

LaGO - A solver for mixed integer nonlinear programming

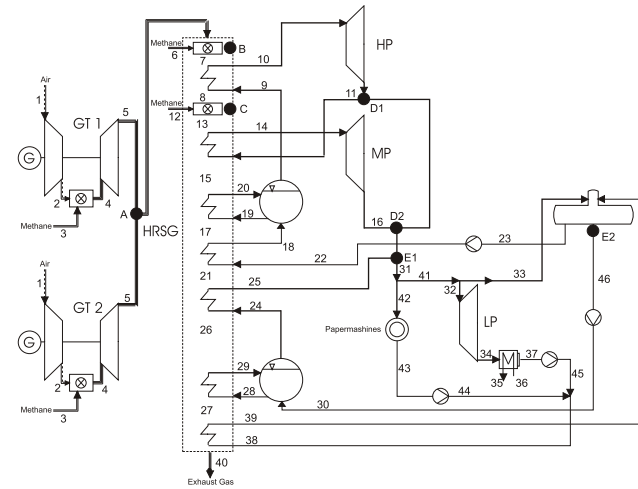
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June 1 2005

Problem formulation

MINLP:

$$\begin{aligned} \min \quad & f(x, y) \\ \text{s.t.} \quad & g(x, y) \leq 0 \\ & h(x, y) = 0 \\ & x \in [\underline{x}, \bar{x}] \\ & y \in [\underline{y}, \bar{y}] \text{ integer} \end{aligned}$$



MINLP:

- $n \ll 5.000$
- large problems are **structured**
- many possible applications
- in contrast to MIP, not used very much in practice

production planning, man power planing, scheduling, blending, refinery optimization, process design, engineering design, investment/de-investment, network design, financial optimization

Algorithmic overview of LaGO

- First branch-cut-and-price system for MINLP
- Deformation, rounding and Lagrange heuristics
- Block-separable reformulation:
$$\min\{c^T x \mid Ax + b \leq 0, x_{J_k} \in G_k, k = 1, \dots, p\}$$
- Convex and polyhedral relaxations:
replace G_k by a convex set or polyhedron $\hat{G}_k \supseteq G_k$
Optimal relaxations by solving dual problems
 $|J_k|$ influences the quality and computational cost of a relaxation/lower bound

Convex relaxations using underestimators

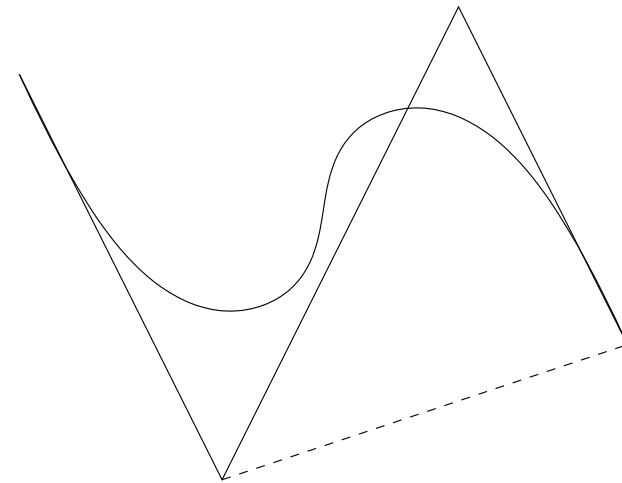
Bézier-underestimators [No96]:

- Bézier representation:

$$p(x) = \sum_{i=0}^l a_i x^i = \sum_{i=0}^l b_i \cdot B_i(x)$$

- Convex hull property:

$$p(x) \in \text{conv}\{b_i\}$$

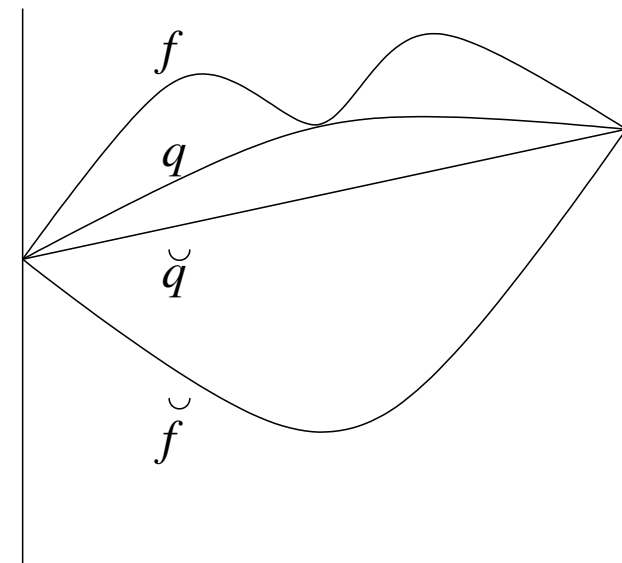


α -underestimators (sampling-technique) [NoAIVi03]

- Underestimation by a nonconvex quadratic form q

- $\check{q}(x) = q(x) + \alpha(x - \underline{x})^T(\bar{x} - x)$

- \check{q} often better than \check{f}



Both methods produce consistent bounds.

