

§3.2 Some Properties of Weasel

In §3.1 we introduced the limit Q_∞ of certain weasel iterations $\langle Q_i \rangle$ of length ∞ . We begin this section by extending that notion and developing the theory of long iterations more fully.

Def Let Q be a mouse or weasel.

By a long iteration of Q we mean an iteration $\langle Q_i \mid i < \infty \rangle$ with indices $\langle \nu_i, \alpha_i \rangle_{i < \infty}$.

(*) $E_{\nu_i} \neq \emptyset$ for arbitrarily large $i < \infty$.

Any other iteration is called short.

(Note An iteration of length ∞ which does not satisfy (*) will eventually become constant & hence can be replaced by an iteration of length $< \infty$).

Lemma 1.1 Let $\langle Q_i \mid i < \omega \rangle$ be a long iteration with indices $\langle \nu_i, d_i \rangle$ and κ_i . Then

$$\bigwedge \alpha \bigvee \beta \bigwedge i \geq \beta \quad \kappa_i \geq \alpha.$$

prf.

Suppose not. Then there is κ s.t.

$\{i \mid \kappa_i = \kappa\}$ is unbounded in ω .

Let κ be the least such & choose

i_0 s.t. $\kappa_j \geq \kappa$ for $j \geq i_0$. Let $i \geq i_0$

s.t. $\kappa_i = \kappa + \nu_i$ is minimal for i with

this property. Let $j > i$ be least s.t.

$\kappa_j = \kappa$. Then $\nu_j = \kappa + Q_j$ and $E^{Q_j} = \emptyset$.

Hence $\nu_j < \nu_i$. Contr! QED ν_i

Def Let $\langle Q_i \mid i < \omega \rangle$ be as above.

Set $\beta_\alpha =$ the least β s.t.

$\kappa_i \geq \alpha$ for all $i \geq \beta$. Set

$$Q_\infty = \bigcup_{\alpha} \bigcap_{\beta} E^{Q_{\beta_\alpha}}$$

Then Q_∞ is a weakl.

[Note $C = \{\alpha \mid \alpha = \beta_\alpha\}$ is a club class.]

If $\pi_{i,j}$ are the iteration maps we define $\pi_{i,\infty}$ ($i < \infty$) by:

$$\pi_{i,\infty}(x) \approx \begin{cases} \pi_{i,\beta_d}(x) & \text{for } d \text{ s.t.} \\ \pi_{i,\beta_d}(x) \in J_d \in E^{\mathcal{Q}_{\beta_d}} \end{cases}$$

This definition extends our earlier one as shown by:

Lemma 1.2 The following are equivalent:

(a) $\bigwedge_{j \geq i} \pi_{0,i}(j) < \pi_j$

(b) $\mathcal{Q}_\infty, \langle \pi_{i,\infty} \rangle = \lim_{i \leq j} (\mathcal{Q}_i, \pi_{i,j})$

(c) $\text{dom}(\pi_{0,\infty}) = \mathcal{Q}$.

(d) $\infty \in \text{dom}(\pi_{0,\infty})$.

The proof is straightforward.

If (a) - (d) hold, we call \mathcal{Q}_∞ a simple iterate of \mathcal{Q} by the iteration $\langle \mathcal{Q}_i \rangle$.

Note Q_∞ can fail to be a simple iterate even if $\langle Q_i \mid i < \infty \rangle$ is a simple iteration. The reason is that Q_∞ is really obtained by an "iteration" of length $\infty + 1$, where "truncation" can occur at the last stage. One can think of a non simple iterate Q_∞ as follows: Form the direct limit Q' of $\langle Q_i \mid i < \infty \rangle$. Q' properly extends ∞ and we truncate it to get Q_∞ . (If ∞ were an inaccessible cardinal we could carry this out explicitly with transitive Q').

The following facts are obvious:

Fact 1 Any long iteration results in a weasel.

Fact 2 Any simple iterate of a weasel (mouse) is a weasel (mouse).

Fact 3 Any short non simple iteration results in a mouse.

Fact 4 If Q is a non simple iterate of a weasel W with iteration map π , then there is a set Q is a non simple iterate of W with iteration map π .

We now prove:

Lemma 2.1 Let W' be a simple iterate of the weasel W with iteration map π and W'' a simple iterate of W' with map π' . Then W'' is a simple iterate of W with iteration map $\pi \pi'$.

proof of Lemma 2.1

If the first iteration is short, the conclusion is trivial, so suppose $W' = W_\infty$, where $\langle W_i \mid i < \infty \rangle$ is a long iteration. Set:

$C = \{ \alpha \mid \beta_\alpha = \alpha \wedge \alpha \text{ is a cardinal in } W \}$,
where $\beta_\alpha =$ the least β s.t. $\kappa_i \geq \alpha$
for $i \geq \beta$ and $|\{ i < \beta \mid E_{\nu_i} \neq \emptyset \}| \geq \alpha$.

