

### §3. The fine structure of the constructible hierarchy

Let  $u$  be a transitive set which is closed under the formation of finite sets (thus, in particular,  $u$  is closed under  $n$ -tuples and  $\text{Fm } u \subset u$ ). In this section, we attempt to determine the p.r. closed levels  $\alpha L_\alpha[u]$  of the constructible hierarchy over  $u$  at which interesting things happen — f.r. when is  $L_\alpha[u]$  admissible? It turns out ~~not~~ that this is the case iff for no  $\beta < \alpha$  ~~there is a map~~  $\Delta_1$ , map of  $L_\beta[u]$  onto  $L_\alpha[u]$  exists.

~~As a corollary we get if  $L_\alpha[u]$  is admissible iff  $\Delta_1$  applies.~~

(Corollary: If  $\alpha$  is a singular cardinal and  $\bar{\alpha} < \alpha$ , then  $L_\alpha[u]$  is admissible.)

(2)

Throughout this section,  $u$  will be a fixed transitive set which is closed under finite subsets.  $L_d[u]$  will always be p.r. closed. 'p.r.' will always mean 'p.r. in parameters from  $u \cup \{u\}$ '. We begin with an observation on p.r. functions:

Lemma 1 There is a p.r. function  $\gamma(x, y)$  which maps  $u \times d$  onto  $L_d$ .  
proof.

There is a p.r. map  $l: \text{On}^2 \longleftrightarrow \text{On}$  s.t.  $\beta, \gamma \leq \langle \beta, \gamma \rangle$  (hence the inverses  $l^{-1}, r$ , defined by  $d = \langle l(d), r(d) \rangle$  are p.r.). To see this, we order  $\text{On}^2$  by:

$$\langle \alpha, \beta \rangle R \langle \gamma, \delta \rangle \iff \max(\alpha, \beta) < \max(\gamma, \delta) \vee$$

v.  $\max = \max \wedge \alpha < \gamma \wedge \beta < \delta$ .

ii.  $\max = \max \wedge \alpha = \gamma \wedge \beta < \delta$ .

(3)

Let  $\langle \rangle : R \longleftrightarrow \in^{\uparrow} \Omega_n$ . To see that  $\langle \rangle$  is p.r., we first define the function  $\langle \langle 0, \beta \rangle \mid \beta \in \Omega_n \rangle$  by :

$$\langle 0, \beta \rangle = \sup_{\nu < \beta} (\langle 0, \nu \rangle + \nu \cdot 2) ,$$

and then set :

$$\langle \nu, \beta \rangle = \langle 0, \beta \rangle + \nu \text{ if } \nu < \beta$$

$$\langle \beta, \nu \rangle = \langle 0, \beta \rangle + \beta + \nu \text{ if } \nu \leq \beta .$$

We can represent  $n$ -tuples of ordinals by :

$$\langle \beta_1, \dots, \beta_n \rangle = \langle \beta_1, \langle \beta_2, \dots, \beta_n \rangle \rangle .$$

Define a function  $h(z, y)$  by :

$$h(\langle \varphi, n, i \rangle, y) =$$

$$= \varphi \left( \overline{\nu_{i_1}, \dots, \nu_{i_m}} / \underline{L_{\gamma_1}[u]}, \dots, \underline{L_{\gamma_n}[u]} \right)$$

if  $\varphi \in \text{Fml}_u^{\Sigma_0}$ ,  $n < \omega$ ,  $i < \omega$ ,

$i = \langle i_1, \dots, i_m \rangle$ ,  $\nu = \langle \gamma_1, \dots, \gamma_n \rangle$ .

$$h(z, y) = 0 \text{ otherwise} .$$

(4)

Then  $h$  is a p.r. function which maps  $u \times d$  onto the set of  $\Sigma$ .

formulae  $\varphi$  containing only the constants  $\underline{x}$  ( $x \in u$ ),  $\underline{L_v[u]}$  ( $v \in d$ ).

We have seen that every  $y \in L_d[u]$  has the form:

$$y = \left\{ x \in L_\beta[u] \mid \models \varphi(v_0/x) \right\},$$

where  $\beta < d$  and  $\varphi$  is such a formula. Hence, we may define the desired function  $\gamma$  by:

$$\gamma(x, \{\beta, s\}) = \left\{ z \in L_\beta[u] \mid \models h(x, s)(v_0/z) \right\}$$

QED.

(5)

Def The function  $r(\vec{x})$  uniformises the relation  $Ry\vec{x}$  iff  $\text{dom}(r) = \text{dom}(R)$  and  $\wedge\vec{x} (\forall y Ry\vec{x} \longleftrightarrow R_{r(\vec{x})}\vec{x})$ .

A structure  $M = \langle |M|, \epsilon, A_1, \dots, A_n \rangle$  is called  $\Sigma_n$ -uniformisable ( $n \geq 1$ ) iff each  $\Sigma_n$  relation  $R$  s.t.  $\text{rng}(R) \subset \Omega^n$  is uniformisable by a  $\Sigma_n$  function.

(This ~~—~~ notion should really be called 'ordinal uniformability'. We use it in preference to the stronger notion because the latter may fail for lack of a nice well ordering of  $|M|$ ).

(6)

Thm 1.  $L_d^{[u]}$  is  $\Sigma_1$  uniformisable.

proof. Let  $R(v, \vec{x})$  be a  $\Sigma_1$  relation s.t.  $\text{rng}(R) \subset \text{On}$ . Let

$$Rv\vec{x} \iff \forall y P_{v,y}\vec{x},$$

where  $P$  is  $\Sigma_0$ . Set:

$$Qv\tau\vec{x} \iff \bigvee_{y \in u} P_{v,y}(\tau, y)\vec{x}.$$

Then  $Q$  is p.r. (in the parameters entering the  $\Sigma_1$  definition of  $R$ ) and:

$$Rv\vec{x} \iff \forall \tau Qv\tau\vec{x}.$$

Set:  $g(\vec{x}) \simeq_u \tau Q\ell(\tau) r(\tau)\vec{x}$

(where  $\langle \ell(\tau), r(\tau) \rangle = \tau$ ). Then  $g$  is  $\Sigma_1$ , since:

$$\begin{aligned} \tau = g(\vec{x}) &\iff Q\ell(\tau) r(\tau)\vec{x} \wedge \\ &\quad \wedge \delta < \tau \wedge Q\ell(\delta) r(\delta)\vec{x}. \end{aligned}$$

Hence,  $r(\vec{x}) \simeq \ell g(\vec{x})$  is a  $\Sigma_1$  function which uniformises  $R$ . QED

(Note : By the proof of ~~Lemmas~~<sup>Thm 1</sup>, we may take  $\alpha$  as being  $\Sigma_1$  in the same parameters which enter the  $\Sigma_1$  definition of  $R_\alpha$ ).

Def  $X \prec_{\sum_m} L_\alpha[u]$  (" $X$  is a  $\Sigma_m$ -elementary submodel of  $L_\alpha[u]$ )  $\iff_{pt}$   
 $\iff_{pt} u \cup \{u\} \subset X$  and for every ~~constant~~  
 $\varphi \in \text{Fml}_X^{\Sigma_m}$  :

$$\models_{\langle X, \in \rangle} \varphi \iff \models_{\langle L_\alpha[u], \in \rangle} \varphi.$$

In other words,  $X \prec_{\sum_m} L_\alpha[u]$  iff  
 $u \cup \{u\} \subset X$  and for each  $R \subset L_\alpha^{[u]}{}^{m+1}$   
which is  $\Sigma_1$  in parameters from  $X$ :

$$\forall y R_y \vec{x} \iff \forall y \in X R_y \vec{x} \quad \text{for } \vec{x} \in X.$$

In particular, if  $X \subseteq \sum_{\alpha_1} L_\alpha[u]$ , then

$X$  is closed under  $\Sigma_1$  functions definable with parameters from  $X$ . Since p.r. functions are  $\Sigma_1$  in the parameter  $u$ ,  $X$  is p.r. closed.