Optimal convex relaxation of block-separable quadratic problems

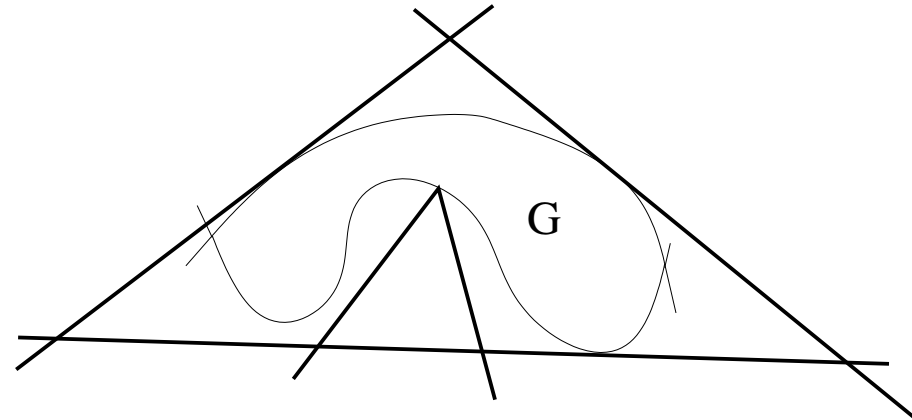
[No03]

- Reformulation (Q) of the original problem (P) by elimination of linear terms
- Formulation of dual problem to (Q) as an **eigenvalue optimization problem**:
$$\max_{\mu} \min_x L(x; \mu) = \max_{\mu} \sum_k \lambda_1(A_k(\mu)) + c(\mu)$$
- Solution of the eigenvalue optimization problem using a subgradient or bundle method
- Proof of the **dual equivalence** of (P) and (Q)

Polyhedral relaxation I

- **Linearization** (if g_i is convex):
$$g_i(\hat{x}) + \nabla g_i(\hat{x})^T (x - \hat{x}) \leq 0$$
- **Knapsack-cut** (by solving a separation problem):
$$b^T x \geq \underline{b} = \min_{x \in Z} b^T x,$$
$$Z = G \quad \text{or}$$
$$Z = \{x \in X \mid g_i(x) \leq 0\}$$
- **Interval gradient cut (new)**:
$$g_i(\hat{x}) + \min_{d \in [\underline{d}, \bar{d}]} d^T (x - \hat{x}) \leq 0$$

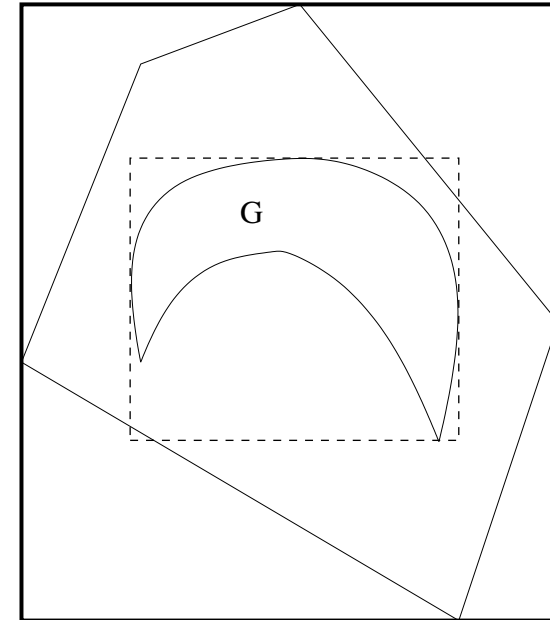
where $\nabla g_i(x) \in [\underline{d}, \bar{d}]$ für alle $x \in [\underline{x}, \bar{x}]$



Nonconvex polyhedral outer approximation $\hat{G} \supset G$

Polyhedral relaxation II

- **Level-cut:** $c^T x \leq \bar{v}$, where $\bar{v} \geq \text{val}(\text{MINLP})$
- **Box-reduction (constraint propagation)** improves the quality of cuts



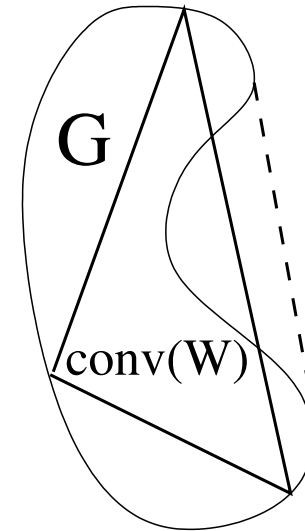
A polyhedral relaxation can also be used for sensitivity analysis and multicriterial optimization