($|A| =$ order type of A). Now let

$\langle W'_i \mid i \leq \theta \rangle$ be the iteration from W' to W'' . If $\theta = \infty$, define C' the way C was defined. If $\theta < \infty$, set:

$C' =$ the set of cardinals α in W'
s.t. $\alpha > \nu'_i$ for $i < \theta$.

[Here ν_i, κ_i and ν'_i, κ'_i are indices of the iterations $\langle W_i \rangle, \langle W'_i \rangle$ resp.]

Set: $\tilde{C} = C \cap C'$. For $\alpha \in \tilde{C}$:

$J_\alpha^{E^{W_\alpha}} = J_\alpha^{E^{W'}}$ is an iterate of $J_\alpha^{E^W}$
with iteration map $\pi \upharpoonright J_\alpha^{E^W}$

and:

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$J_\alpha^{EW'} = J_\alpha^{EW''}$ - 7 -
 is an iterate of J_α^{EW}
 with iteration map $\pi' \uparrow J_\alpha^{EW'}$.

These iterations are short & hence
 $J_\alpha^{EW''}$ is a short iterate of J_α^{EW}
 with map $\pi' \pi \uparrow J_\alpha^{EW}$. Let

$\langle \tilde{\nu}_i^\alpha \mid i < \alpha \rangle$ be the indices of the
normal iteration from J_α^{EW} to $J_\alpha^{EW''}$.

If $\alpha < \beta$, the models coincide
 up to α , but $\kappa_i^\beta \geq \alpha$ for $i \geq \alpha$.

Thus $\tilde{\nu}_i^\alpha = \tilde{\nu}_i^\beta$ for $i < \alpha$. So!

$\tilde{\nu}_i = \tilde{\nu}_i^\alpha$ for an $\alpha \in C$ s.t. $i < \alpha$.

It follows easily that W'' is
 an iterate of W by an iteration
 with indices $\langle \tilde{\nu}_i \rangle$, giving
 iteration map $\pi' \pi$.

QED (Lemma 2.1)

An easy modification of this
 proof yields:

Lemma 2.2 Let W^0 be a weasel & let W^i be a simple iterate of W^i with map $\pi_{i,j}$ for $i \leq j < \theta \leq \infty$. Set: $W, \langle \pi_i \rangle = \lim_{i \leq j} (W_i, \pi_{i,j})$.

Then W is a simple iterate of W_i with iteration map π_i ($i < \theta$).

Using Fact 4 above we get:

Lemma 2.3 There is no sequence $\langle Q_i \mid i < \omega \rangle$ s.t. Q_{i+1} is a non-simple iterate of Q_i for $i < \omega$.

prf. Suppose not.

If Q_0 is a weasel, replace it by a mouse $Q'_0 = Q_0 \upharpoonright \alpha_0$ s.t. Q_1 is a non-simple iterate of Q'_0 . If Q_1 is a weasel, replace it by $Q'_1 = Q_1 \upharpoonright \alpha_1$ s.t. Q_2 is a nonsimple iterate of Q'_1 . Then Q'_1 is a nonsimple iterate of Q'_0 . Proceeding in this way we get a sequence $\langle Q'_i \mid i < \omega \rangle$ of mice s.t. Q'_{i+1} is ~~an~~ a nonsimple iterate of Q'_i for $i < \omega$. Contradiction!

(Q.E.D.) / Lemma 2.3,

[Note It is also easy to prove:

Lemma 2.4 Let Q' be a nonsimple iterate of $Q + Q''$ a nonsimple iterate of Q' with maps π, π' resp. Then Q'' is a nonsimple iterate of Q with map $\pi' \pi$,

However, the corresponding theorem for "mixed" sequences of simple and nonsimple iterates is false.]

Now write $\langle Q, \nu \rangle R \langle \bar{Q}, \bar{\nu} \rangle$ to mean that Q is an iterate (long or short) of \bar{Q} with map π and that either Q is a nonsimple iterate, or $\nu < \pi(\bar{\nu})$. By Lemma 2.3 it is impossible that $\langle Q_{i+1}, \nu_{i+1} \rangle R \langle Q_i, \nu_i \rangle$ for $i < \omega$. But then a virtual repetition of the proof of § 2.1 Lemma 3 yields:

Lemma 3 Let W_0, \dots, W_n be words
 + let W_i be an iterate of W_{i-1}
 with iteration map π_i for $1 \leq i \leq n$,
 Set: $\pi = \pi_n \circ \dots \circ \pi_1$. Let
 $\sigma: W_0 \xrightarrow{\Sigma_1} W_n$. Then W_i is a
 simple iterate of W_{i-1} ($1 \leq i \leq n$)
 and $\sigma(\xi) \geq \pi(\xi)$ for all ξ .

