop free.

The proof is by induction on i. We leave the details to the reader.

Note. h < i implies  $\nu_h < \lambda_i$ , since  $\nu_h < \nu_i$  is a successor cardinal in  $M_i$ ; hence  $\nu_h \notin [\lambda_i, \nu_i)$ .

**Definition 3.4.3.** Let  $I = \langle \langle M_i \rangle, \langle \nu_i \rangle, \langle \pi_{ij} \rangle, T \rangle$  be a normal iteration.

- lh(I) denotes the length of I
- If  $\eta \leq lh(I)$  we set:

$$I|\eta =: \langle \langle M_i | i < \eta \rangle, \langle \nu_i | i + 1 < \eta \rangle, \langle \pi_{ij} | i \leq_T i < \eta \rangle, T \cap \eta^2 \rangle.$$

**Definition 3.4.4.** Let  $I = \langle \langle M_i \rangle, \ldots, T \rangle$  be a normal iteration of limit length  $\eta$ . By a *well founded cofinal branch* in I we mean a branch b in T such that

- $\sup b = \eta$
- b has at most finitely many truncation points
- Let  $i \in b$  such that  $b \setminus i$  is truncation free. Then

$$\langle M_j | j \in b \rangle, \langle \pi_{hi} | i \leq h \leq j \text{ in } b \rangle$$

has a well founded direct limit.

We leave it to the reader to prove:

**Lemma 3.4.3.** Let  $I = \langle \langle M_i \rangle, \langle \nu_i \rangle, \langle \pi_{ij} \rangle, T \rangle$  be a normal iteration of limit length  $\eta$ . Let b be a well founded cofinal branch in I. I has a unique extension I' of length  $\eta + 1$  such that  $I'|\eta = I$  and  $T'''\{\lambda\} = b$ . (Moreover, if  $i \in b$  and  $b \setminus i$  is drop free then:

$$M'_n, \langle \pi'_{h,n} | h \in b \setminus i \rangle$$

is the transitivized direct limit of

$$\langle M_h | h \in b \setminus i \rangle, \langle \pi_{h,i} | h \in b \setminus i \rangle.$$

**Note.** We use Theorem 3.3.23 to show that  $M'_{\eta}$  is a premouse.

**Note**. It will be easier to talk about such limits if we have a notion of direct limit which can be applied to directed systems of *partial maps*. This could be defined quite generally, but the following version suffices for our purposes: Let  $S = \langle S, \langle \rangle$  be a linear ordering. Let  $\mathbb{A}_i$  be a model and let  $\pi_{ij}$  be a partial injection of  $\mathbb{A}_i$  to  $\mathbb{A}_j$  for  $i \leq j$  in S. Assume that the maps commute (i.e.  $\pi_{ij}\pi_{\kappa i} = \pi_{\kappa j}$ ) and that for sufficiently large  $i \in S$  we have:

$$\pi_{ij}$$
 is a total map on  $\mathbb{A}_8$  for all  $j \geq i$  in  $I$ .

Let S' be the set of such *i*. We call:

$$\mathbb{A}, \langle \pi_i | i \in S \rangle$$

a direct limit of:

$$\langle \mathbb{A}_i | i \in S \rangle, \langle \pi_{ij} | i \leq j \text{ in } S \rangle$$

iff:

$$\mathbb{A}, \langle \pi_i | i \in S'$$

in a direct limit of:

$$\langle \mathbb{A}_i | i \in S' \rangle, \langle \pi_{ij} | i \leq j \text{ in } S' \rangle$$

and  $\pi_h$  is defined by:  $\pi_h = \pi_i \pi_{hi}$  for  $h \notin S'_1 i \in S$ .

In §3.2 we defined  $\mathbb{N}$  to be a  $\Sigma^*$ -ultrapower of M by F with  $\Sigma^*$ -extension  $\pi$  (in symbols  $\pi : M \to_F^* N$ ) iff F is close to M and  $\pi : M \to_F^{(n)} N$  where  $n \leq \omega$  is maximal such that  $\operatorname{crit}(F) < \rho_M^n$ . Theorem 3.2.17 said that in this case  $\pi$  is  $\Sigma^*$ -preserving. We shall now show that in a normal iteration  $E_{\nu_i}^{M_i}$  is always close to  $M_i^*$ . In order to utilize the full strength of this fact, we shall formulate it not only for normal iteration, but also for potential normal iteration in the following sense:

Let  $I = \langle \langle M_i \rangle, \langle \nu_i \rangle, \langle \pi_{ij} \rangle, T \rangle$  be a normal iteration of length i + 1. If we attempt to extend I to an I' of length i + 2 by appointing the next  $\nu_i$ , we call this attempt a *potential normal iteration*. The formal definition is:

**Definition 3.4.5.** A potential normal iteration of length i + 2 is a structure

$$\mathfrak{T}' = \langle \langle M_j | j \le i \rangle, \langle \nu_j | j \le i \rangle, \langle \pi_{ij} | i \le j \le i \rangle, T' \rangle$$

where:

- $I = \langle \langle M_j \rangle, \langle \nu_j | j < i \rangle, \langle \pi_{ij} \rangle, T \rangle$  is a normal iteration of length i + 1, where  $T = T' \cap (i + 1)^2$
- $E_{\nu_i}^{M_i} \neq \emptyset$  and  $\nu_i > \nu_j$  for j < i
- $hT'j \leftrightarrow (hTj \lor (h \leq_T \xi \land j = i))$  where:

 $\xi = T'(i+1) =:$  the least  $\xi$  such that  $\kappa_i < \lambda_{\xi}$ .

If I' is a potential iteration and  $\xi = T'(i+1)$ , we define  $M_i^* = M_{\xi}U$  is in the usual way, (but we do not yet know whether  $M_i^*$  is extendable by  $E_{\nu_i}^{M_i}$ ).