Optimal polyhedral relaxations using column-generation

Find inner approximation points

$W = \{w_1, w_2, \dots\} \subset \text{conv}(G)$ such that

$$\begin{aligned} \min\{c^T x \mid Ax + b \leq 0, x \in \text{conv}(W)\} = \\ \min\{c^T x \mid Ax + b \leq 0, x \in \text{conv}(G)\} \end{aligned}$$



Algorithm:

1. Solve the restricted master problem (RMP):

$$\hat{\mu} = \operatorname{argmax}_{\mu} \min_x \{c^T x + \mu^T (Ax + b) \mid x \in \text{conv}(W)\}$$

2. Solve the Lagrange problem: $w = \operatorname{argmin}\{(c^T + \hat{\mu}^T A)x \mid x \in G\}$,
set $W = W \cup \{w\}$ and add new columns to RMP.

Equivalent to dual cutting plane method and therefore convergent

Remarks

- Lagrange problem decomposes into **sub-problems**, which can be solved by an arbitrary **global solver** (branch+cut, populations heuristic).
- Each solution of a sub-problem generates a **Lagrangian cut**, which is added to the outer approximation
- By **comparing the outer and inner approximation** it can be determined how good the approximation of a sub-problem is.
- In contrast to bundle methods, relaxations of stochastic programs and optimal control problems can be **updated efficiently** after refining **szenarios or grids** respectively.

Heuristics

1. **Deformation heuristic** by successively solving:

$$\min\{(1 - t_k)\check{f}(x) + t_k f(x) \mid (1 - t_k)\check{g}(x) + t_k g(x) \leq 0\}$$

2. **Rounding heuristic** by rounding some components of the solution of an inner or outer approximation and backtracking
3. **Lagrange heuristic** by combining **several** inner approximation points $x_{J_k} \in W_k$, where $c^T x + \delta \|Ax + b\|_+$ is small, and projection onto $\{x \mid Ax + b \leq 0, x_B \text{ binary}\}$

Branch-Cut-and-Price

Lower bounds:

(B1) dual bounds: column generation (BCP) or eigenvalue optimization

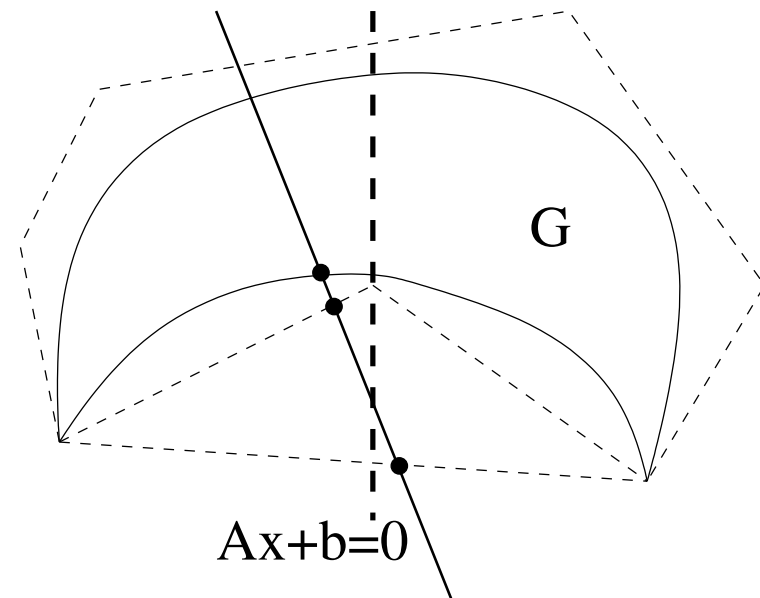
(B2) LP-bounds (with Knapsack and linearization cuts)