Clearly,  $\in^r X$  satisfies the axiom of extensionality, since, if  $x, y \in X$  and  $x \neq y$ , then  $\forall z (z \in x \leftrightarrow z \in y)$ , hence  $\forall z \in X (z \in x \leftrightarrow z \in y)$ . Therefore there exists a map  $\pi$  of  $X$  onto a transitive set  $v$  s.t.

$$\pi : \langle X, \in \rangle \xrightarrow{\sim} \langle v, \in \rangle,$$

If  $\varphi$  is a  $\Sigma_1$  formula (without constants), then

$$(+)\quad \models_{L_\alpha[u]} \varphi(\vec{x}) \longleftrightarrow \models_v \varphi(\vec{\pi(x)}) \text{ for } \vec{x} \in X.$$

We may conclude:

$$(++) \pi f(\vec{x}) = f(\pi(\vec{x})) \text{ for all p.r. } f,$$

since, by the stability lemma, each p.r.  $f$  has a  $\Sigma_1$  definition which is absolute with respect to p.r. closed domains; i.e. there is a  $\Sigma_0$  formula  $\varphi_f$  (with ~~one~~ constants  $\underline{x}$  ( $x \in u \cup \{u\} \cup \{\omega\}$ )) s.t.

$$y = f(\vec{x}) \iff \forall z \models \varphi_f(z, \underline{y}, \vec{x}) \text{ for all } y, \vec{x}$$

$$y = f(\vec{x}) \iff \forall z \in L_d[u] \models \varphi_f(z, \underline{y}, \vec{x})$$

$$\text{for } \vec{x} \in L_d[u].$$

Hence, if  $\vec{x} \in X$ , we have:

$$y = f(\vec{x}) \iff \forall z \models \varphi_f(z, \underline{y}, \vec{x}) \\ \in L_d[u]$$

$$\iff \forall z \in u \models \varphi_f(z, \underline{\pi(y)}, \underline{\pi(\vec{x})})$$

$$\rightarrow \pi(y) = f(\pi(\vec{x})). \quad \text{QED}(++)$$

(Note) (++) implies that  $\sigma$  is p.r. closed

By (++) , we get :

Lemma 2 If  $X \preceq_{\Sigma_1} L_\alpha[u]$ ,  $\pi : \langle X, \in \rangle \rightleftarrows \langle \text{U}, \in \rangle$

and  $\text{U} \subset \text{U}'$ , then  $\text{U}_\beta \leq \alpha \Rightarrow \text{U} = L_\beta$ .

Proof.

Since  $\langle L_\nu[u] \mid \nu \in \text{On} \rangle$  is p.r., we have

by (++) :  $\pi L_\nu[u] = L_{\pi(\nu)}[u]$ .

For all  $x \in X$ , we have :

$\forall \nu \in L_\alpha[u] \quad x \in L_\nu[u]$ , hence

$\forall \nu \in X \quad x \in L_\nu[u]$ , hence

$\forall \nu \in X \quad \pi(x) \in L_{\pi(\nu)}[u]$ .

Let  $\beta$  be the least ordinal not in  $\text{U}$ .

Then  $\beta = \pi``\text{On} \cap X$ . Hence :

$$\text{U} = \pi``X = \bigcup_{\nu \in X} L_{\pi(\nu)}[u] = L_\beta[u]$$

QED

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Using the fact that  $L_\beta[u] \subset L_\alpha[u]$ , we can strengthen (++) to:

Lemma 3 Let  $X, \pi$  be as in Lemma 2 and let  $f$  be a function which is  $\Sigma_1$  without parameters (or at most parameters  $x \in X$  s.t.  $\pi(x) = x$ ). Then, whenever  $\vec{x} \in X$  and  $f(\vec{x})$  is defined, so is  $f(\pi(\vec{x}))$  and

$$\pi f(\vec{x}) = f(\pi(\vec{x})).$$

Proof:

Let  $\varphi$  be a  $\Sigma_1$  formula defining  $f$  in  $L_\alpha[u]$  (containing at most constant  $x$  s.t.  $\pi(x) = x$ ). Then

$$y = f(\vec{x}) \longleftrightarrow \models_{L_\alpha[u]} \varphi(y, \vec{x})$$

$$\longleftrightarrow \models_{L_\beta[u]} \varphi(\pi(y), \pi(\vec{x}))$$

$$\rightarrow \models_{L_\alpha[u]} \varphi(\pi(y), \pi(\vec{x}))$$

$$\rightarrow \pi(y) = f(\pi(\vec{x})).$$

QED

Lemma 4 There is a  $\Sigma_1$  function  $h$

s.t.  $\text{dom}(h) \subset u \times L_d[u]$  and

$$\forall x \in L_d[u] \quad x \in h''(u \times \{\bar{x}\}) \subseteq_{\Sigma_1} L_d[u].$$

Proof.

Define  $s(z, x)$  by:

$$s(\langle \varphi, i, j \rangle, x) = \varphi(v_i v_j / u, z)$$

if  $\varphi \in \text{Fml}_u^{\Sigma_1}, i, j < \omega, i \neq j$ .

$$s(z, x) = \underline{0} \text{ otherwise.}$$

Then  $s$  is p.r. and maps  $u$  onto  $\text{Fml}_u^{\Sigma_1}$ .

Note that, since  $\models^{\Sigma_0}$  is p.r.,  $\models^{\Sigma_1}_{L_d[u]}$

is  $\Sigma_1$ . Set:

$$R_{\varphi} \leftrightarrow_p \models^{\Sigma_1} \varphi(v_0 / z).$$

and let  $\sigma$  uniformise  $R$ . Let

$\sigma$  be  $\Sigma_1$  in the parameter  $p$  and

set:  ~~$h(\langle z, w \rangle, x) = \sigma(\langle z, w \rangle, \langle x, p \rangle)$~~

$$h(\langle z, w \rangle, x) = \sigma(z, \sigma(w, s(w, \langle x, p \rangle)))$$

if  $\langle z, w \rangle \in u$ ; (otherwise undefined).

Clearly,  $h \in \Sigma_1$ .

Let  $x \in L_d[u]$ . Set  $X = h''(u \times \{x\})$ .

Claim  $X \prec_{\Sigma_1} L_d[u]$

Let  $A \subset L_d[u]$  be  $\Sigma_1$  in parameters from  $X$ . We must show:

$$\forall y A_y \leftrightarrow \forall y \in X A_y.$$

Let  $\vec{z} \in X$  be the parameters of  $A$ ; since  $z_i = h(w_i, x)$  ( $w_i \in u$ ),  $A$  is  $\Sigma_1$  in parameter from  $u \cup \{u, x, p\}$ .

Assume  $\forall y A_y$ . Then  $y = \gamma(z, v)$  for some  $z \in u$ ,  $v \in d$ . The set

$A' = \{v \mid A_{\gamma(z, v)}\}$  is  $\Sigma_1$  in parameters from  $u \cup \{u, x, p\}$ . Hence there is a  $\varphi$  with constants from  $u \cup \{u, x, p\}$  s.t.

$$A' \vdash \models^{\Sigma_1} \varphi(\underline{z}).$$

Set  $y' = \gamma(z, \tau(\varphi))$ . Then  $A_{y'}$ .

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But  $y \in X$ , since, letting  $\varphi = \sigma(\omega, \langle x, p \rangle)$ ,

$$y = \gamma(z, \sigma(\sigma(\omega, \langle x, p \rangle))) = h(\langle z, \omega \rangle, x)$$

~~QED~~ QED

Thm 2 The following conditions are equivalent:

- (i) There is a  $\Sigma_1$  acu s.t.  $a \notin L_d[u]$
- (ii) There is a  $\Sigma_1$  map from a subset of  $u$  onto  $L_d[u]$ .

proof.

(iii)  $\rightarrow$  (i) is trivial, since  $a = \{x \mid x \notin f(x)\}$  is  $\Sigma_1$  but not an element of  $L_d[u]$ , for if not, we should have:

$$x \in a \longleftrightarrow x \in f(z) \quad \text{---} f$$

for some  $z$ ; hence:

$$\star z \in a \quad \text{---} \longleftrightarrow z \notin a.$$

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(ii)  $\rightarrow$  (iii). Let  $a \in u$  be  $\Sigma_1$ ,  $a \notin L_\alpha[u]$ , let  $a$  be  $\Sigma_1$  in  $x$ . Set:  $x = h''(u \times \{x\})$   
 let  $\pi : \langle x, \in \rangle \leftrightarrow \langle L_\beta[u], \in \rangle$ . Then,  
 if  $\varphi(z, x)$  is the  $\Sigma_1$  definition of  $a$ ,  
 we have:

$$\begin{aligned} z \in a &\longleftrightarrow \models_{L_\alpha[u]} \varphi(z, x) \\ &\longleftrightarrow \models_{L_\beta[u]} \varphi(z, \pi(x)). \end{aligned}$$

Hence  $a$  is  $\Sigma_1$  in  $L_\beta[u]$ . But this means  
 that  $\beta = \alpha$ , since otherwise  
 $a \in L_{\beta+1}[u] \subset L_\alpha[u]$ . Let  $h$  be  $\Sigma_1$ ,  
 in the parameter  $p$ ; in particular let:

~~#~~  $y = h(z, x) \longleftrightarrow H(p, y, z, x),$

where  $H$  is  $\Sigma_1$  without parameters.