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Note. (a)-(d) in the definition of normal iteration continue to hold. ((d) is trivial if  $\xi = i$ . If  $\xi < i$ , then  $\tau_i < \lambda_{\xi}$  and  $J_{\lambda_{\xi}}^{E^{M_{\xi}}} = J_{\lambda_{\xi}}^{E^{M_i}}$ ). But then  $M_i^*$  is defined and  $\tau_i \in M_i^*$  is a cardinal in  $M_i^*$ . Let  $n \leq \omega$  be maximal such that  $\kappa_i < \rho_{M_i^*}^n$ . It is easily seen that, if the  $\Sigma_0^{(n)}$  extension:

$$\pi': M_i^* \longrightarrow_{E^{M_i} \nu_i}^{(n)} M'$$

exists, we can turn I' into a normal iteration of length i + 2 by setting:

$$M_{i+1} = M', \ \pi_{\xi,i+1} = \pi'$$

We now prove a basic fact about normal iteration:

**Theorem 3.4.4.** Let I be a potential normal iteration of length i + 2. Let  $\xi = T(i+1)$ . Then  $E_{\nu_i}^{M_i}$  is close to  $M_i^*$ .

Before proving this we note the obvious corollary:

**Corollary 3.4.5.** Let I be a normal iteration. If h = T(i+1) in I, then:

$$\pi_{h,i+1}: M_i^* \to_{E_{\nu_i}}^* M_i$$

**Lemma 3.4.6.** Let I be a normal iteration. Let  $h = T(i+1), i+1 \leq_T j$ , where  $(i+1, j)_T$  has no truncation point. Then:

$$\pi_{h,j}: M_i^* \longrightarrow_{\Sigma^*} M_j$$
 strongly.

In particular  $\pi_{h,j}$   $P_{M_i}^n \subset P_{M_j}^n$  for  $\rho^{n+1} = \rho^{\omega}$  in  $M_i^*$ .

**Proof.** By induction on j using Lemma 3.2.26, Lemma 3.2.27 and Lemma 3.2.28.

QED(Lemma 3.4.6)

We shall derive Theorem 3.4.4 from an even stronger statement:

**Lemma 3.4.7.** Let I be a potential normal iteration of length i + 2. Then

$$\mathbb{P}(\tau_i) \cap \underline{\Sigma}_1(M_i || \nu_i) \subset \underline{\Sigma}_1(M_i^*).$$

We first show that Lemma 3.4.7 implies theorem 3.4.4. Since  $F = E_{\nu_i}$  is weakly amenable, we need only show that  $F_{\alpha} \in \underline{\Sigma}_1(M_i^*)$  for  $\alpha < \lambda_i$ , where:

$$F_{\alpha} = \{ x \subset \kappa_i | x \in M_i | | \nu_i \land \alpha \in F(x) \}$$

Let  $k \in M_i || \nu_i \text{ map } \tau_i \text{ onto } J_{\tau_i}^E$ . Then  $k \in M_i^*$ , since either i = T(i+1) and  $M^* \supset M_i || \nu_i$ , or else h = T(i+1) < i, whence follows:  $k \in J_{\lambda_h}^{E^{M_i}} = J_{\lambda_h}^{E^{M_i^*}} \subset M_i^*$ . Set:

$$\tilde{F}_{\alpha} = \{\xi < \tau_i | k(\xi) \in F_{\alpha}\}$$

Then  $\tilde{F}_{\alpha} \subset \mathbb{P}(\tau_i)$  is  $\underline{\Sigma}_1(M_i^*)$  by Lemma 3.4.7. Hence  $F_{\alpha} = k'' \tilde{F}_{\alpha} \in \underline{\Sigma}_1(M_i^*)$ . QED

We now prove Lemma 3.4.7. Suppose not. Let I be a counterexample of length i + 2, where i is chosen minimally. Let h = T(i + 1). Then:

(1) h < i

**Proof:** Suppose not. Then  $M_i^* = M_i || \mu$  where  $\mu \ge \nu$ . Hence  $\underline{\Sigma}_1(M_i || \nu_i) \subset \underline{\Sigma}_1(M_i^*)$ . Contradiction!

(2)  $\nu_i = \operatorname{On}_{M_i}$  and  $\rho_{M_i}^1 \leq \tau_i$ .

**Proof:** Suppose not. Let  $A \subset \tau_i$  be  $\underline{\Sigma}_1(M_i || \nu_i)$ . Then  $A \in \mathbb{P}(\tau_i) \land M_i \subset J_{\lambda_n}^{E^{M_i}}$ , since  $\lambda_h > \tau_i$  is inaccessible in ???. But  $J_{\lambda_n}^{E^{M_i}} = J_{\lambda_n}^{E^{M_i}} \subset M_i^*$ . Contradiction!

(3) i is not a limit ordinal.

**Proof:** Suppose not. Then  $\sup\{\operatorname{crit}(\pi_{li})(\leq i\} = \sup_{l < i} \lambda_l$ , so we can pick  $L \leq i$  such that  $\operatorname{crit}(\pi_{l,i}) > \lambda_h > \tau_i$  and  $\pi_{l,i}$  is a total function on  $M_l$ . Then  $\pi_{l,i} : M_l \to_{\Sigma_1} M_i$ , where  $M_i = \langle J_{\nu_i}^{E_i}, F \rangle$ , where  $F \neq \emptyset$ . Hence  $M_l = \langle J_{\overline{\nu}}^{\overline{E}}, \overline{F} \rangle$  where  $\overline{F} \neq \emptyset$ . Let  $A \subset \tau_i$  be  $\Sigma_1(M_i)$  such that  $A \notin \Sigma_1(M_i^*)$ . We can assume l to be chosen large enough that  $p \in \operatorname{rng}(\pi_{li})$ , where A is  $\Sigma_1(M_i)$  in the parameter p. Thus  $A \in \Sigma_1(M_l)$ . Clearly  $\overline{\nu} > \nu_j$  for all j < l, since  $\nu_j \in M_l = \langle J_{\overline{\nu}}^{\overline{F}}, \overline{F} \rangle$ .