(B3) by a convex relaxation

All bounds are **consistent** and we have $(B1) \geq (B2) \geq (B3)$

Upper bounds by heuristics

Several **branching strategies**

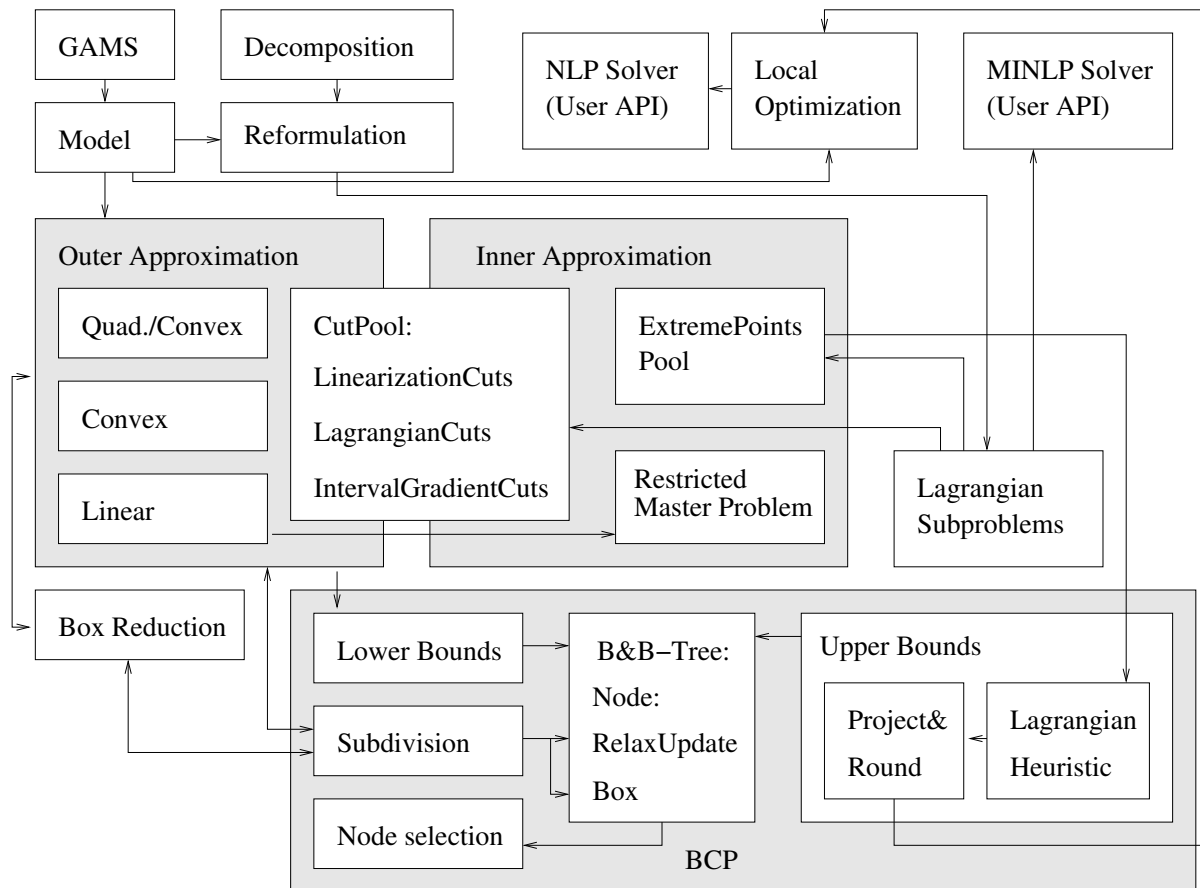


C++ Library

Together with Stefan Vigerske, more than 40000 lines of code

Object oriented design similar to COIN/BCP and ABACUS

Interfaces: GAMS, AMPL, COIN, SNOPT, CPLEX, ARPACK, NOA, FILIB, TNT, METIS



Numerical experiments

- **MaxCut** experiments, $n \leq 1000$: performance similar to specialized solvers
- **GAMS-MinlpLIB** experiments, $n \leq 500$, **Comparison with BARON**, November 2004, BC-Algorithmus of LaGO using LP-bounds and BARON with default parameters:

	total number	obj LaGO better	obj equal	obj BARON better
LaGO much faster :	3	3	-	-
LaGO faster :	1	-	1	-
Both solvers same performance :	5	1	3	1
BARON faster :	9	-	8	1
BARON much faster :	28	5	17	6
Only BARON solution :	5	-	-	5
Both solvers no solution :	1	-	1	-
Number of problems :	52	9	30	13

DFG-Project

Design of Complex Energy Conversion Systems,
Technical University of Berlin (Institute for Energy Engineering) and
Humboldt-University Berlin (Department of Mathematics)

[Ahadi-Oskui, Alperin, Cziesla, Nowak, Tsatsaronis, 2001-2004]

MINLP, $n = 1300$: BCP and specialized heuristic

Found acceptable solution in reasonable time

BARON and SBB were not able to find a solution

Final Remarks

- <http://www.mathematik.hu-berlin.de/~eopt/LaGO/documentation/> .
- Book: [Relaxation and Decomposition Methods for Mixed-Integer Nonlinear Optimization](#)
Birkhäuser Verlag, to appear
- Possible improvements:
 - Reduction of duality gap by [nonconvex polyhedral inner approximation](#) (MIP master problem)
Li, D., Sun, X. L., Wang, J., and McKinnon, K. (2002). A convergent lagrangian and domain cut method for nonlinear knapsack problems. Technical report, SEEM2002-10, Department of Systems Engineering & Engineering Management, The Chinese University of Hong Kong.
 - consistent bounds and branching
- New DFG project planned
- Integration into GAMS planned