~~Set:  $h'(z, x) = \pi^{-1} h(z, \pi^{-1}(x))$~~  By (i),  
 ~~$\pi^{-1} h(z, \pi^{-1}(x))$~~ .

~~we get:~~

~~$\pi(y) = h'(z, \pi(x)) \longleftrightarrow H(p, \pi(p), \pi(y), \pi(x)) \in$~~

Set:  $h'(z, x) \simeq \pi h(z, \pi^{-1}(x))$ . By (+);  
 $\pi(y) = h'(z, \pi(x)) \leftrightarrow H(\pi(p), \pi(y), z, \pi(x))$ .

Thus,  $h'$  is  $\Sigma_1$  in  $\pi(p)$ , and

$$h'' u \times \{\pi(x)\} = \pi'' X = L_d[u].$$

Set  $f(z) \simeq h'(z, \pi(x))$ . Then  $\text{dom}(f) \subset u$ ,  
 $f$  is  $\Sigma_1$  and  $f'' u = L_d[u]$ . QED

As a corollary of Thm 2, we obtain:

Thm 3 The following conditions are equivalent:

- (a) There is a  $\Delta_1$  set  $a \subset u$  s.t.  $a \notin L_d[u]$
- (b) There is a  $\Delta_1$  map of  $u$  onto  $L_d[u]$ .

proof:

(b)  $\rightarrow$  (a) follows as before

We now prove (a)  $\rightarrow$  (b)

(1)

By Thm 2, there exists a  $\Sigma_1$  map  $f'$  s.t.  $\text{dom}(f') \subset u$  and  $f'''u = L_\alpha[u]$ . We must replace  $f'$  by a  $\Sigma_1$  map which is defined on the whole of  $u$ . Since  $A$  is  $\Delta_1$ , we have:

$$z \in a \iff \forall y A_0 y z$$

$$z \notin a \iff \forall y \neg A_1 y z,$$

where  $A_0, A_1$  are  $\Sigma_0$ . In particular,

$$\wedge_{z \in u} \forall y (A_0 y z \vee A_1 y z).$$

~~Set:  $G y z \iff A_0 y z \vee A_1 y z$ .~~

~~Let  $g$~~

Set:  $G y z \iff \forall y \in L_\gamma[u] (A_0 y z \vee A_1 y z)$ .

Let  $g$  uniformise  $\circ G$ . Then

$g''u$  is unbounded in  $L_\alpha[u]$ , since if  $g''u \subset L_\gamma[u]$ , ~~for some~~,  $\gamma < \alpha$ , we would have:  $a \in L_{\gamma+1}[u] \subset L_\alpha[u]$ .

Since  $f'$  is  $\Sigma_1$ , we have:

$$y = f'(x) \iff \forall z F z \rightarrow y x,$$

where  $F$  is  $\Sigma_0$ . Set:

$$\tilde{f}(w, x) = \begin{cases} y & \text{if } \forall z \in L_y[u] F z \rightarrow y x \\ & \text{and } y \in L_y[u] \\ 0 & \text{if not.} \end{cases}$$

Then  $\tilde{f}$  is p.r.

Set:  $f(\langle z, w \rangle) = \tilde{f}(g(z), w)$  if  $\langle z, w \rangle \in u$   
 $f(x) = \emptyset$  otherwise.

Then  $f''u = f'''u = L_d[u]$ . QED

## Non projectible admissible sets

Def Call  $M = \langle |M|; \in, A_1, \dots, A_n \rangle$

non projectible iff  $M$  is admissible and satisfies the stronger replacement axiom:

$$\Lambda u \vee v \Lambda x \in u (\forall y \varphi \longleftrightarrow \forall y \in v \varphi)$$

where  $\varphi$  is  $\Sigma_0$ .

One easily establishes the following

lemma Let  $M$  be admissible; then the following are equivalent

(a)  $M$  is non projectible

(b)  $x \in M \rightarrow x \cap A \in M$  for every

$\Sigma_1$  set  $A$ .

(c)  $x \in M \rightarrow f''x \in M$  for every

$\Sigma_1$  map  $f$ .

We wish to characterise the  $\alpha$  s.t.

$L_\alpha[u]$  is non projectible. Our major tool in this endeavour will be:

Lemma 6 Let  $h$  be as in Lemma 4.

Let  $h$  be  $\Sigma_1$  in the parameter  $x$ .

~~Let  $\alpha \in L_\alpha[u]$  be transitive~~

Let  $v \subset L_\alpha[u]$  be transitive, closed under finite sets, and let

$u \cup \{x\} \subset v$ . Then

$$\forall \beta \leq \alpha \quad h''u \times v = L_\beta[u].$$

Proof. Let  $X = h''u \times v$ .

Obviously,  $X \prec_{\Sigma_1} L_\alpha[u]$ . Let

$\pi : X \rightarrow L_\beta[u]$ . Since  $\pi \upharpoonright v = \text{id} \upharpoonright v$ ,

we have:  $\pi h(z, w) \simeq h(z, w)$  for  $z \in u, w \in v$ . Hence  $\pi \upharpoonright X = \text{id} \upharpoonright X$ ;

$$\alpha X = \pi''X = L_\beta[u].$$

QED

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Thm 4  $L_\alpha[u]$  is non-projectible

iff there is a normal function

$\langle d_\nu \mid \nu < \lambda \rangle$  ( $\text{Lim}(\lambda)$ ) s.t.  $d = \sup_\nu d_\nu$ ,

and  $L_{d_\nu}[u] \preceq_{\Sigma_1} L_\alpha[u]$  for  $\nu < \lambda$ .

proof.

( $\leftarrow$ ) Let  $\varphi$  be a  $\Sigma_0$  formula. Let

$v \in L_\alpha[u]$ . Then  $v \in L_{d_\nu}[u]$  for

some  $\nu$ . For all  $x \in v$ , we have:

$$\models_{L_\alpha[u]} \forall y \varphi(y, x) \leftrightarrow \models_{L_{d_\nu}[u]} \forall y \varphi(y, x).$$

Hence, for  $w = L_{d_\nu}[u]$ :

$$\models_{L_\alpha[u]} \forall x \in w (\forall y \varphi \leftrightarrow \forall y : \varepsilon \subseteq w \varphi).$$

The remaining admissibility axioms hold trivially by the fact that  $\alpha$  is a limit ordinal.

( $\rightarrow$ ) Since the set of  $\beta < \alpha$  s.t.  ~~$L_\beta[u] \prec \sum_1 L_\alpha[u]$~~

$L_\beta[u] \prec \sum_1 L_\alpha[u]$  is closed, we need only show that it is unbounded.

~~Let~~ Let  $r < \alpha$ . Claim There is  $\beta < \alpha$

s.t.  $r < \beta$  and  $L_\beta[u] \prec \sum_1 L_\alpha[u]$ .

~~Let~~  ~~$\gamma > r$  be a limit ordinal s.t.~~

~~Let  $x \in L_\gamma[u]$ , where~~ Let  $h$  be as in Lemma 5 and let  $\gamma > r$  be a limit ordinal s.t.  $x \in L_\gamma[u]$ , where  $h$  is  $\Sigma_1$  in the parameter  $x$ . By Lemma 5:

$$h''(u \times L_\gamma[u]) = L_\beta[u] \prec \sum_1 L_\alpha[u]$$

for some  $\beta \leq \alpha$ . But, by the non projectibility of  $L_\alpha[u]$ :

$$h''(u \times L_\gamma[u]) \in L_\alpha[u],$$

hence  $\beta < \alpha$ .