Extend I|l + 1 to a potential iteration I' of cf length l + 2 by setting  $\nu_l = \overline{\nu}$ . Since  $\operatorname{crit}(\pi_{l,i}) > I_i$ , it follows easily that  $\tau'_l = \tau_i, \kappa'_l = \kappa_i$ , where  $\tau_l, \kappa'_l$  are defined in the usual way. But then  $M_i^* = (M'_l)^*$  and  $A \in \underline{\Sigma}_1(M_i^*)$  by the minimality of *i*. Contradiction! QED (3)

Now let  $i = j + 1, \xi = T(i)$ . Since  $\pi_{\xi,i} : M_j^* \to_{\Sigma_1} M_i = \langle J_{\nu_i}^E, E_{\nu} \rangle$  where  $E_{\nu_i} \neq \emptyset_i$  we have:

- (4)  $M_j^* = \langle J_{\overline{\nu}}^{\overline{E}}, \overline{E}_{\overline{\nu}} \rangle$  where  $\overline{E}_{\overline{\nu}} \neq \emptyset$ .
- (5)  $\tau_i < \kappa_j$

**Proof:**  $\tau_i < \lambda_j$  since  $\tau_i = \kappa_i^{+M_i}$  and  $\kappa_i < \lambda_h \leq \lambda_j$ , where  $\lambda_j$  is inaccessible in  $M_i$ . But obviously  $\kappa_i, \tau_i \in \operatorname{rng}(\pi_{\xi,i})$  by (4) where  $[\kappa_j, \lambda_j) \cap \operatorname{rng}(\pi_{\xi i}) = \emptyset$ . QED (5)

(6)  $\pi_{\xi i}: M_j^* \to_{E_{\nu_j}} M_i$  is a  $\Sigma_0$  ultrapower.

**Proof:** Suppose not. Then  $\kappa_j < \rho_{M_j^*}^1$ . Hence  $\pi_{\xi,i}$  is  $\Sigma_0^{(1)}$ -preserving. Hence  $\pi_{\xi i}'' \rho_{M_i^*}^1 \subset \rho_{M_i}^1$ . Hence  $\tau_i = \pi_{\xi i}(\tau_j) < \rho_{M_i}^1$ , contradicting (2). QED (6)

But then:

(7)  $\mathbb{P}(\tau_i) \cap \underline{\Sigma}_1(M_i) \subset \mathbb{P}(\tau_i) \cap \underline{\Sigma}_1(M_i^*).$ 

**Proof:** Let  $A \subset \tau_i$  be  $\Sigma_1(M_i)$  in the parameter p. Let  $p = \pi_{\xi_i}(f)(\alpha)$ , where  $f : \kappa_i \to M_i^*, f \in M_i^*$ , and  $\lambda < \lambda_j$ . Then

$$A(\xi) \leftrightarrow \bigvee xA'(\zeta, x, p)$$

where A' is  $\Sigma_0(M_i)$ . Let  $\overline{A}'$  be  $\Sigma_0(M_j^*)$  by the same  $\Sigma_0$  definition. Then, since  $\pi_{\xi i}$  takes  $M_i^*$  cofinally to  $M_i$  by (6), we have

$$A(\zeta) \leftrightarrow \bigvee u \in M_j^* \bigvee x \in \pi_{\xi,i}(u) A'(\zeta, x, p).$$

By the minimality of *i* we know that  $(E_{\nu_j})_{\alpha} \in \underline{\Sigma}_1(M_j^*)$  for  $\alpha < \lambda_j$ . But then:

$$A(\zeta) \leftrightarrow \bigvee u \in m_j^* \{ \gamma < \kappa_j | \overline{A}'(\zeta, x, f(\gamma)) \} \in (E_{\nu_i})_{\alpha}.$$
  
Hence A is  $\underline{\Sigma}_1(M_j^*)$ . QED (7)

Now extend  $I|\xi + 1$  to a potential iteration I' of length  $\xi + 2$  by setting  $\nu'_{\xi} = \overline{\nu}$ , where  $M_j^* = M_{\xi} || \overline{\nu} = \langle J_{\overline{\nu}}^{\overline{E}}, \overline{E}_{\overline{\nu}} \rangle$ . Then  $\kappa_i = \kappa'_{\xi}$  and  $\tau_i = \tau'_{\xi}$ , since  $\pi_{\xi i} \upharpoonright \kappa_j = \text{id}$ . Hence  $h = T(i+1) = T'(\xi+1)$  and  $M_i^* = (M_{\xi}^*)'$ . By the minimal choice of i we conclude

$$\mathbb{P}(\tau_i) \cap \underline{\Sigma}_1(M_i^*) \subset \underline{\Sigma}_1(M_i^*).$$

Hence  $\mathbb{P}(\tau_i) \cap \underline{\Sigma}_1(M_i) \subset \underline{\Sigma}_i(M_i^*)$  by (7). Contradiction! QED (Lemma 3.4.7)

### 3.4.3 Padded iterations

Normal iterations are often used to "compare" two premice M and M'. The comparison iteration or coiteration consists of a pair  $\langle I, I' \rangle$  of iteration I of M and I' of M'. When we have reached  $M_i, M'_i$ , we proceed as follows: We look for the least point of difference — i.e. the least  $\nu$  such that  $M_i || \nu \neq M'_i || \nu$ . Then  $J_{\nu}^{EM_i} = J_{\nu}^{EM'_i}$  and  $E_{\nu}^{M_i} \neq E_{\nu}^{M'_i}$ . Then at least one of  $E_{\nu}^{M_i}, E_{\nu}^{M'_i}$  is an extender. If both are extenders, we continue on the I-side with the index  $\nu_i = \nu$ . However, if, say,  $E_{\nu}^{M_i}$  is an extender and  $E_{\nu}^{M'_i} = \emptyset$ , we iterate by  $\nu_i = \nu$  on the *I*-side and on the *I'*-side do nothing. We then call *i* an *inactive point* on the *I'*-side and set:  $M'_{i+1} = M'_i, \pi'_{i,i+1} = \text{id}$  with i = T'(i+1) in *I*. Thus *i* is active on one or the other side and we have achieved:  $M_{i+1} ||\nu = M'_{i+1}||\nu = \emptyset$ . (This is called "iterating away the least point of difference".) At a limit  $\lambda$  we choose on either side a well founded branch and continue with that.