Cor 3.1 If W is a simple iterate of \bar{W} ,
 it cannot be a nonsimple iterate
 of \bar{W} . In fact, if $\bar{W} = W_0, \dots, W_n = W$
 is a chain s.t. W_i is an iterate
 of W_{i-1} for $1 \leq i \leq n$, then W_i is a
 simple iterate of W_{i-1} for $1 \leq i \leq n$.

Cor 3.2 If W' is a simple iterate
 of W , then the iteration map is
 unique.

Def $\pi_{WW'}$ = the unique iteration
 map from W to W' if W' is
 a simple iterate of W .

Def $W = \lim_{i < \theta} W_i$ means:

$W, \langle \pi_{W_i, W} \rangle = \lim_{i \leq j < \theta} (W_i, \pi_{W_i, W_j})$,
 where W_j is a simple iterate of W_i
 for $i \leq j < \theta \leq \infty$.

Thus we can deal with weasels in virtually the same way as mice. We carry this thought further by extending the relations \sim_* , $<_*$ of § 2.4 to weasels. We there defined $M \sim_* N$ to hold iff M, N possess a common simple iterate. Since we cannot quantify over weasels this definition is unavailable, so we instead define:

Def $W \sim_* W'$ iff W, W' coiterate to a common simple iterate.

The force of the missing quantifier is then given by:

Lemma 4.1 Let W, W' have a common simple iterate W'' . Then $W \sim_* W'$,
 prf. Suppose not,

let W, W' coiterate to \tilde{W} and suppose
 e.g. that \tilde{W} is a simple iterate of
 W but not of W' . Coiterate W'', \tilde{W}
 to W^* . Since W^* is an iterate of \tilde{W} ,
 it is not a simple iterate of W' by
 Cor 3.1. Hence W^* is not a simple
 iterate of W'' . But then W^* is not
 a simple iterate of W by Cor 3.1, hence
 not a simple iterate of \tilde{W} . Contr!

QED (Lemma 4.1)

If W is a weasel and Q a mouse,
 we regard $W \sim_* Q, Q \sim_* W$ as
 false (since they are, indeed, false
 in the sense of the above definition)

Def Let W be a weasel and Q a
 mouse or weasel.

$W <_* Q$ iff W, Q coiterate
 to a W' which is a simple
 iterate of W but not of Q

Finally:

Def Let Q be a mouse & W a weasel

$Q <_* W$ iff there is a mouse Q' which is a simple iterate of Q and a nonsimple iterate of W .

(In other words: $Q <_* W \iff Q <_* W \upharpoonright \alpha$ for some α)

By a proof very much like that of Lemma 4.1 we get:

Lemma 4.2 Let Q, Q' be mice or weasels. Let Q'' be a simple iterate of Q and let $Q' = Q_0, \dots, Q_n = Q''$ be s.t. Q_i is an iterate of Q_{i-1} for $1 \leq i \leq n$ and Q_i is a nonsimple iterate of Q_{i-1} for some i . Then $Q <_* Q'$.

The proof is left to the reader.

Using these lemmas we can virtually repeat the proofs of §2.4 Lemma 1, Lemma 2 to get:

Lemma 4.3 $<^*$ is a linear ordering of mice and weasels modulo the congruence relation. In other words:

$$(a) Q \sim_* Q$$

$$(b) Q \sim_* Q' \rightarrow Q' \sim_* Q$$

$$(c) Q \sim_* Q' \sim_* Q'' \rightarrow Q \sim_* Q''$$

$$(d) Q <_* Q' \vee Q \sim_* Q' \vee Q' <_* Q$$

$$(e) Q \sim_* Q' \rightarrow Q \not<_* Q'$$

$$(f) Q <_* Q' \rightarrow Q' \not<_* Q$$

$$(g) Q <_* Q' <_* Q'' \rightarrow Q <_* Q''$$

If W is ^{weakly} universal, then by definition $W \not<_* Q$ for every mouse Q . But then $W \not<_* W'$ for any weasel W' , since otherwise $W <_* W' \mid d$ for some d . Hence $Q \leq_* W$, whenever Q is a mouse or weasel. If W' is a non-^{weakly} universal weasel, then $W' \not<_* W$, since otherwise

There is a mouse \mathcal{Q} s.t., $W' <_* \mathcal{Q}$
and we would have: $W <_* \mathcal{Q}$,
Hence:

Lemma 4.4 The ^{weakly} universal weasels
comprise the maximal elements
in the ordering $<_*$.

Open Question Is every weakly
universal weasel universal?