QED

Thm 5  $L_\alpha[u]$  is non projectible iff

there is no  $\Sigma_1$  function which, for some  $\gamma < \alpha$ , maps a subset of  $L_\gamma[u]$  onto  $L_\alpha[u]$ .

Proof.

( $\rightarrow$ ) trivial

( $\leftarrow$ ) Let  $L_\alpha[u]$  be projectible

Then there is a  $\Sigma_0$  relation  $R$  and a  $v \in L_\alpha[u]$  s.t. for each  $x < \alpha$  there is an  $y \in v$  with:

$$\forall y R_{yx} \text{ but } \neg \forall y \in L_\gamma[u] R_{yx},$$

Let  $r$  uniformize the relation:

$$\forall y \in L_\gamma[u] R_{yx}.$$

Then  $g''v$  is unbounded in  $\alpha$ .

Let  $h$  be as in Lemma 4. Let  $h$ ,  $r$  be  $\Sigma_1$  in the parameter  $x$  and let  $v, x \in L_\gamma[u]$ , where

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$\gamma$  is a limit ordinal. By Lemma 5:

$$h'' u \times L_\gamma[u] = L_\beta[u] \preceq_{\Sigma_1} L_\alpha[u],$$

In particular,  $g'' v \subset L_\beta[u]$ ; hence  $\beta = \alpha$ , since  $g'' v$  is unbounded in  $\alpha$ .

~~Set T~~ Set  $f(\langle x, y \rangle) \simeq h(x, y)$

for  $x \in u$ ,  $y \in L_\gamma[u]$ . Then  $f$  is  $\Sigma_1$ ;

$\text{dom}(f) \subset L_\gamma[u]$  and  $\text{rng}(f) = L_\alpha[u]$

QED

We now come to the Thm announced at the outset of this section:

Thm 6  $L_\alpha[u]$  is admissible iff

there is no  $\Delta_1$  function which, for some  $\gamma < \alpha$ , maps  $L_\gamma[u]$  onto  $L_\alpha[u]$ .

Proof.

( $\rightarrow$ ) trivial

( $\leftarrow$ ) Let  $L_d[u]$  not be admissible.

Then there is a  $\Sigma_1$  relation  $R$  s.t.

$\forall y \forall x R_{yx}$  but for some  $v \in L_d[u]$ ,

there is no  $x < d$  with:  ~~$\forall y \forall x R_{yx} \wedge$~~

$\forall x \in v \forall y \in L_y[u] R_{yx}$ . Let  $r$  uniformize the relation

$$\forall y \in L_y[u] R_{yx}.$$

Then  $r``av$  is unbounded in  $d$ , or is

$\Sigma_1$  and defined everywhere. By

Thm 5, there is a  $\delta < d$  and

a  $\Sigma_1$  f s.t.  $\text{dom}(f) \subset L_\delta[u]$  and  $\text{rng}(f) = L_d[u]$ . Let:

$$y = f(x) \iff \forall z F_{zyx},$$

where  $F$  is  $\Sigma_0$ . Set:

$$\tilde{f}(v, x) =_{\text{def}} \begin{cases} y & \text{if } y \in L_y[u] \text{ and} \\ & \forall z \in L_z[u] F_{zyx} \\ \emptyset & \text{if not} \end{cases}$$

Then  $\tilde{f}$  is p.n.

~~Take  $\alpha$  as a limit ordinal large~~

Set:  $\bar{f}(\langle x, y \rangle) = \tilde{f}(r(x), y)$ ,  ~~$f(\cdot)$  is other~~

Then  $\bar{f}$  is defined everywhere and

$\bar{f}''_{\cup X L_\beta[u]} = f''_{L_\beta[u]}$ . If  $\gamma$  is  
a limit ordinal and  $\cup \subset L_\gamma[u]$ , then  
 $\bar{f}$  maps  $L_\gamma[u]$  onto  $L_\alpha[u]$ . QED

## The projectum

Def  $\alpha^* = \inf \beta$  s.t. there  
is a  $\Sigma_1(L_\alpha[u])$  function mapping  
a subset of  $\beta - \{L_\beta[u]\}$  onto  $L_\alpha[u]$ .

$\alpha^*$  is called the projectum of  $\alpha$ .

By Thm 5,  $L_\alpha[u]$  is non-projectible  
iff  $\alpha = \alpha^*$ .

Thm If  $\alpha^* > 0$ , then  $L_{\alpha^*}[u]$  is  
~~admissible~~ non-projectible.

proof. ~~If  $\alpha^* = \alpha$ ,  $L_{\alpha^*}[u]$  is~~

If  $\alpha^* = \alpha$ , the theorem is trivial.

Now let  $\alpha^* < \alpha$ . There is no  $f \in L_\alpha[u]$   
mapping a  $\gamma < \alpha^*$  onto  $\alpha^*$ , for then:

$$\begin{aligned} \text{--- } g(\langle z, v \rangle) &= \gamma(z, f(v)) \quad \text{if } z \in u, v < \gamma \\ g(z) &\neq 0 \quad \text{if not} \end{aligned}$$

would map  ~~$L_\gamma[u]$~~  onto  $L_{\alpha^*}[u]$ .

By composition, we would obtain  
a  $\Sigma_1$  map of  $L_\gamma[u]$  onto  $L_\alpha[u]$ .

(28)

But this means that  $\alpha^*$  is p.r. closed, for, as we shall show in an appendix, whenever  $\gamma$  is p.r. closed and  $\beta$  is the first p.r. closed ordinal after  $\gamma$ , each  $\gamma < \beta$  is 1-1 mappable into  $\gamma$  by a map  $f \in L_\beta$ . If  $\alpha^*$  were not p.r. closed, we should have  $\gamma < \alpha^* < \beta \leq \alpha$  for such a pair  $\gamma, \beta$ ;  ~~$f \in L_\beta^{[\kappa]}$~~ , hence some  $f \in L_\beta^{[\kappa]} \subset L_\alpha^{[\kappa]}$  would map  $\gamma$  onto  $\beta$ . But, since  $\alpha^*$  is p.r. closed, we may apply Thm 5 to  ~~$L_\alpha^{[\kappa]}$~~  to conclude that  $L_{\alpha^*}^{[\kappa]}$  is non-projectible, for otherwise there would be  $\gamma < \alpha^*$  mappable onto  $\alpha^*$  by an  $f \in L_{\alpha^*+1}^{[\kappa]}$ .

QED

## $\Sigma_n$ -admissibles

Def  $M = \langle |M|, \in, A_1, \dots, A_n \rangle$  is called  $\Sigma_n$ -admissible ( $n \geq 1$ ) iff  $M$  is admissible and satisfies the replacement axiom:

$$\Lambda x V y \varphi \rightarrow \Lambda u V v \Lambda x \in u V y \in v \varphi$$

for  $\Sigma_{n-1}$ -formulas  $\varphi$ .

(Thus, 'admissible' = ' $\Sigma_\infty$ -admissible')

Def  $M$  is called  $\Sigma_n$ -non projectible iff  $M$  is admissible and satisfies:

$$\Lambda u V v \Lambda x \in u (V y \varphi \leftrightarrow V y \in v \varphi)$$

for  $\Sigma_{n-1}$ -formulas  $\varphi$ .

We can readily establish:

(1)  $M$  is  $\Sigma_n$ -admissible iff  $\langle M, \mathbb{F}_M^{\Sigma_{n-1}} \rangle$  is admissible

(2)  $M$  is  $\Sigma_n$ -non-projectible iff  $\langle M, \mathbb{F}_M^{\Sigma_{n-1}} \rangle$  is non-projectible

(3) If  $M$  is  $\Sigma_n$ -admissible, then  $R$  is  $\Sigma_n$  iff  $R$  is  $\Sigma_1$  in  $\Sigma_{n-1}$  relations.

■ Thus, all the theorems of §1 carry over to  $\Sigma_n$ -admissibles. Some of the theorems in this section carry over.

In particular, we shall obtain slightly weaker ~~analogous~~ analogues of Thm 4 - Thm 6.

(31)

Lemma 7 If  $\langle L_\alpha[u], A \rangle$  is admissible, then  $\langle L_\alpha[u], A \rangle$  is  $\Sigma_2$ -uniformizable.

Proof.

Let  $R$  be  $\Sigma_1$ ,  $\text{rng}(R) \subset d$

Let  $Rx \vec{x} \iff \forall y P_{y \rightarrow \vec{x}}$ ,

where  $P \in \overline{\Pi}_1$ .

Set:  $p(\vec{x}) \simeq_{\mu\delta} \forall z \in L_x P_z l(\delta) \vec{x}$ ,

where  $\langle l(\delta), r(\delta) \rangle = \delta$ .

Then  $p$  is  $\Sigma_2$ , since:

$$\forall y = p(\vec{x}) \iff \underbrace{\forall z \in L_y P_z l(\delta) \vec{x}}_{\overline{\Pi}_1} \wedge$$

$$\wedge \underbrace{\forall \tau < \delta \forall z \in L_\tau P_z l(\tau) \vec{x}}_{\Sigma_1},$$

Set:  $r(\vec{x}) \simeq l(p(\vec{x}))$ . Then  $r$  uniformizes  $R$ . QED

(Note: This proof also goes through on the assumption:  $\Sigma_1$  in  $\Sigma_1 = \Sigma_2$ )

Since, if  $L_d[u]$  is  $\Sigma_m$ -admissible,  
 $\langle L_d[u], \models^{\Sigma_{m-1}} \rangle$  is admissible and  
 $\Sigma_1(\langle L_d[u], \models^{\Sigma_{m-1}} \rangle) = \Sigma_m(L_d[u])$ ,  
we get:

Corollary 7a If  $L_d[u]$  is  $\Sigma_m$

admissible, then  $L_d[u]$  is  $\Sigma_{m+1}$ -uniformizable.

Lemma 8 If  $L_d[u]$  is  $\Sigma_m$ -uniformizable, then there is a  $\Sigma_m$  function  $h$  s.t.  $\text{dom}(h) \subset u \times L_d[u]$  and

$$\wedge x (x \in h''(u \times \{x\}) \leq_{\Sigma_m} L_d[u]).$$

~~Lemma~~

Lemma 8 is proved exactly like Lemma 4, which is a special case of it.

Lemma 2 obviously holds with  $\Sigma_m$  in place of  $\Sigma_1$  ( $m \geq 1$ ), since

$X \prec_{\sum_m} L_d[u]$  implies  $X \prec_{\Sigma_1} L_d[u]$ .

Lemma 3 does not hold, but we do get the weaker form:

Lemma 9 If  $X \prec_{\sum_m} L_d[u]$  and

$\pi : \langle X, \in \rangle \xrightarrow{\sim} \langle L_d[u], \in \rangle$ , then

for every  $\Sigma_1$ -f (which is  $\Sigma_1$  in parameter  $x \in X$  s.t.  $\pi(x) = x$ ):

$$\pi f(\vec{x}) \simeq f(\pi(\vec{x})) \text{ for } \vec{x} \in X.$$

The proof is obvious.

Using Lemmas 8, 9 in place of Lemmas 4, 3, we get

Thm 7 If  $L_d[u]$  is  $\Sigma_m$ -uniformizable, then the  $\Sigma_n$  analogues of Thm 2, Thm 3 hold.

The analogues of

The proofs of Thm 2, Thm 3 can be repeated word for word to obtain Thm 7.

By ~~common~~ Lemma 7, then, the  $\Sigma_{n+1}$  analogues of Thms 2, 3 hold whenever  $L_d[u]$  is  $\Sigma_n$  admissible. We shall show later that this result can be greatly strengthened. The hypothesis of Thm 7 is always satisfied. But first we turn to the question of criteria for  $\Sigma_m$  admissibility + non ~~projective~~ projectibility.

The  $\Sigma_m$  analogue of Thm 6 does not hold. Fr. ins. letting  ~~$L_\omega$~~   $L_{\omega_\omega}[u]$  admit no function mapping an element onto the entire domain, yet

$L_{\omega_\omega}[u]$  is not admissible, since  
 $\langle \omega_n \mid n < \omega \rangle$  is  $\Sigma_2$  (understanding  
 $\omega_\omega$  in the sense of  $L[u]$ ).

The analogues of Thms 4, 5, 6 do hold,  
however, on the assumption that,  
for some  $\beta < \alpha$ ,  ~~$L_\beta[u]$~~   $L_\beta[u]$   
can be mapped onto each  
 $x \in L_\alpha[u]$  by an  $f \in L_\alpha[u]$ . Since  
 $L_\alpha[u] = L_\alpha[L_\beta[u]]$ , it suffices  
to prove this for the case:  
 $\beta = 0$  ( $L_\beta[u] = u$ ).

Def  $L_\alpha[u]$  is  $u$ -dense iff for  
all  $\gamma < \alpha$  there is an  $f \in L_\alpha[u]$   
mapping  $u$  onto  $\gamma$ .

By Lemma 1,  $u$ -density is equivalent to the condition, that  $u$  can be mapped onto each  $x \in L_d[u]$  by an  $f \in L_d[u]$ .

Lemma 10 If  $L_d[u]$  is  $u$ -dense

and  $X \prec_{\Sigma_m} L_d[u]$ , then

$$\forall \beta \leq d \quad X = L_\beta[u].$$

proof. By Lemma 2 it suffices to show that  $X$  is transitive.

Let  $x \in X$ . We wish to show:  $x \subset X$ . The statement:

$$\forall f \quad f: u \xrightarrow{\text{onto}} x$$

holds in  $\langle L_d[u], \in \rangle$ , hence in  $\langle X, \in \rangle$ . Thus there is an

$f \in X$  s.t.  $f: u \xrightarrow{\text{onto}} x$ . But then  $f(z) \in X$  for each  $z \in u$ ;

hence:  $x = f''u \subset X$ . QED

Using Lemma 10 in place of Lemma 6, we can repeat the proofs of Thms 4, 5, 6 to obtain:

(\*) If  $L_d[u]$  is  $\Sigma_m$  uniformizable, then the  $\Sigma_m$  analogues of ~~Thms~~ Thms 4, 5, 6 hold.

(The proofs can be repeated word for word).

But this enables us to prove the  $\Sigma_m$  analogues of those Thms outright. We use induction on  $m$ . For  $m=1$  the Thms are proven. Now suppose the Thms to hold for  $m$ . Then either  $m$  is admissible, or else the Thms hold trivially for all  $m \geq m$ . But if  $n$  is admissible, then by Lemma 7  $L_d[u]$  is  $m+1$  uniformizable and the Thms hold for  $m+1$  by (\*).

Thus:

Thm 8 If  $L_d[u]$  is  $u$ -dense, then the  $\Sigma_n$  analogues of Thms 4, 5, 6 hold for  $n \geq 1$ .

.....

### $u$ -uniformizability

Def A function  $\#(z, \vec{x})$  is called a  $u$ -uniformization of a relation  $R_y \vec{x}$  iff  $\text{dom}(r) = u \times \text{dom}(R)$ ,  $\text{rng}(r) \subset \text{rng}(R)$  and  $\forall y R_y \vec{x} \leftrightarrow \forall z \in u R_{r(z)} \vec{x}$ .

Def  $M = \langle |M|, \in, A_1, \dots, A_n \rangle$  (s.t.  $u \in M$ ) is  $\Sigma_n$   $u$ -uniformizable iff every  $\Sigma_n$  relation is  $u$ -uniformizable by a  $\Sigma_n$  function.

Until now we have worked with the notion of ordinal uniformisability (i.e., uniformisability of relations with ordinal range) rather than  $\alpha$ -uniformisability. ~~(1, 2)~~. However, ordinal uniformisability implies  $\alpha$ -uniformisability for  $L_\alpha[u]$  (and, indeed, the efficacy of ordinal uniformisability as a tool depends on this fact).

Lemma 11 ~~If~~  $L_\alpha[u]$  is  $\Sigma_m$  ordinal uniformisable, then  $L_\alpha[u]$  is  $\alpha$ -uniformisable.

Proof

~~Let  $R_{y\vec{x}}$  be  $\Sigma_m$ . Let~~

~~$R_{y\vec{x}} \leftrightarrow \forall z R_{zy\vec{x}}$  where  $P \in \Pi_m$ .~~

~~Set:  $\bullet P' \vec{v}\vec{x} \leftrightarrow_P P(v), (v)_1, \vec{x}$ ,~~

~~where  $(x, y)_0 = x$ ,  $(x, y)_1 = y$ . Set:~~

~~$G_{y\vec{z}\vec{x}} \leftrightarrow_P P(y(z, x)) \vec{x} \wedge z \in u$ .~~

proof of Lemma 11.

We first show that each  $\Pi_{n-1}$ -relation is  $u$ -uniformizable by a  $\Sigma_n$  function. Let  $R$  be  $\Pi_{n-1}$ .

Set:  $G \ni \vec{x} \longleftrightarrow \forall z \in u \ R\gamma(z, v) \vec{x}$ ,

$G$  is  $\Sigma_n$ . Let  $g$  uniformize  $G$ .

We may assume w.l.o.g. that  $R \neq \emptyset$ , hence that  $y \in \text{rng}(R)$ .

Set:

$$r(z, \vec{x}) \simeq \begin{cases} \gamma(z, g(\vec{x})) & \text{if } R\gamma(z, g(\vec{x})) \vec{x} \\ y & \text{if } \neg R\gamma(z, g(\vec{x})) \vec{x} \end{cases}.$$

Then  $r$  uniformizes  $R$ . ~~Now let~~

Now let  $R$  be  $\Sigma_n$ . Let:

$$R\gamma \vec{x} \longleftrightarrow \forall z \ Pz\gamma \vec{x},$$

where  $P$  is  $\Pi_{n-1}$ . Set:

$$P'(z, y) \vec{x} \longleftrightarrow Pz\gamma \vec{x}$$

and let  $p$  uniformize  $P'$ . Set:

$$r(w, \vec{x}) \simeq (p(w, \vec{x}))_1$$

(where  $\langle\langle z, y \rangle\rangle_0 = z$ ,  $\langle\langle z, y \rangle\rangle_1 = y$ ).

Then  $r$   $\alpha$ -uniformizes  $R$  QED

All previous theorems in which ordinal uniformizability was mentioned as an assumption hold on the (apparently) weaker assumption of  $\alpha$ -uniformizability. In particular:

Lemma 12 If  $L_d[u]$  is  $\Sigma_n$   $\alpha$ -uniformizable, then there is a  $\Sigma_n$  function  $h$  s.t.  $\text{dom}(h) \subset u \times L_d[u]$  and

$$\forall x \in L_d[u] \quad (x \in h''(u \times \{x\})) \leq_{\Sigma_n} L_d[u]$$

proof. We imitate the proof of Lemma 4.

Letting  $s(z, x)$  s.t.  $s : u \times \{x\} \xrightarrow{\text{onto}} Fm|_{u \cup \{u, x\}}^{\Sigma_n}$

be as before, we set

$$R_x \varphi \longleftrightarrow Fm|_{L_d[u]}^{\Sigma_n} \varphi(v_0/x)$$

and let  $r$  uniformize  $R$ . Set:

$$h(\langle\langle z, w \rangle\rangle, x) \simeq r(z, s(w, \langle x, p \rangle)),$$

then  $r$   $\alpha$ -uniformizes  $L_d[u]$  QED

Carrying through the earlier proofs, again virtually without change, we get:

Thm 9 Let  $L_\alpha[u]$  be  $\Sigma_m$   $u$ -uniformizable.

Then the following are equivalent

(a) There is a  $\Sigma_m$  set  $a \subset u$  s.t.

$$a \notin L_\alpha[u]$$

(b) There is a  $\Sigma_m$  map  $f$  s.t.

$$\text{dom}(f) \subset u \text{ and } f''u = L_\alpha[u].$$

Thm 10 Let  $L_\alpha[u]$  be  $\Sigma_m$   $u$ -uniformizable.

Then the following are equivalent

(a) There is a  $\Delta_m$  set  $a \subset u$  s.t.  $a \notin L_\alpha[u]$

(b) There is a  $\Delta_m$  map of  $u$  onto  $L_\alpha[u]$ .

We now prove:

Thm 11  $L_\alpha$  is  $\Sigma_m$   $u$ -uniformizable  
 $(n \geq 1)$ ,

The proof of Thm 11 extends over several lemmas.<sup>9</sup> From now on, we shall write 'uniformisable' to mean ' $\alpha$ -uniformisable'.

Lemma 13 Let  $L_\alpha[u]$  be admissible and let  $A \subset L_\alpha[u]$  be s.t.

$$x \in L_\alpha[u] \rightarrow A \cap x \in L_\alpha[u].$$

Then  $\frac{\Sigma_0}{A}$  is  $\Delta_1$  in  $\langle L_\alpha[u], A \rangle$ . Moreover,  $R$  is  $\Sigma_1$  in  $\langle L_\alpha, A \rangle$  iff  $R$  is  $\Sigma_1$  in  $\langle L_\alpha[u], \frac{\Sigma_0}{u} \rangle$ .

proof.

We first show that  $\frac{\Sigma_0}{A}$  is  $\Delta_1$ .

Set:  $a(x) =_{\text{df}} A \cap x$ .

$L_\alpha[u]$  is closed under  $a$ .  $a$  is  $\Sigma_1$ , since

$$y = a(x) \iff y \subset x \wedge \forall z \in x (z \in y \iff A z).$$

Thus  $\frac{\Sigma_0}{A}$  is  $\Delta_1$ , since

$$\models_A^{\Sigma_0} \varphi \leftrightarrow \models^{\Sigma_0} \langle c(\varphi), a(c(\varphi)) \rangle .$$

But, by the same argument,

$$\models_{\models_A^{\Sigma_0}}^{\Sigma_0} \text{ is } \Delta_1 \text{, since :}$$

$$\models_A^{\Sigma_0} \psi \leftrightarrow \models \langle c(\psi), \models_{\models_A^{\Sigma_0}}^{\Sigma_0} \langle c(\psi), a(c(\psi)) \rangle \rangle \psi$$

This establishes the second part  
of the lemma. QED

Using Lemma 13, we can repeat the  
proofs of Thm 1 and ~~Lemma 11~~  
to obtain the analogues:

Lemma 14. If  $L_d[u], A$  are as in Lemma 13, then  $\langle L_d[u], A \rangle$  is  $\Sigma_1$  uniformizable.

Since the only two facts used in the proof of ~~Lemma~~ Lemma 12 were:  $\Sigma_m$  uniformizability and the  $\Sigma_m$  definability of  $F^{\Sigma_m}$ , we may repeat the proof to obtain:

Lemma 15. If  $L_d[u], A$  are as in Lemma 13 and if  $\langle L_d[u], A \rangle$  is  $\Sigma_m$  uniformizable, then there is a  $\Sigma_m$  Skolem function (i.e. an  $h$  s.t.  $\text{dom}(h) \subset u \times L_d[u]$ , and  $\lambda x (x \in h^{cc}(u \times L_d[u]) \prec_{\Sigma_m} \langle L_d[u], A \rangle)$ .

In particular, by Lemma 14, there is a  $\Sigma_1$  Skolem function.

Def  $\langle L_d[u], A \rangle$  is called feasible

iff for every  $\Delta_1$  set  $B$  we have:  
 $x \in L_d[u] \rightarrow B \cap x \in L_d[u]$

$\langle L_d[u], A \rangle$  is called  $\Sigma_m$ -feasible if

this holds for every  $\Delta_m B$ .

Lemma 16 Let  $\langle L_\alpha[u], A \rangle$  be  $\Sigma_m$ -feasible but not  $\Sigma_m$  admissible. Let  $\langle L_\alpha[u], A \rangle$  be  $\Sigma_m$  uniformisable. Then a relation  $R$  is  $\Sigma_1$  in  $\Sigma_m$  iff  $R$  is  $\Sigma_{m+1}$ .

Proof.

( $\leftarrow$ ) trivial, since each  $\Sigma_{m+1}$  relation is  $\Sigma_1$  in  $\Sigma_m$

( $\rightarrow$ ) Since  $\langle L_\alpha[u], A \rangle$  is not  $\Sigma_m$  admissible, there is a  $\Pi_{m+1}$  relation  $R$  and a  $\beta < \alpha$  s.t.  $\forall x \forall y R y x$  but for each  $\gamma < \beta$ :  $\forall x \in L_\beta[u] \exists y \in L_\gamma[u] R y x$ .

Set:  $G_{y x} \leftrightarrow_{\text{df}} \exists y \in L_\gamma[u] R y x$

and let  $g$  uniformise  $G$ . Then

$g$  is  $\Delta_m$ ,  $\text{dom}(g) = u \times L_\beta[u]$  and

$g'' u \times L_\beta[u] = L_\alpha[u]$ . Let  $h$  be a

$\Sigma_m$  Skolem function for  $\langle L_\alpha[u], A \rangle$ .

( $h$  exists by Lemma 15).

Let :

$$y = h(z, x) \longleftrightarrow \forall w \exists v \forall q z = x,$$

where  $H$  is  $\text{PT}_{M-1}$ . Set :

$$h^*(\langle w, z \rangle, x) = \begin{cases} y & \text{if } y \in L_g(w) \wedge \\ & \wedge \forall v \in L_g(w) H v y = x \\ \emptyset & \text{if not} \end{cases}$$

Then  $\text{dom}(h^*) = (u \times L_\beta[u]) \times L_\alpha[u]$  and

~~h^\*(x)~~

$$h^{**}(u \times L_\beta[u]) \times \{x\} = h^{**} u \times \{x\}$$

for all  $x \in L_\alpha[u]$ .

For  $\gamma < \alpha$  set :

$$\bar{\gamma} = (u \times L_\beta[u]) \times L_\gamma[u]$$

$$e(\gamma) = \{\langle x, y \rangle \mid x, y \in \bar{\gamma} \wedge h^*(x) \in h^*(y)\}$$

$$a(\gamma) = \{x \mid x \in \bar{\gamma} \wedge A h^*(x)\}.$$

Since, for each  $\gamma < \alpha$ ,  $e(\gamma), a(\gamma)$  are  $\Delta_m$  subsets of  $\bar{\gamma}$ , we have :

$$e(\gamma), a(\gamma) \in L_\alpha[u] \quad \text{for } \gamma < \alpha.$$

Let  $m(x) =_D \langle v, e(v), a \rangle$ , where

$v \in v$  and for some  $\pi$ :

$$\pi : \langle \bar{x}, e(x), a(x) \rangle \xrightarrow{\sim} m(x).$$

By the admissibility of  $L_d[u]$  (Thm 3, Thm 6), we may conclude that  $m(x)$ ,  $\pi$  are elements of  $L_d[u]$ . This follows by the recursion theorem, since the factorisation of  $e(x)$  by extensional equivalence is certainly in  $L_d[u]$  and the ~~factorised~~ factorised  $e(x)$  is well founded.

Thus,  $L_d[u]$  is closed under the function  $m(x)$ . We show now that  $m$  is  $\Delta_{m+1}$ .  $e$  is  $\Sigma_{m+1}$ , since:

$$y = e(r) \longleftrightarrow y \in \bar{x}^z \text{ --- } \wedge$$

$$\wedge \forall z, w \in \bar{x} \underbrace{(\langle z, w \rangle \in y \leftrightarrow h^*(z) \in h^*(y))}_{\Delta_m}$$

$\Pi_m$

(49)

Similarly,  $a(\bar{x})$  is  $\Delta_{m+1}$ .

This means that  $m$  is  $\Delta_{m+1}$ , since:

$$y = m(\bar{x}) \longleftrightarrow \forall \pi (\pi : \langle \bar{x}, e(\bar{x}), a(\bar{x}) \rangle \rightsquigarrow y)$$

$\pi : x \rightsquigarrow y$  being  $\Delta_1$ .

To establish the lemma, we need only show that relations  $\Sigma_0$  in  $\Sigma_m$  relations are  $\Sigma_{m+1}$ .

Let the formula  $\varphi$  be  $\Sigma_0$  in  $\Sigma_m$  (i.e. built up from  $\Sigma_m$  formulae by sentential operations and bounded ~~quantifications~~ quantifications).

Then

$$\vdash_{\langle L_d[u], A \rangle} \varphi \longleftrightarrow \forall \bar{x} < d (\lim(\bar{x}) \wedge \varphi \in L_{\bar{x}}[u] \wedge \underset{m(\bar{x})}{\models} \varphi)$$

QED

Note that the assumption:  $\Sigma_{m+1} = \Sigma_1$  in  $\Sigma_m$  can be used alternatively to  $\Sigma_m$  admissibility to carry out the proof of Lemma 7, hence:

Lemma 17 If  $\langle L_d[u], A \rangle$  is  $\Sigma_m$ -feasible, then  $\langle L_d[u], A \rangle$  and  $\Sigma_m$ -uniformizable, then  $\langle L_d[u], A \rangle$  is  $\Sigma_{m+1}$  uniformizable.

We are now ready to prove Thm 11. We proceed by induction on  $m$ . For  $m=1$  the theorem is proven. We now suppose it to hold for  $m$  and prove it for  $m+1$ .

Case I.  $L_d[u]$  is  $\Sigma_m$ -feasible.

The conclusion follows by Lemma 17.

If Case I fails, there is a  $\beta < \alpha$  s.t. a  $\Delta_n$   $a \in L_\beta[u]$  exists with  $a \notin L_\alpha[u]$ . Let  $\beta$  be the least such. By Thm 10, there is a  $\Delta_1$  map from  $L_\beta[u]$  onto  $L_\alpha[u]$ .

Case II.  $\beta = 0$  (hence  $L_\beta[u] = u$ ).

We first show that each  $\Delta_n$  relation is uniformizable by a  $\Sigma_m$  function. Let  $R_y \vec{x}$  be  $\Delta_n$ .

Assume (w.l.o.g.)  $y \in \text{rng}(R)$ .

Set:  $r(z, \vec{x}) \leftarrow \begin{cases} f(z) & \text{if } Rf(z) \vec{x} \\ y & \text{if } \neg Rf(z) \vec{x} \end{cases}$

Then  $r$  uniformizes  $R$ . If  $R \in \Sigma_m$ , there is  $\Pi_{m-1} P$  s.t.

$$R_y \vec{x} \leftrightarrow \forall z Pz y \vec{x}.$$

$$\text{Set } P'(z, y) \vec{x} \leftrightarrow_P Pz y \vec{x}.$$

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Let  $p$  uniformize  $P'$  and set:

$$r(\omega, \vec{x}) = (p(\omega, \vec{x})), \quad \text{QED Case II}$$

Case III  $\beta > 0$ .

Then  $L_\beta[u]$  will be admissible by the same argument which demonstrated that the projectum of  $d$  is admissible.

Lemma 18 If  $A$  is  $\Delta_m(L_\alpha[u])$  and  $A \subset L_\beta[u]$ , then each  $R \in L_\beta[u]^\sim$  which is  $\Sigma_m(\langle L_\beta[u], A \rangle)$  is  $\Sigma_{m+1}(L_\alpha[u])$ .

proof.

It suffices to show: If  $R$  is  $\Sigma_0(\langle L_\beta, A \rangle)$ , then  $R \in \Sigma_{m+1}(L_\alpha)$ .

Let  $\varphi$  be a  $\Sigma_0$  formula of  $\langle L_\beta, A \rangle$ .

Then

$$\models_A \varphi \longleftrightarrow \models_{\langle C(\varphi), A \cap C(\varphi) \rangle} \varphi.$$

But  $a(u) = A \cap u$  is a  $\Sigma_2(L_\alpha[u])$

function which is defined on all  $L_\beta[u]$ . Hence  $\models_A^{\Sigma_0}$  is  $\Sigma_2(L_\alpha[u])$

QED

Letting  $f : L_\beta \xrightarrow{\text{onto}} L_\alpha$  be  $\Delta_m(L_\alpha)$ ,  
pick  $A \subset L_\beta$  in such a way that:

$$\{(x, y) \mid f(x) \in f(y)\}, \quad f^{-1} \upharpoonright L_\beta[u],$$

$$\models_f^{\Pi_m} =_{\text{pt}} \{\varphi \in \text{Fml}_{L_\beta}^{\Pi_{m-1}} \mid \models_{L_\alpha} \bar{f}(\varphi)\}$$

are  $\Sigma_0$  in  $A \langle L_\beta, A \rangle$ . (Setting

$$\bar{f}(\varphi(\vec{x})) =_{\text{pt}} \varphi(f(\vec{x})).$$

Then every  $\Sigma_m(L_\alpha[u])$  relation

$$R \subset L_\beta[u]^m \text{ is } \Sigma_1(\langle L_\beta, A \rangle).$$

Using an obvious abbreviation, we have then:

$$\Sigma_m(L_\alpha) \subset \Sigma_1(\langle L_\beta, A \rangle) \subset \Sigma_{m+1}(L_\alpha).$$

On which side, if any, of this chain of conclusions does the identity lie?

We consider two cases:

Case 1 There is a  $\gamma < \beta$  and a  $\Sigma_m(L_\alpha)$  function  $g$  s.t.  $\text{dom}(g) = L_\gamma[u]$  and  $g''L_\gamma[u]$  is unbounded in  $\alpha$ .

In this case, we prove that, for an appropriate choice of  $A$ :

$$\Sigma_1(\langle L_\beta, A \rangle) = \Sigma_{n+1}^2(L_\alpha).$$

But, by Lemma 14,  $\langle L_\beta, A \rangle$  is  $\Sigma_1$  uniformizable.

Case 2 Case 1 fails.

In this case we show that

$$\Sigma_m(L_\alpha) = \Sigma_1(\langle L_\beta, A \rangle).$$

But then  $\langle L_\beta, A \rangle$  is feasible and, by Lemmas 94, 17,  $\langle L_\beta, A \rangle$  is  $\Sigma_2$  uniformizable, whereby:

$$\sum_{n+1}^{\alpha}(L_\alpha) = \Sigma_2(\langle L_\beta, A \rangle).$$

In either case, we may conclude that, if  $R \subset L_\beta[u]^{m+1}$  is  $\Sigma_{n+1}(L_\alpha)$ , then  $R$  is uniformizable by a  $\Sigma_{n+1}(L_\alpha)$  function. Now let  $R \subset L_\alpha[u]^{m+1}$ .

$$\text{Set: } R' \xrightarrow{y} \xrightarrow{x} R f(y) \xrightarrow{f(x)}.$$

Let  $r'$  uniformize  $R'$ . Let  $f'$

uniformize:  $f(x) = y$  and set:

$$r(\langle z, w \rangle, \vec{x}) \simeq f(r'(z, f'(w_1, x_1), \dots, f'(w_m, x_m)),$$

Then  $r$  uniformizes  $R$ .

Thus, it remains only to prove the assertions made in Cases 1, 2.

Lemma 19. Let  $\kappa < \beta$  and let there be a  $\Delta_m(L_\alpha)$  function  $g$  which maps  $L_\kappa[u]$  onto  $L_\alpha[u]$ . Then  $A$  can be so chosen that every  $R \subset L_\beta[u]^m$  which is  $\Sigma_{m+1}(L_\alpha)$  is  $\Sigma_1(\langle L_\beta, A \rangle)$ .

Proof. It suffices to show: If  $R \subset L_\beta[u]^m$  is  $\Sigma_m(L_\alpha)$ , then  $R$  is  $\Delta_m(\langle L_\beta, A \rangle)$ . For this, it suffices that

$$\{\varphi \in \text{Fml}_{L_\beta[u]}^{\Sigma_m} \mid F_{L_\beta[u]} \varphi\} \quad \text{is}$$

in  $\Delta_1(\langle L_\beta, A \rangle)$ . Let  $h$  be — a  $\Sigma_m$  Shalem function for  $L_\alpha[u]$ .

~~Let  $\varphi \in \text{Fml}_{L_\beta[u]}^{\Sigma_m} \rightarrow V \models \varphi$~~

Let  $y = h(z, x) \iff \forall v H v y z x$ ,

where  $H$  is  $\Pi_{n-1} (L_d[u])$ . Define:

$$h^*(\langle w, z \rangle, x) = \begin{cases} y & \text{if } y \in L_{g(w)}[u] \text{ and} \\ & \forall v \in L_{g(w)} H v y z x \\ & \text{o if not} \end{cases}$$

$h^*(z, x) = o$  in all other cases.

Then  $\text{dom}(h^*) = L_x[u] \times L_d[u]$  and

$$\forall x \quad h^{**}(L_x[u] \times \{x\}) = h^{**}(u \times \{x\}).$$

Choose  $A$  in such a way that  $R$  is  $\Sigma_0 \dot{\cup} (\langle L_\beta, A \rangle)$ , where:

$$Rxy \iff x, y \in L_x \times L_\beta \wedge h^*(x) \in h^*(y).$$

$R$  is  $\Delta_m(L_\alpha)$ .

$$\text{Set: } \bar{\gamma} = L_x \times L_\gamma \quad \left. \right\} \text{ for } \gamma < \beta.$$

$$e(\gamma) = \bar{\gamma}^2 \cap R$$

Then  $\forall \gamma < \beta e(\gamma) \in L_\beta[u]$ , since

$e(\gamma)$  is a  $\Delta_m(L_\alpha)$  subset of  $\bar{\gamma}^2$ .

The function  $e(\bar{x})$  is  $\Sigma_1(L_\beta, A)$ , since:

$$e = e(\bar{x}) \leftrightarrow e \in \bar{x}^2 \wedge \forall xy \in \bar{x} (\langle x, y \rangle \in e \leftrightarrow R_{xy}).$$

Set:  $m(\bar{x}) = \text{that } \langle v, \epsilon \rangle \text{ s.t. } v \in v$

$$\text{and } \langle v, \epsilon \rangle \leftrightarrow \cancel{\langle \bar{x}, e(\bar{x}) \rangle}$$

$$\leftrightarrow \langle h^*\bar{x}, \epsilon \rangle.$$

Imitating the methods of the proof of Lemma 6, we get: The function  $m$  is  $\Sigma_1(\langle L_\beta, A \rangle)$  and is defined everywhere. But this means

that  $\{\varphi \in \text{Fml}_{L_\beta}^{\Sigma_m} \mid \models_{L_\alpha}^\Sigma \varphi\}$  is  ~~$\Sigma_1(L_\beta, A)$~~ ,

~~$\Delta_1(L_\beta, A)$~~ , since  $\bar{x}$ ,

setting  $\bar{x}(\varphi) = \mu \bar{x} (\text{fim}(\bar{x}) \wedge \varphi \in L_\beta[u])$ ,

we have:

$$\models_{L_\alpha}^{\Sigma_m} \varphi \leftrightarrow \models_{m(\bar{x}(\varphi))}^{\Sigma_m} \varphi$$

for  $L_\beta[u] - \text{fmlae } \varphi$ .

QED

Lemma 20 If the hypothesis of Lemma 19 fails, then every  $\Sigma_1(\langle L_\beta, A \rangle)$  relation is  $\Sigma_m(L_\alpha)$  (hence  $\langle L_\beta, A \rangle$  is feasible).

proof.

It suffices to show: If  $R$  is  $\Sigma_0(\langle L_\beta, A \rangle)$ , then  $R$  is  $\Delta_m(L_\alpha)$ . We show this by induction on the defining formula of  $R$ , using

(\*) If  $R_y \vec{x}$  is  $\Sigma_m(L_\alpha)$ , then

so is:

$$\bar{R}_y \vec{x} \leftrightarrow_{P_f} y \in L_\beta[u] \wedge \exists \vec{z} R \vec{z} \vec{x}.$$

proof of (\*):

Let  $R_y \vec{x} \leftrightarrow \forall u P_u y \vec{x}$ , where

$P$  is  $\Pi_{m-1}$ . For  $y \in L_\beta[u]$ , we

have:

$$\cancel{\exists \vec{z} R \vec{z} \vec{x}} \rightarrow \forall u \cancel{\exists \vec{z} R \vec{z} \vec{x}} \cancel{V_{\vec{z} \in L}}$$

$$\lambda z \in y \vee u P \vec{z} \rightarrow \forall x \exists a \lambda z \in y \forall u \in L_x P \vec{z} \vec{x}.$$

since otherwise, letting  $p(w, z)$  uniformize the relation:

~~P'~~  $\rightarrow$

$$P' z \leftrightarrow_{\text{nt}} \forall u \in L_z [u] P u z \vec{x},$$

$p$  would map  $u \in y$  unboundedly into  $a$ . (Contradiction!)

Hence:

$$\lambda z \in y P \vec{z} \vec{x} \leftrightarrow \forall w \lambda z \in y \forall u \in w P u z \vec{x}.$$

We apply the same reduction to  $\forall u \in w P u z \vec{x}$  etc. until we are left with a  $\Sigma_m$  formula. QED