If all goes well, we eventually reach a point *i* such that  $M_i, M'_i$  or one of  $M_i = M'_i$  is a proper segment of the other.

In order to carry this out we need a slightly more flexible definition of "normal iteration", which admits inactive points. We therefore define:

**Definition 3.4.6.** A padded normal iteration of length  $\mu$  is a sequence:

$$I = \langle \langle M_i | i < \mu \rangle, \langle \nu_i | i \in A \rangle, \langle \pi_{ij} | i \leq_T j \rangle, T \rangle$$

such that:

- (1)  $A \subset \{i : j+1 < \mu\}$  is called the set of *active points* in *I*.
- (2) (a)-(h) of the previous definition hold, where (d)-(f) both require the assumption:  $i \in A$ .
- (3) Let  $h < j < \mu$  such that  $[h, j) \cap A = \emptyset$ . Then:
  - $h \leq_T j, M_k = M_j, \pi_{hj} = \text{id.}$
  - $i \leq h \longrightarrow (i \leq_T h \longleftrightarrow i <_T j)$  for  $i < \mu$ .
  - $j \leq i \longrightarrow (j \leq_T i \longleftrightarrow h <_T i)$  for  $i < \mu$ . (In particular, if  $i+1 < \mu, i \notin A$ , then  $i = T(i+1), M_i = M_{i+1}$ , and  $\pi_{i,i+1} = \text{id.}$ )

**Note.** This gives a new way of potentially extending I of length i + 1. Instead of appointing  $\nu_i$ , we could set:  $i \notin A, M_{i+1} = M_i$ .

All previous results go through a *mutatis mutandis*. We shall often use the term "normal iteration" so as to include padded normal iteration. We then call normal iterations in the sense of our previous definition *strict*. We can turn a padded iteration into a strict iteration simply by omitting the inactive points.

Conversely, we can turn a strict iteration into a padded iteration simply by inserting inactive points. The relevant lemmas are: **Lemma 3.4.8.** Let  $I = \langle \langle M_i \rangle, \langle \nu_i \rangle, \langle \pi_{ij} \rangle, T \rangle$  be a (possibly padded) normal iteration of length  $\mu$ . Let A be the set of active points in I. Set:

$$A' =: \{i : i \in A \lor i + 1 = \mu\}$$

Let  $B \subset \mu$  such that  $A' \subset B$ . Let f be the monotone enumeration of B. Then:

$$I' = \langle \langle M_{f(i)} \rangle, \langle \nu_{f(i)} \rangle, \langle \pi_{f(i), f(j)} \rangle, T' \rangle$$

is a normal iteration , where  $T' = \{\langle i, j \rangle : f(i)Tf(j)\}$ . (Moreover I' is strict if B = A').

**Proof.** (a)-(i) are satisfied by I'.

Conversely:

**Lemma 3.4.9.** Let  $I, \mu$  be as above. Let  $f : \mu \longrightarrow \mu'$  be monotone such that  $\operatorname{lub} f^{"}\mu = \mu'$  if  $\mu$  is a limit ordinal. Set:  $\overline{f}(i) = \operatorname{lub} f^{"}i$  for  $i < \mu$ . For  $i < \mu'$  set:

 $\xi_i = \text{ that } \xi \text{ such that either } \overline{f}(\xi) \leq i \leq f(\xi), \text{ or else } \xi + 1 = \mu \text{ and } f(\xi) < i.$ 

Define:

$$I' = \langle \langle M'_i \rangle, \langle \nu'_i \rangle, \langle \pi'_{ij} \rangle, T' \rangle$$

by:

$$M'_{i} = M_{\xi_{i}}, \pi'_{ij} = \pi_{\xi_{i},\xi_{j}}, T' = \{ \langle i, j \rangle : \xi_{i}T\xi_{j} \}$$

and:

$$\nu_i' = \begin{cases} \nu_{\xi_i} & \text{if } i = f(\xi_i) \\ \text{otherwise undefined} \end{cases}$$

Then I' is a normal iteration.

**Proof:** I' satisfies (a)-(i).

Note. Lemma 3.4.9 enables to recover I form the I' in Lemma 3.4.8.

We leave the proof to the reader.

#### 3.4.4 *n*-iteration

In a normal iteration we always take  $\Sigma^*$  ultrapowers. For technical reasons, however, we may sometimes want to bound the degree of preservation of our ultraproducts. In a 0-*iteration* for instance, we would use the ordinary  $\Sigma_0$ 

ultrapower to pass from  $M_i$  to  $M_{i+1}$ , as long as no  $h \leq_T i+1$  is a truncation point. If, on the other hand, we have reached a truncation point  $h \leq_T i+1$ , we then revert to the full  $\Sigma^*$ -ultrapowers. More generally:

**Definition 3.4.7.** Let  $n \leq \omega$ . By a normal *n*-iteration of *M* of length  $\mu$  we mean

 $\langle \langle M_i | i < \mu \rangle, \langle \nu_i | i + 1 < \mu \rangle, \langle \pi_{ij} | i \leq_T \rangle, T \rangle,$ 

where (a) - (e) and (g), (h) in the definition of "normal iteration" hold, and in addition:

(f) Let h = T(i+1). If  $\tau_i$  is a cardinal in  $M_h$  and  $\pi_{jh}$  is a total map on  $M_j$  for jTh, then  $\pi_{h,i+1} : M_h \to_{E_{\nu_i}}^{(m)} M_{i+1}$ , where  $m \leq n$  is maximal such that  $\kappa_i < \rho_{M_h}^m$ .

Otherwise  $\pi_{h,i+1}: M_i^* \to_{E_{\nu_i}}^{(m)} M_{i+1}$ , where  $M_i^*$  is defined as before and  $m \leq \omega$  is maximal such that  $\kappa_i < \rho_{M_i^*}^m$ .

Note. An  $\omega$ -iteration is then the same as a normal iteration n the sense of our previous definition. We also call such iterations \*-*iterations*, since we then always take the  $\Sigma^*$  ultrapowers. \*-iterations are the ones we are interested in.

It is easily seen that the conclusions of Lemma 3.4.2 hold for normal niterations. Lemma 3.4.3 also holds for these iterations and Lemma 3.4.7holds *mutatis mutandis*. We leave this to the reader. More suprising is:

**Theorem 3.4.10.** Theorem 3.4.4 holds for normal n-iterations.

Before proving this, we again note some consequences. It follows easily that:

**Corollary 3.4.11.** Let I be a normal n-iteration. Let h = T(i+1). Let m be maximal such that  $\kappa_i < \rho_{M_i^*}^m$ . Assume either that  $m \leq n$  or that there is a  $j \leq_T i + 1$  which is a drop point. Then:

$$\pi_{h,i+1}: M_i^* \to_{E_{\nu_i}}^* M_{i+1}.$$

In all other cases we have:

$$\pi_{h,i+1}: M_i^* \to_{E_{\nu_i}}^{(n)} M_{i+1}.$$

But then by induction on i we get:

**Corollary 3.4.12.** Let I be as above. Let  $\pi_{ij}$  be a total map on  $M_i$ . If there is a drop point j such that jTi, then  $\pi_{ij}$  is  $\Sigma^*$ -preserving. Otherwise it is  $\Sigma_0^{(n)}$ -preserving.

As before, we derive Lemma 3.4.10 from:

**Lemma 3.4.13.** Let  $I = \langle \langle M_i \rangle, \langle \nu_i \rangle, \langle \pi_{ij} \rangle, T \rangle$  be a potential *n*-iteration of length i + 2. Then  $\mathbb{P}(\tau_i) \cap \underline{\Sigma}_i(M_i || \nu_i) \subset \underline{\Sigma}_1(M_i^*)$ .

The derivation of Lemma 3.4.10 from Lemma 3.4.13 is exactly as before. We prove Lemma 3.4.13. Almost all steps in the proof of Lemma 3.4.7 go through as before. The only difficulty occurs in the proof of (6), where we derived that  $\pi_{\xi,i}$  is  $\Sigma_0^{(1)}$ -preserving from:  $\kappa_j < \rho_{M_j^*}^1$ . If  $n \ge 1$ , this is unproblematical. Now assume n = 0. If there is a drop point  $l \le_T i$ , then  $\pi_{\xi,i}$  is  $\Sigma^*$ -preserving and there is nothing to prove. Now suppose there is no such drop point.

By the definition of "0-iteration" we then have:  $\pi_{\xi,i}: M_j^* \to_{E_{\nu_j}}^0 M_i$ , which was to be proven.

All other steps in the proof go through. QED (Lemma 3.4.13)

This proves Theorem 3.4.10.

The concept "padded n-iteration" is defined exactly as before. As before, every padded iteration can be converted into a strict iteration by omitting the inactive points, and every strict iteration can be expanded to a padded iteration by inserting inactive points. We leave this to the reader.

## 3.4.5 Copying an iteration

Suppose that I is a normal iteration of a premouse M and  $\sigma: M \to_{\Sigma^*} M'$ , where M' is a premouse. We can attempt to "copy" I onto an iteration I'of M' by repeating the same steps modulo  $\sigma$ . We define:

**Definition 3.4.8.** Let  $I = \langle \langle M_i \rangle, \langle \nu_i \rangle, \langle \pi_{ij} \rangle, T \rangle$  be a strict normal iteration of M. Let  $\sigma : M \to_{\Sigma^*} M$ , where M' is a premouse. We call  $I' = \langle \langle M'_i \rangle, \langle \nu'_i \rangle, \langle \pi'_{ij} \rangle, T' \rangle$  a copy of I induced by  $\langle \sigma, M' \rangle$  with copying map  $\langle \sigma_i | i < lh(I) \rangle$  iff the following hold:

- (a) lh(I') = lh(I) and T' = T
- (b)  $\sigma_i: M_i \to_{\Sigma^*} M'_i$  and  $\sigma_0 = \sigma$
- (c)  $\sigma_i \pi_{li} = \pi'_{li} \sigma_j$  for  $l \leq_T i$
- (d)  $\sigma_i \upharpoonright \lambda_l = \sigma_l \upharpoonright \lambda_l$  for  $l \leq i$

(e)  $\nu'_i = \sigma_i(\nu_i)$  for  $\nu_i \in M_i$ . Otherwise  $\nu'_i = \text{On} \cap M'_i$ .

Note. This definition can easily be extended to padded normal iterations. (b) - (e) are then stipulated for active points, and for inactive points we stipulate:

(f) If *i* is inactive in *I*, it is inactive in *I'* and  $\sigma_{i+1} = \sigma_i$ .

We shall often formulate our definitions and theorems for strict iteration, leaving it to the reader to discover — mutatis mutandis — the correct version for padded iterations. In particular, the remaining theorems in this section will assume strictness.

We also define:

**Definition 3.4.9.**  $\langle I, I', \langle \sigma_i | i < lh(I) \rangle \rangle$  is a *duplication* iff I, I' are normal iterations and I' is a copy of I with copying maps  $\langle \sigma_i \rangle$ .

**Lemma 3.4.14.** Let I' be a copy of I with copying maps  $\langle \sigma_i \rangle$ . Let h = T(i+1).

- (i) If i + 1 is a drop point in I, then it is a drop point in I' and  $M'_i^* = \sigma_h(M_i^*)$ .
- (ii) If i + 1 is not a drop point in I, it is not a drop point in I'. (Hence  $M_i^* = M_h, M'_i^* = M'_h$ .)
- (iii) Let  $F = E_{\nu_i}^{M_i}, F' = E_{\nu'_i}^{M'_i}$ . Then:

$$\langle \sigma_h \upharpoonright M_i^*, \sigma_i \upharpoonright \lambda_i \rangle : \langle M_i^*, F \rangle \to \langle M'_i^*, F' \rangle$$

as defined in §3.2.

(iv) 
$$\sigma_{i+1}(\pi_{h,i+1}(f)(\alpha)) = \pi'_{h,i+1}\sigma_h(f)(\sigma_i(\alpha))$$
 for  $f \in \Gamma^*(\kappa_i, M_i^*)\alpha < \lambda_i$ 

(v) 
$$\sigma_j(\nu_i) = \nu'_i \text{ for } i < j.$$

#### **Proof:**

(i) Let h = T(i+1). Then  $M_i^* = M_h ||\mu$ , where  $\mu \in M_h$  is maximal such that  $\tau_i$  is a cardinal in  $M_h ||\mu$ . But  $\tau'_i = \sigma_i(\tau_i) = \sigma_h(\tau_i)$  by (d), (e). Hence  $\sigma_h(\mu) = \mu'$ , where  $\mu'$  is maximal such that  $\tau'_i$  is a cardinal in  $M'_h$ , and  $\sigma_h(M_h ||\mu) = M'_h ||\mu'$ .

- (ii) If  $\tau$  is a cardinal in  $M_h$ , then  $\tau'_i = \tau_h(\tau)$  is a cardinal in  $M'_h$ , since  $\sigma_h$  is  $\Sigma_1$ -preserving.
- (iii) Clearly  $\sigma_h \upharpoonright M_i^* : M_i^* \to_{\Sigma^*} M_i'^*$  by (i) and (ii). Now let  $x \in \mathbb{P}(\kappa_i) \cap M_i^*$ and  $\alpha_1, \ldots, \alpha_n < \lambda_0$ . Since  $\sigma_i : M_i \longrightarrow M_i'$  is  $\Sigma^*$ -preserving we have:

$$\langle \vec{\alpha} \rangle \in F(x) \leftrightarrow \langle \sigma_i(\vec{\alpha}) \rangle \in F'(\sigma_i(x)).$$

But  $\sigma_i(x) = \sigma_h(x)$ , since by (d) we have:  $\sigma_i \upharpoonright J_{\lambda_n}^{E^{M_i}} = \sigma_h \upharpoonright J_{\lambda_h}^{E^{M_h}}$ .

(iv) If  $f \in M_i^*$ , then by (c):

$$\sigma_{i+1}\pi_{h,i+1}(f) = \pi'_{h,i+1}\sigma_h(f).$$

Otherwise  $f(\xi) \simeq G(\xi, q)$  where  $q \in M_i^*$  and G is a good  $\Sigma_1^{(n)}(M_i^*)$  function for an n such that  $\kappa_i < \rho_{M_i^*}^{n+1}$ . But then:

$$\sigma_{i+1}\pi_{h,i+1}(f)(\xi) \simeq G'(\xi,\sigma_{i+1}\pi_{h,i+1}(q))$$
$$\simeq G'(\xi,\pi'_{h,i+1}\sigma_h(q))$$
$$\simeq \pi'_{h,i+1}\sigma_h(f)$$

where G' is  $\Sigma_1^{(n)}(M_i'^*)$  by the same good definition.

(v) If j > i + 1, then  $\nu_i < \lambda_{i+1}$  and  $\sigma_j(\nu_i) = \sigma_{i+1}(\nu_i)$ . But letting h = T(i+1), we have:

$$\sigma_{i+1}(\nu_i) = \sigma_{i+1}\pi_{h,i+1}(\tau_i) = \pi'_{h,i+1}\sigma_h(\tau_i),$$

where

$$\sigma_h(\tau_i) = \sigma_i(\tau_i) = \tau'_i$$
, since  $\tau_i < \lambda_h$ .

Hence  $\sigma_{i+1}(\nu_i) = \pi'_{h,i+1}(\tau'_h) = \nu'_i$ .

QED (Lemma 3.4.14)

It is apparent from Lemma 3.4.14 that there is only one way to extend a copy of I|i+1 to a copy of I|i+2. Moreover, the copying map  $\sigma_i$  is unique. Similarly, if  $\eta$  is a limit ordinal and I' is a copy of  $I|\mu$  with copying maps  $\langle \sigma_i | i < \eta \rangle$ , ther is only one way to extend I' to a copy of  $I|\eta+1$ , for then:

$$M', \langle \pi'_{i,\eta} | iT\eta \rangle$$

is the direct limit of:

$$\langle M'_i | i < \eta \rangle, \langle \pi'_{ij} | i \leq_T j <_T \eta \rangle,$$

and  $\sigma_{\eta}$  is defined by:

$$\sigma_{\eta}\pi_{i\eta} = \pi'_{i\eta}\sigma_i$$
 for  $i <_T \eta$ .

Hence, by induction on lh(I) we get:

**Lemma 3.4.15.** Let I be a normal iteration of M. Let  $\sigma : M \to_{\Sigma^*} M'$ . Then there is at most one copy I' of I induced by  $\sigma$ . Moreover, the copying maps  $\sigma_i$  are unique.

Now suppose that I is a normal iteration of length i + 1 and I' is a copy of I with copying maps  $\langle \sigma_h | h \leq i \rangle$ . Extend I to a potential iteration  $\tilde{I}$  of length i + 2 by appointing  $\nu_i$ . Extend I' to a potential iteration  $\tilde{I}'$  by appointing:

$$\nu_i' = \begin{cases} \sigma_i(\nu_i) \text{ if } \nu_i \in M_i \\ \operatorname{On} \cap M_i' \text{ if } \nu_i = \operatorname{On} \cap M_i \end{cases}$$

We call  $\langle \tilde{I}, \tilde{I}', \langle \sigma_j || \leq i \rangle \rangle$  a *potential duplication* of length i + 2. The formal definition is:

**Definition 3.4.10.** Let I, I' be potential iteration of length i+2.  $\langle \tilde{I}, \tilde{I}', \langle \sigma_j | j \leq i \rangle$  is a *potential duplication* of length i+2 iff

- $\langle \overline{I}, \overline{I}', \langle \sigma_j | j \leq i \rangle \rangle$  is a duplication of length i+1, where  $\overline{I} = \widetilde{I} | i+1, \overline{I}' = I' | i+1$ .
- $\sigma_i(\nu_i) = \nu'_i$  if  $\nu_i \in M_i$ . Otherwise  $\nu'_i = \operatorname{On} \wedge M'_i$ .

**Note.** It is then easily seen that T(i+1) = T'(i+1). We also know that  $E_{\nu_i}^{M_i}$  is close to  $M_i^n$  and  $E_{\nu'_i}^{M'_i}$  is clost to  $M'_i$ . The following theorem is an analogue of theorem 3.4.7

**Lemma 3.4.16.** Let  $\langle I, I', \langle \sigma_i \rangle \rangle$  be a potential duplication of length i + 2. Let h = T(i + 1). Then:

$$\langle \sigma_h \upharpoonright M_i^*, \sigma_i \upharpoonright \lambda_i \rangle : \langle M_i^*, F \rangle \to^* \langle {M'}_i^*, F' \rangle$$

(as defined in §3.2) where  $F = E_{\nu_i}^{M_i}, F' = E_{\nu'_i}^{M'_i}$ .

Before proving the theorem, we note some of its consequences. It gives us exact criteria for determining whether the copying process can be continued one step further.

**Lemma 3.4.17.** Let I be a normal iteration of M of length i + 2. Let  $\sigma: M \to M'$  induce a copy I' of I|0+1 with copying maps  $\langle \sigma_i | j \leq i \rangle$ . Set:

$$\nu_i' = \begin{cases} \sigma_i(\nu_i) \text{ if } \nu_i \in M_i \\ \operatorname{On} \cap M_0' \text{ if } \nu_i = \operatorname{On} \cap M_i \end{cases}$$

Then  $\sigma$  induces a copy of I iff  $M'_i^*$  is  $\Sigma^*$ -extendible by  $E_{\nu'_i}^{M'_i}$ .

**Proof:** If  $M'_i^*$  is not extendible, then no such copy can exist. Now let  $M'_i^*$  be extendible. Let  $\pi'_{h,i+1}: M'_i^* \to_{E'_{n'}}^* M'_{i+1}^*$ . By theorem 3.4.16 and

Lemma 3.2.23 it follows that there is a unique  $\sigma: M_{i+1} \to_{\Sigma^*} M'_{i+1}$  such that  $\sigma \pi_{h,i+1} = \pi'_{h,i+1} \cdot \langle \sigma_h \upharpoonright M^*_i \rangle$ , where h = T(i+1). Set:  $\sigma_{i+1} =: \sigma$ . This gives us the copy I'' of I with copying maps  $\langle \sigma_j | j \leq 0 + 1 \rangle$ .

QED (Lemma 3.4.17)

We also have:

**Lemma 3.4.18.** Let I be a normal iteration of M of length  $\eta + 1$ , where  $\eta$  is a limit ordinal. Let  $\sigma : M \to_{E^*} M'$  induce a copy I' of  $I|\eta$ . We can extend I' to a copy of I induced by  $\sigma$  iff  $b = T''\{\eta\}$  is a well founded branch in I'.

The proof is left to the reader.

We also note:

**Lemma 3.4.19.** Let I be a normal iteration of limit length. Let I' be a copy of I. If b is a cofinal well founded branch in I', then it is a cofinal well founded branch in I.

The proof is left to the reader.

We now turn to the proof of theorem 3.4.16. As with theorem 3.4.7 we derive it from an even stronger lemma:

**Lemma 3.4.20.** Let  $\langle I, I', \langle \sigma_i \rangle \rangle$  be a potential duplication of length i + 2. Let  $A \subset \tau_i$  be  $\Sigma_1(M_i || \nu_i)$  in a parameter p. Let  $A' \subset \tau'_i$  be  $\Sigma_1(M_i || \nu_i)$  in  $\sigma_i(p)$  by the same definition. Then A is  $\Sigma_1(M_i^*)$  in a parameter q and A' is  $\Sigma_1(M_i^*)$  in  $\sigma_h(q)$  by the same definition, where h = T(i+1).

The derivation of theorem 3.4.16 from lemma 3.4.20 is a virtual repetition of the proof of theorem 3.4.4 from lemma 3.4.7. We leave it to the reader.

Lemma 3.4.20 is proven by a virtual repetition of the proof of lemma 3.4.7, making changes as necessary. We give a brief sketch of the proof:

Suppose not. Let  $I, I', \nu_i, \nu'_i$  be counterexamples of length i + 1, where i is chosen minimally. Let h = T(i + 1) = T'(i + 1). Then:

(1) h < i. Suppose not. Then  $M_i || \nu_i \subset M_i^*$  and  $M'_i || \nu'_i \subset M'_i^*$  as before. If  $\nu_i \in M_i^*$ , then  $\sigma_i(M || \nu_i) = M'_i || \nu'_i$ . Hence  $A \in M_i^*$  and  $\sigma_i(A) = A'$ . Contradiction!

- (2)  $\nu_i = \operatorname{On}_{M_i}$  and  $\rho_{M_i}^i \leq \tau_i$ . Otherwise, as before  $A \in \mathbb{P}(\tau_i) \cap M_i^*, A' \in \mathbb{P}(\tau_i) \cap M'_i^*$  and  $\sigma_h(A) = \sigma_i(A) = A'$ . Contradiction!
- (3) i is not a limit cardinal. The proof of this is a virtual repetition of the argument given in the proof of lemma 3.4.7. We leave it to the reader.

Now let  $i = j + 1, \xi = T(i)$ . Exactly as before we have:

(4) 
$$M_j^* = \langle J_\nu^E, E_\nu \rangle, {M'}_j^* = \langle J_{\nu'}^{E'}, E'_{\nu'} \rangle$$
 where  $E_\nu, E'_\nu \neq \emptyset$ .

- (5)  $\tau_i < \kappa_j$ .
- (6)  $\pi_{\xi,i}: M_j^* \to_{E_{\nu_j}} M_i$  is a  $\Sigma_0$  ultrapower (and therefore cofinal). Similarly for  $\pi'_{\xi,i}: M'_j^* \to_{E_{\nu'_j}} M'_i$ . By the minimality of  $\sigma$  we know that for all  $\alpha < \lambda_j$ ,  $(E^{M_j}_j)_{\alpha}$  is  $\Sigma_1(M_j^*)$  in a parameter r and  $(E^{M_i}_{\nu'_i})_{\sigma_i(\alpha)}$  is  $\Sigma_1(M'_j^*)$  in  $\sigma_{\xi}(r)$  by the same definition. Using this we can repeat the argument in the proof of Lemma 3.4.7 to get:
- (7) A is  $\Sigma_1(M_i^*)$  in a q and A' is  $\Sigma_1(M'_i^*)$  in  $\sigma_{\xi}(q)$  by the same definition.

Now extend  $I|\xi + 1$  to a potential iteration  $\tilde{I}$  of length  $\xi + 2$  by setting  $\tilde{\nu}_{\xi} = \nu$ , where  $\nu$  is as in (4). Extend  $I'|\xi + 1$  to  $\tilde{I}'$  by setting  $\tilde{\nu}_{\xi} = \nu'$  where  $\nu'$  is as in (4). Then  $\kappa_i = \tilde{\kappa}_{\xi}, \tau_i = \tilde{\tau}_{\xi}, \kappa'_i = \tilde{\kappa}_{\xi}, \tau'_i = \tilde{\tau}'_{\xi}$  as before. Hence  $h = \tilde{T}(\xi + 1) = \tilde{T}'(\xi + 1)$  and  $M_i^* = \tilde{M}_{\xi}^*, M'_i^* = \tilde{M}'_{\xi}^*$ . By this minimality of *i* we conclude that *A* is  $\Sigma_1(M_i^*)$  ia a *q* and *A'* is  $\Sigma_1(M'_i^*)$  in  $\sigma_h(q)$  by the same definition. Contradiction! QED (Lemma 3.4.20)

#### **3.4.6** Copying an *n*-iteration

**Definition 3.4.11.** Let  $I = \langle \langle M_i \rangle, \langle \nu_i \rangle, \langle \pi_{ij} \rangle, T \rangle$  be a normal *n*-iteration  $(n \leq \omega)$ . Let  $\sigma : M \to_{\Sigma^M}, M'$ . We call:

$$I' = \langle \langle M'_i \rangle, \langle \nu'_i \rangle, \langle \pi'_{ij} \rangle, T' \rangle$$

a copy (or n-copy) of I induced by  $\langle \sigma, M' \rangle$  iff I' is an n-iteration satisfying (a), (c), (d), (e) of the previous definition together with

(b')  $\sigma_0 = \sigma$  and  $\sigma : M_i \to_{\Sigma_1^{(n)}} M'_i$ . Moreover, if some  $h \leq_T i$  is a truncation point, then  $\sigma_i$  is  $\Sigma^*$ -preserving.

The notion "*n*-duplication" and "potential *n*-duplication" are defined as before. Lemma 3.4.14 goes through as before exept (iv) must be reformulated as:

(iv) If no  $l \leq_T i + 1$  is a truncation point and  $\kappa_i < \rho_{M_b}^n$ , then:

$$\sigma_{i+1}(\pi_{h,i+1}(f))(\alpha) = \pi'_{h,i+1}\sigma_i(f)(\sigma_i(\alpha))$$

for  $f \in \Gamma^n_*(\kappa_i, M_h), \alpha < \lambda_i$ . In all other cases the equation holds for

$$f \in \Gamma^*(\kappa_i, M_i^*), \alpha < \lambda_i.$$

Lemma 3.4.15 then holds as before. Theorem 3.4.16 and lemma 3.4.17 - 3.4.19 then go through as before. By theorem 3.4.16 we also get:

**Lemma 3.4.21.** Let  $\langle I, I', \langle \sigma_i \rangle \rangle$  be an *n*-duplication. Let  $i <_T j$  in I such that  $\pi_{ij}$  is total on  $M_i$ .

- (a) If no  $l \leq_T i$  is a truncation point and  $\kappa_i < \rho_{M_i}^n$ , then  $\pi_{ij} : M_i \to_{\Sigma_1^{(n)}} M_i$ .
- (b) In all other cases  $\pi_{ij}$  is  $\Sigma^*$ -preserving.

These lemmas and theorems hold *mutatis mutandis* for padded n-iterations. The details are left to the reader.

# 3.5 Iterability

A mouse is a premouse which is iterable. Iterability is, however, as complex a notion as that of iterating itself. We begin with *normal iterability* which says that any normal iteration of M constructed according to an appropriate strategy, can be continued.

### 3.5.1 Normal iterability

**Definition 3.5.1.** A premouse M has the normal uniqueness property (NUP) iff every normal iteration of M of limit length has at most one cofinal well founded branch. The simplest mice, such as  $0^{\#}, 0^{\#\#}$  etc. are easily seen to have this property. Unfortunately, however, there are mice which do not. If a premouse M does satisfy NUP, then normal iterability can be defined by: